



Technische  
Universität  
Braunschweig

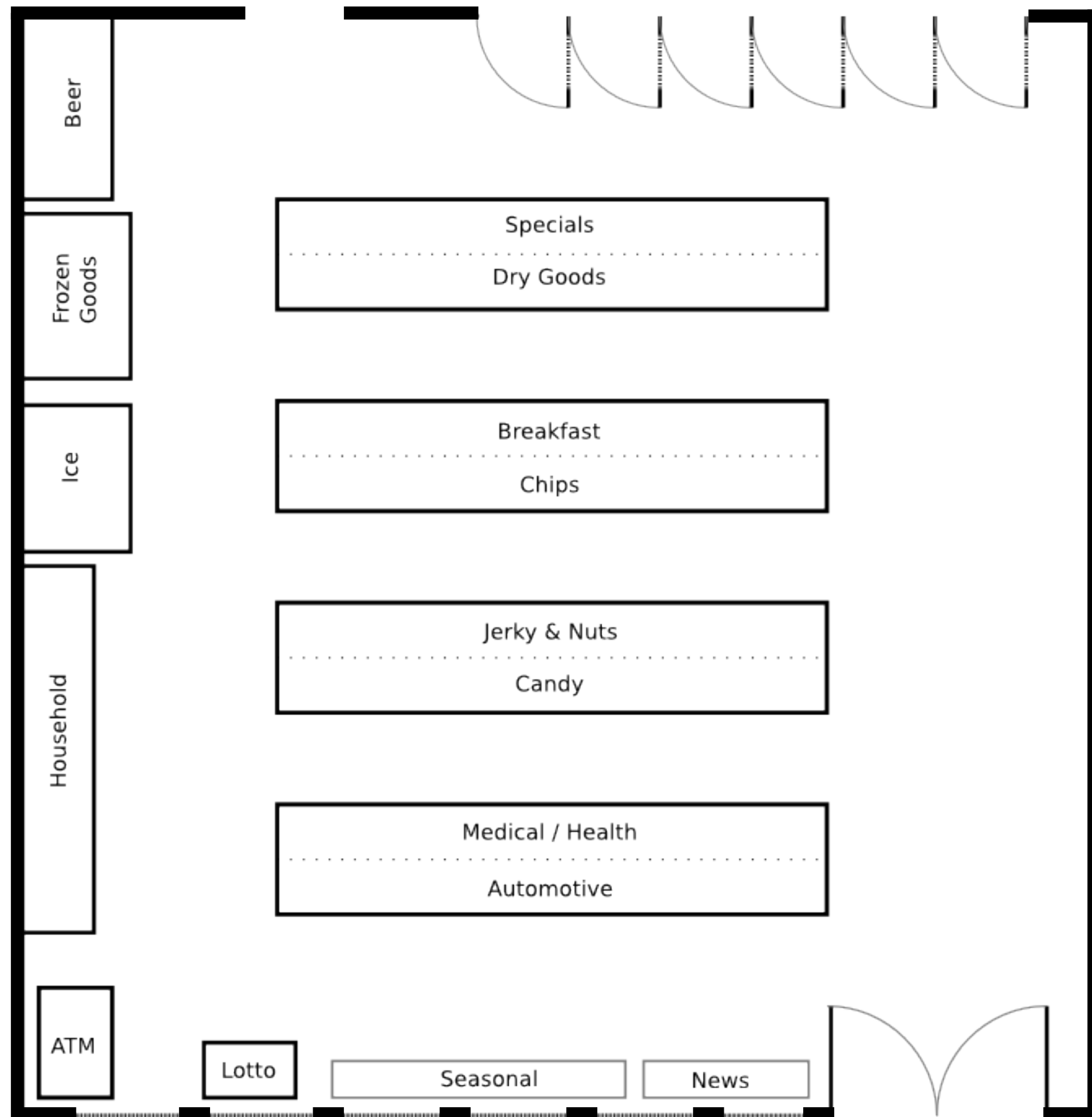


# New Variants of the Floodlight Problem

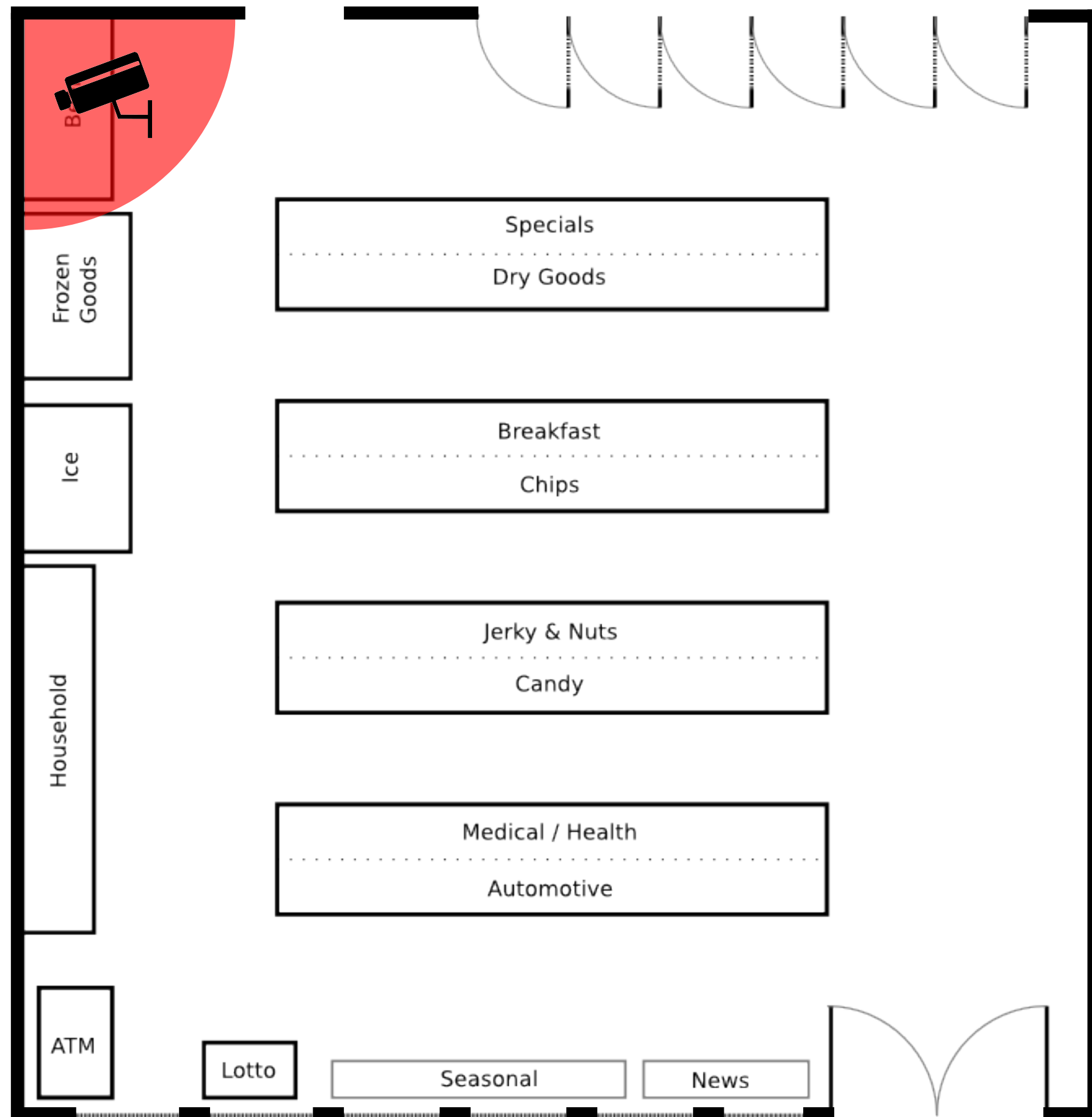
Bachelor's Thesis

Yannic Lieder, September 5, 2018

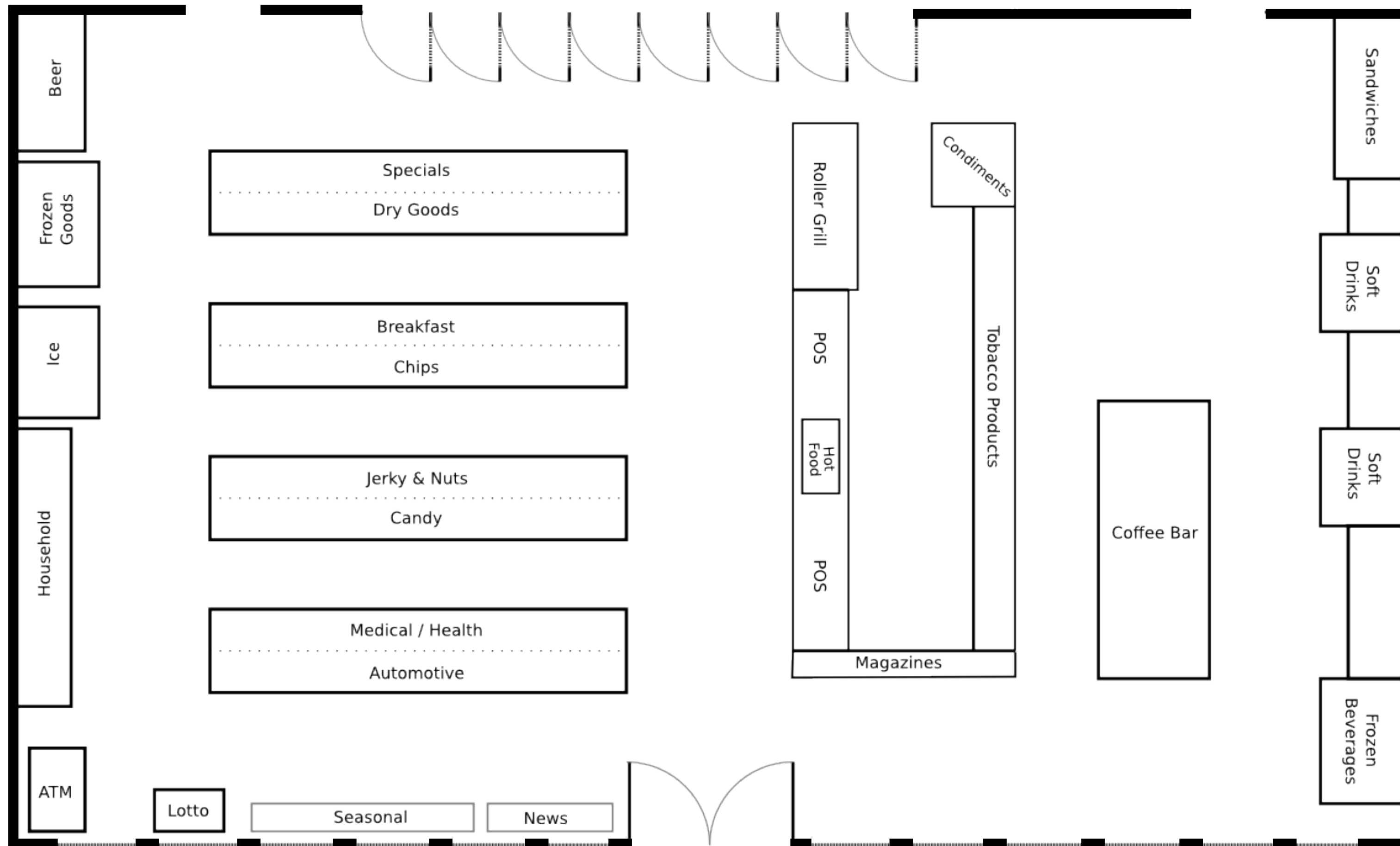
# Introduction



# Introduction

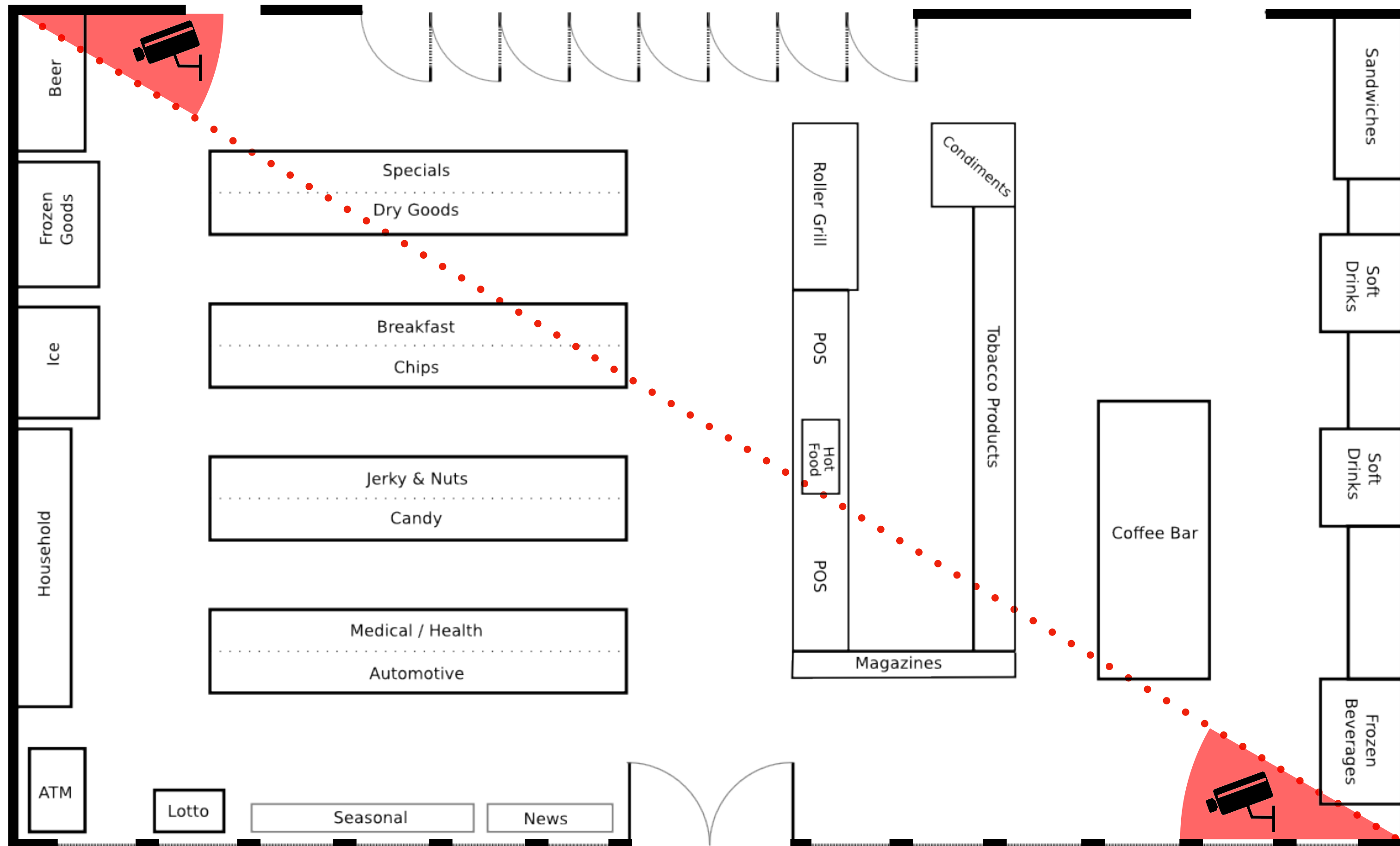


# Introduction





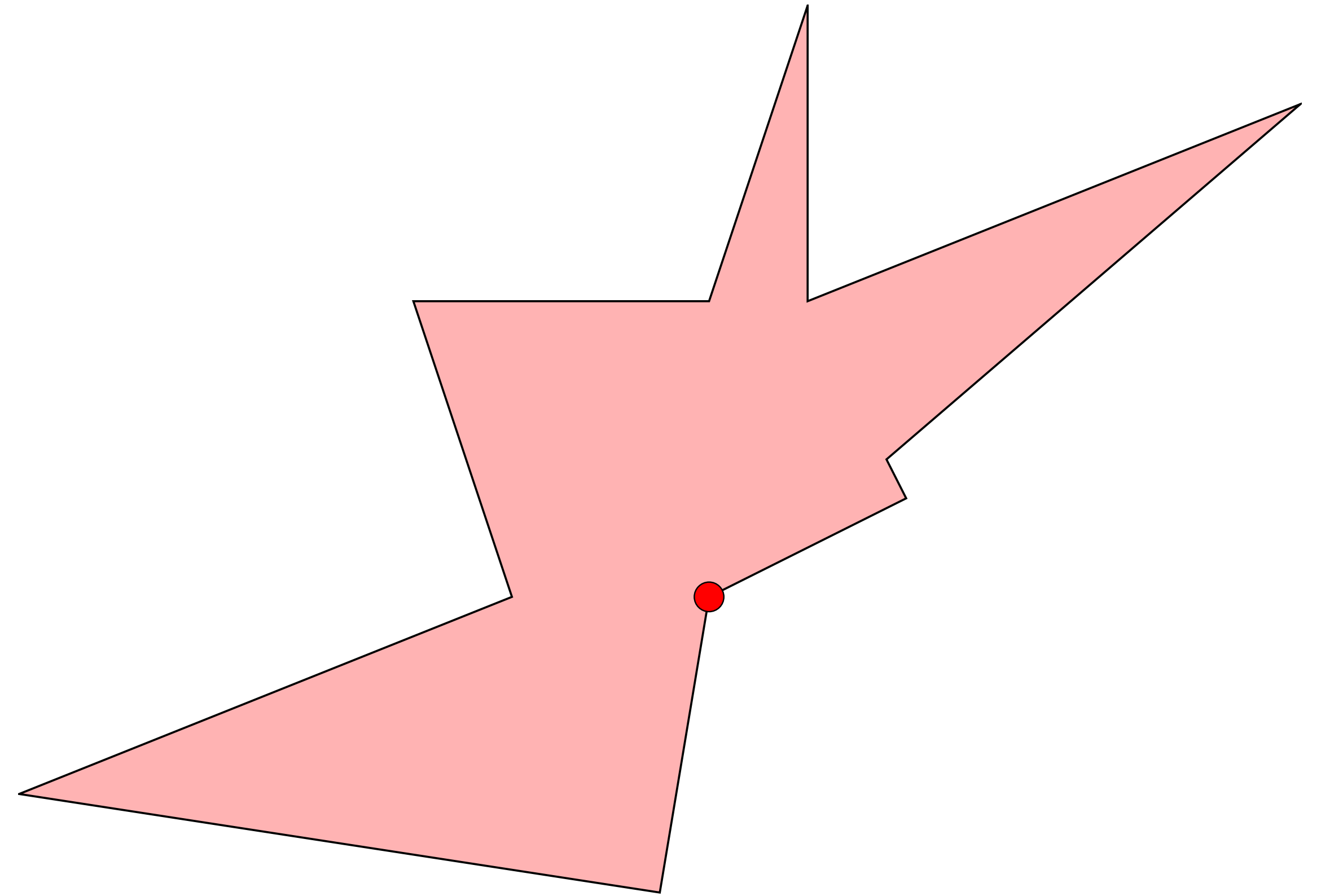
# Introduction



# Background

1973

The Art Gallery Problem



# Background

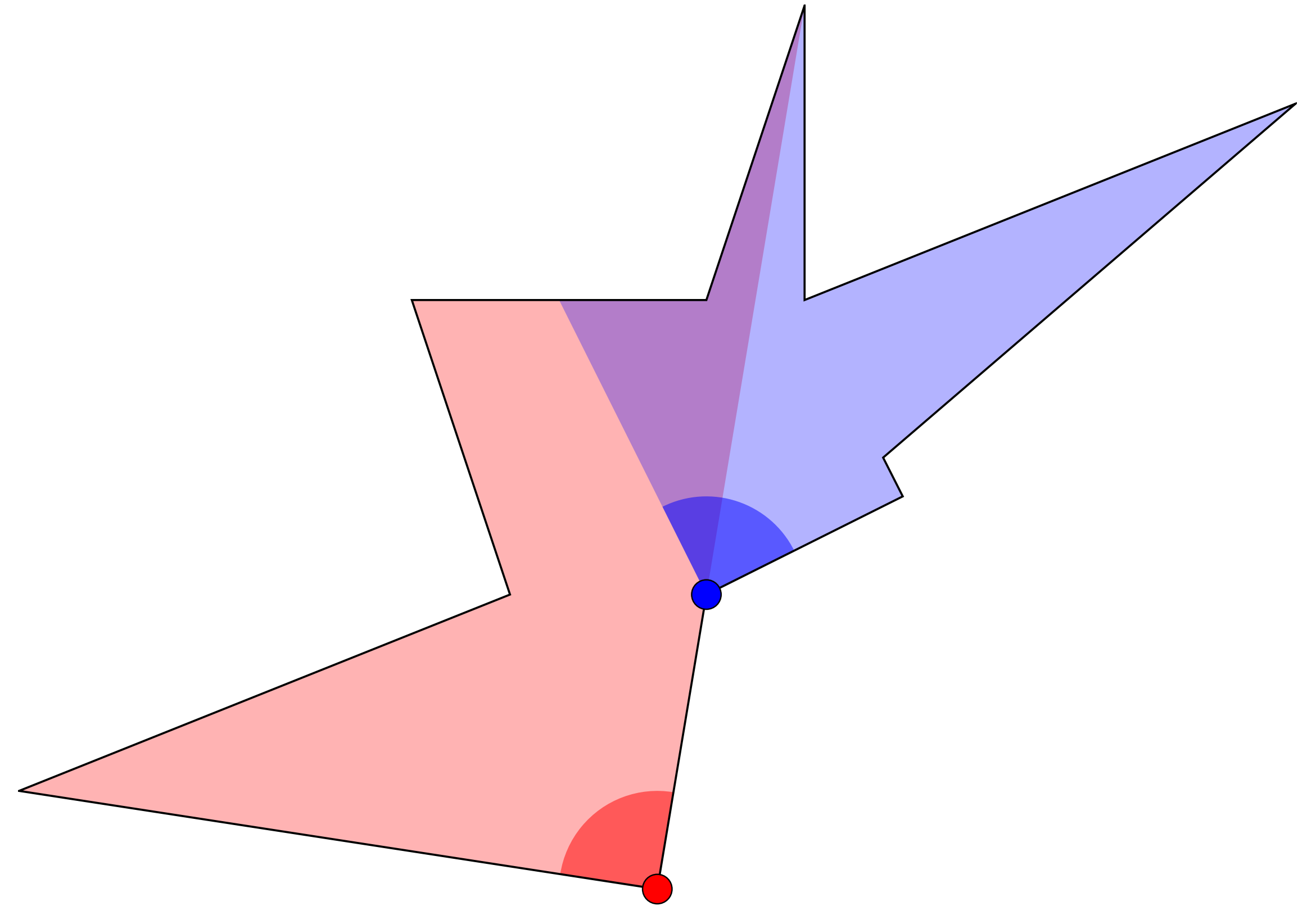
1973

The Art Gallery Problem



1992

The  $\alpha$ -Floodlight Problem



# Background

1973

The Art Gallery Problem



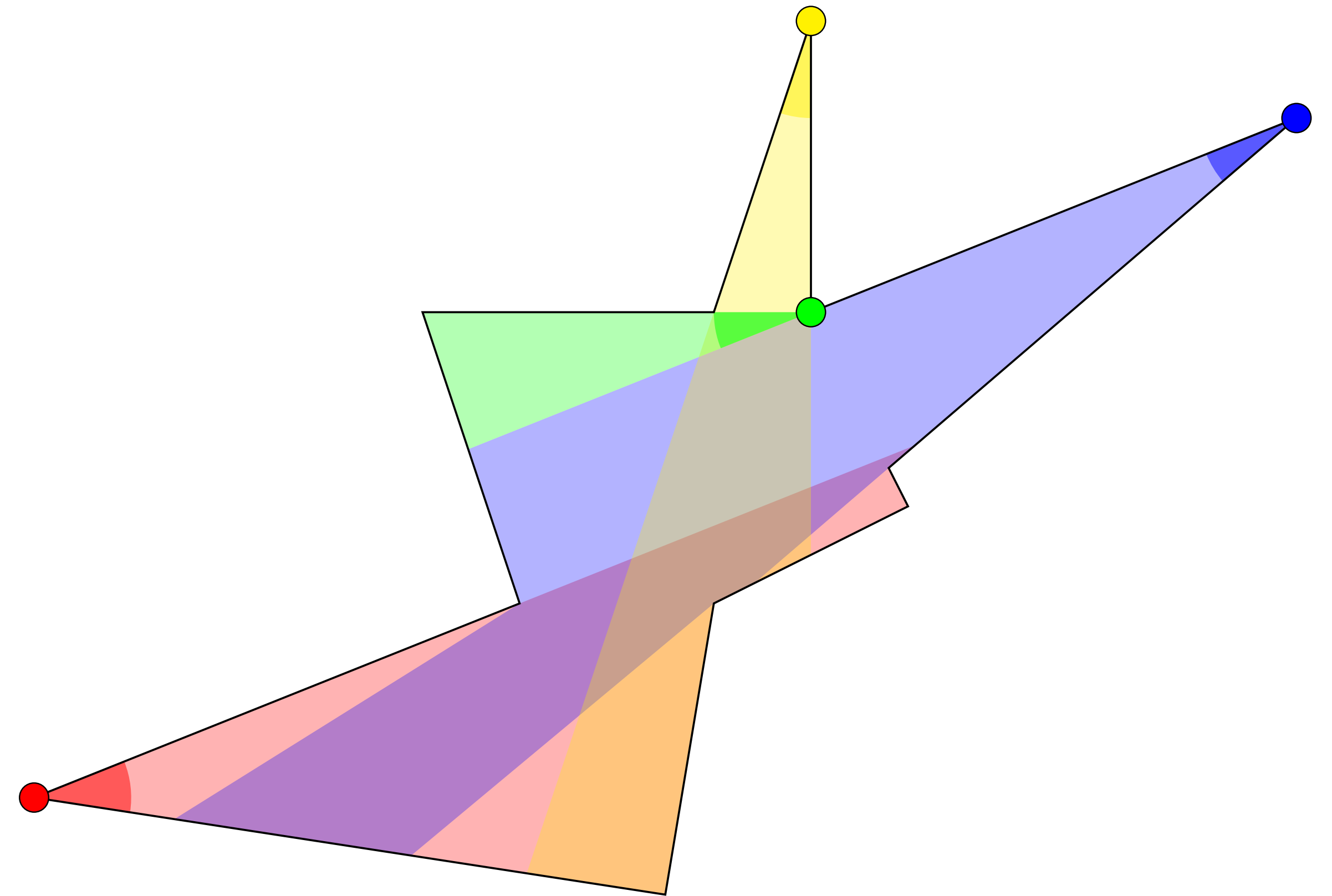
1992

The  $\alpha$ -Floodlight Problem



This  
thesis

The Angular Art Gallery Problem





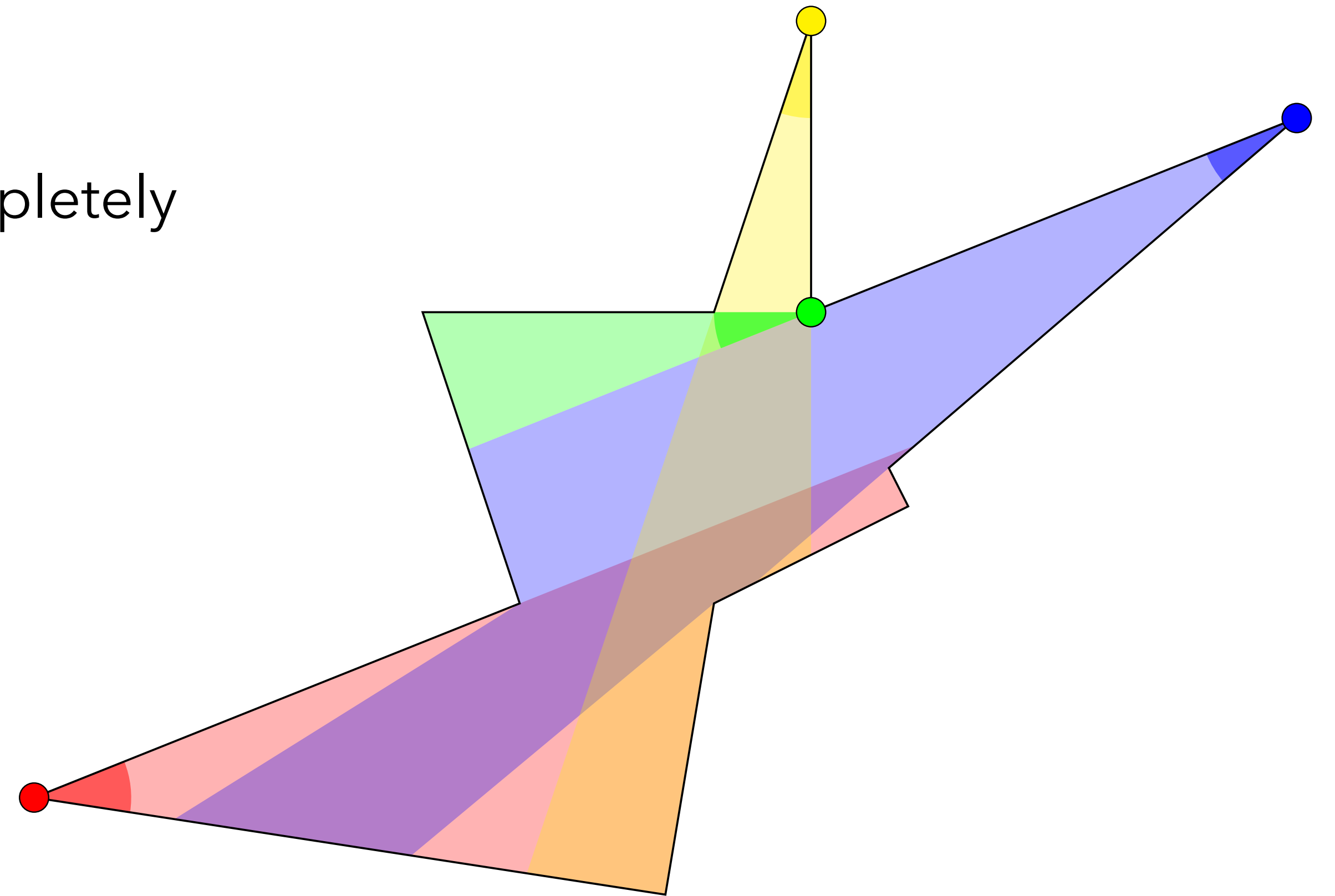
# Problem Definition

## Angular Art Gallery Problem

Instance: A simple polygon  $P$

Wanted: A set of floodlights, covering  $P$  completely

Minimize: The total angle of all floodlights



# Results

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Lower Bound

Upper Bound

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# Results

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	Lower Bound	Upper Bound
Equilateral Triangles	$\frac{\pi}{3}$	$\frac{\pi}{3}$

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	Lower Bound	Upper Bound
Equilateral Triangles	$\frac{\pi}{3}$	$\frac{\pi}{3}$
Histograms	$(n - 1)\frac{\pi}{6} - \epsilon$	$(n - 1)\frac{\pi}{6}$

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Simple Polygons	$(n - 1)\frac{\pi}{6} - \epsilon$	$(n - 2)\frac{\pi}{4}$

---



# Results

	Lower Bound	Upper Bound
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Duality to independent circle packing

# Outline

Introduction

Equilateral Triangles

Histograms

Simple Polygons

Duality

Conclusion

Introduction

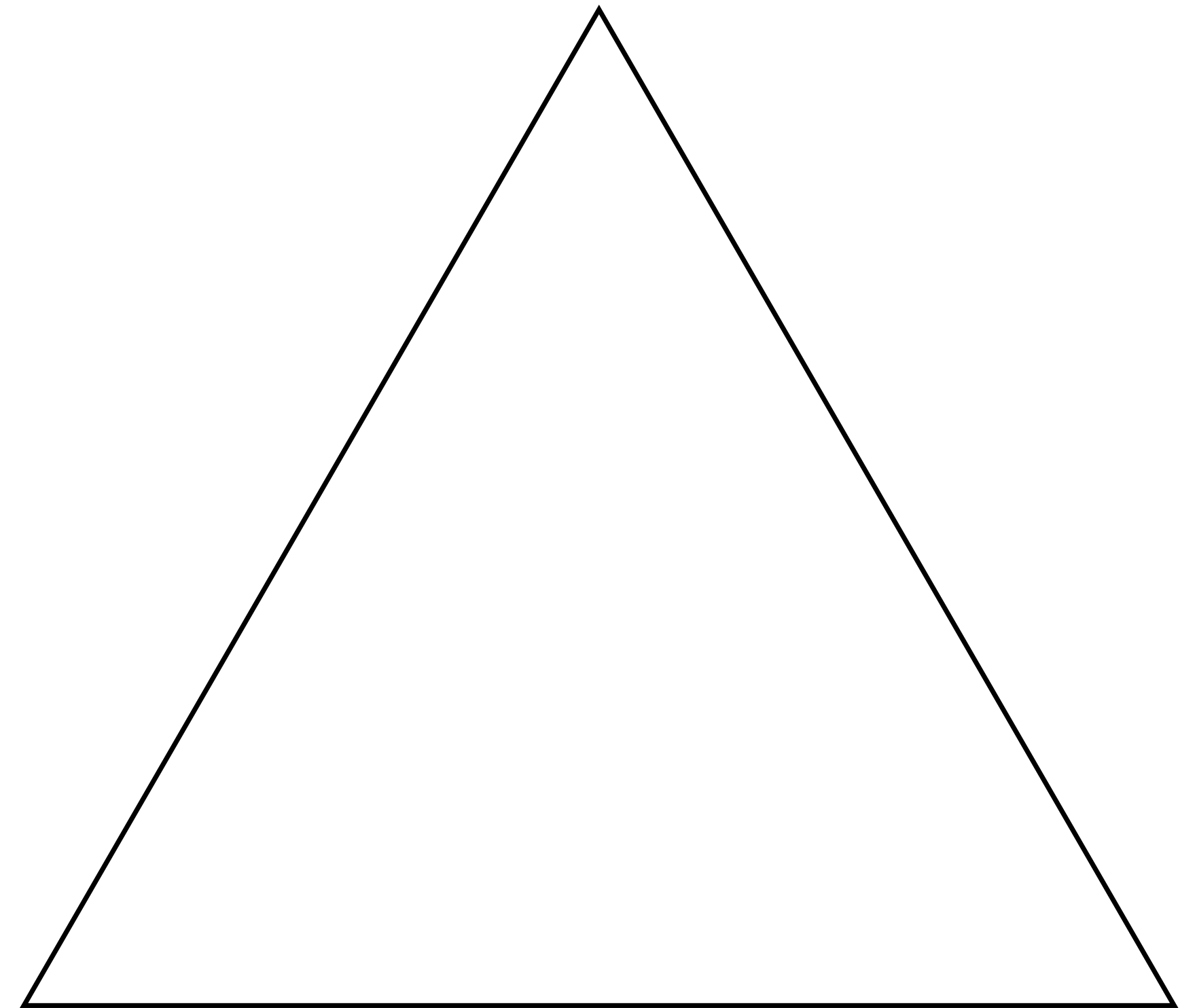
# Optimal Covering of **Equilateral Triangles**

Histograms

Simple Polygons

Duality

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Introduction

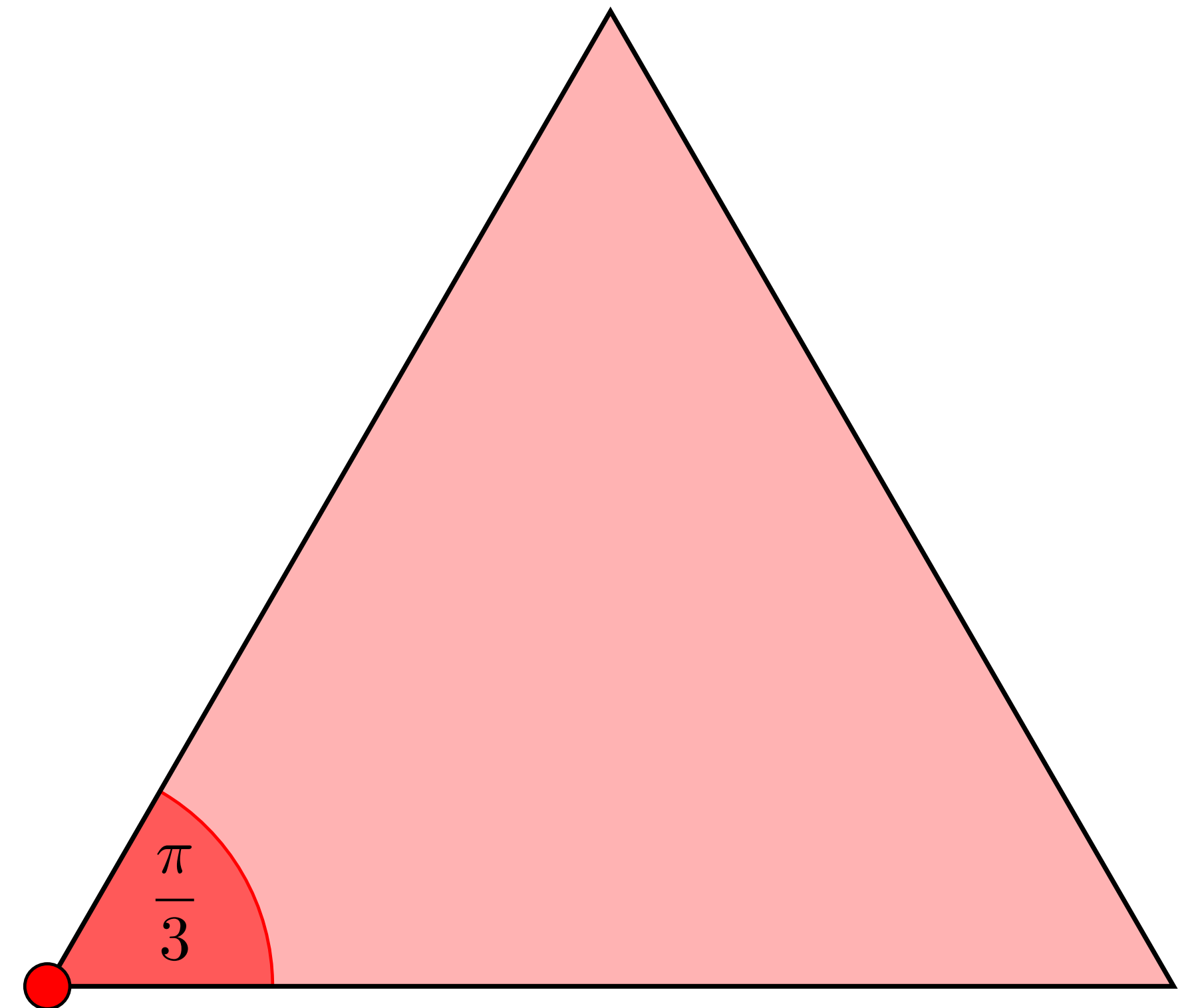
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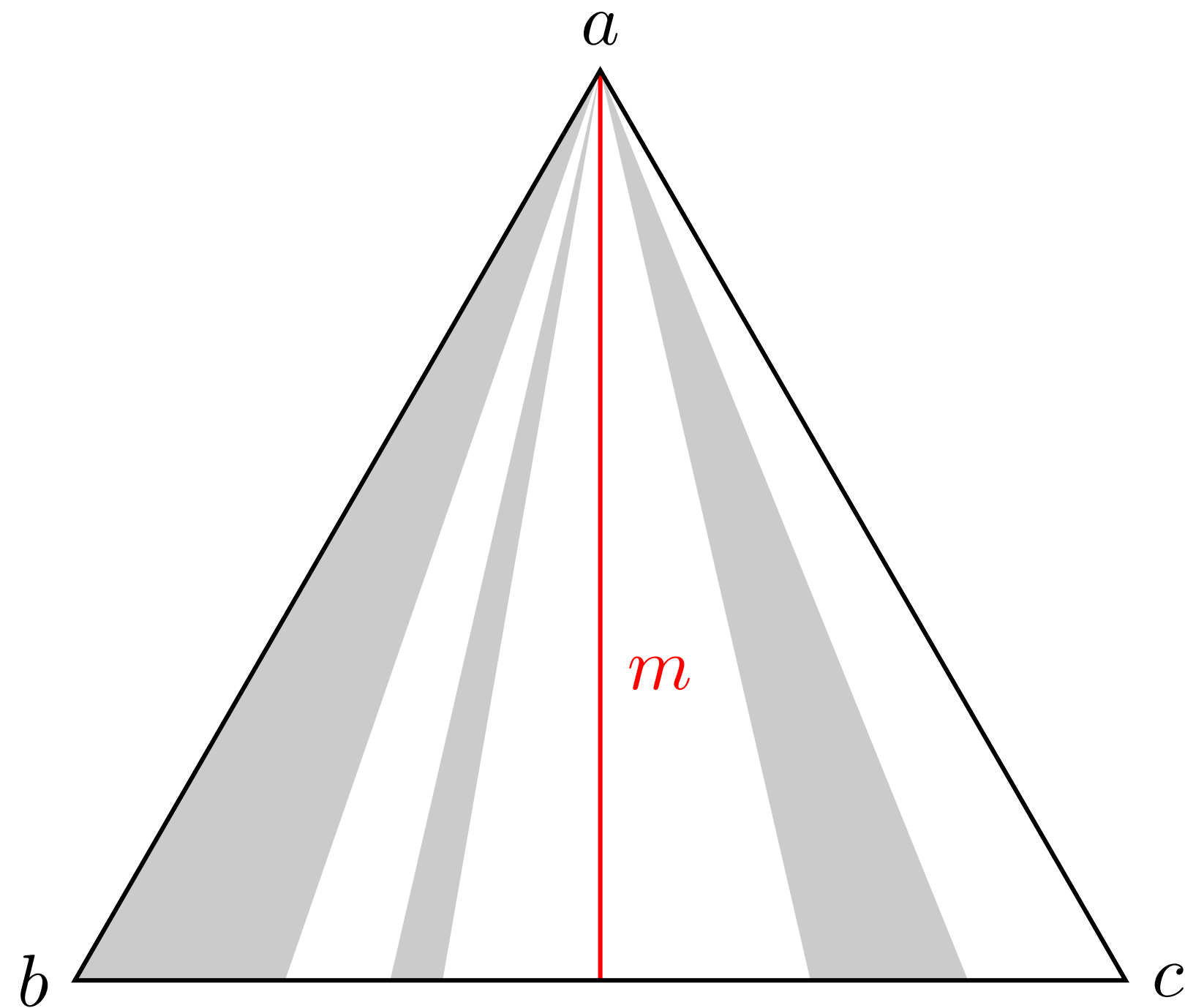
Conclusion



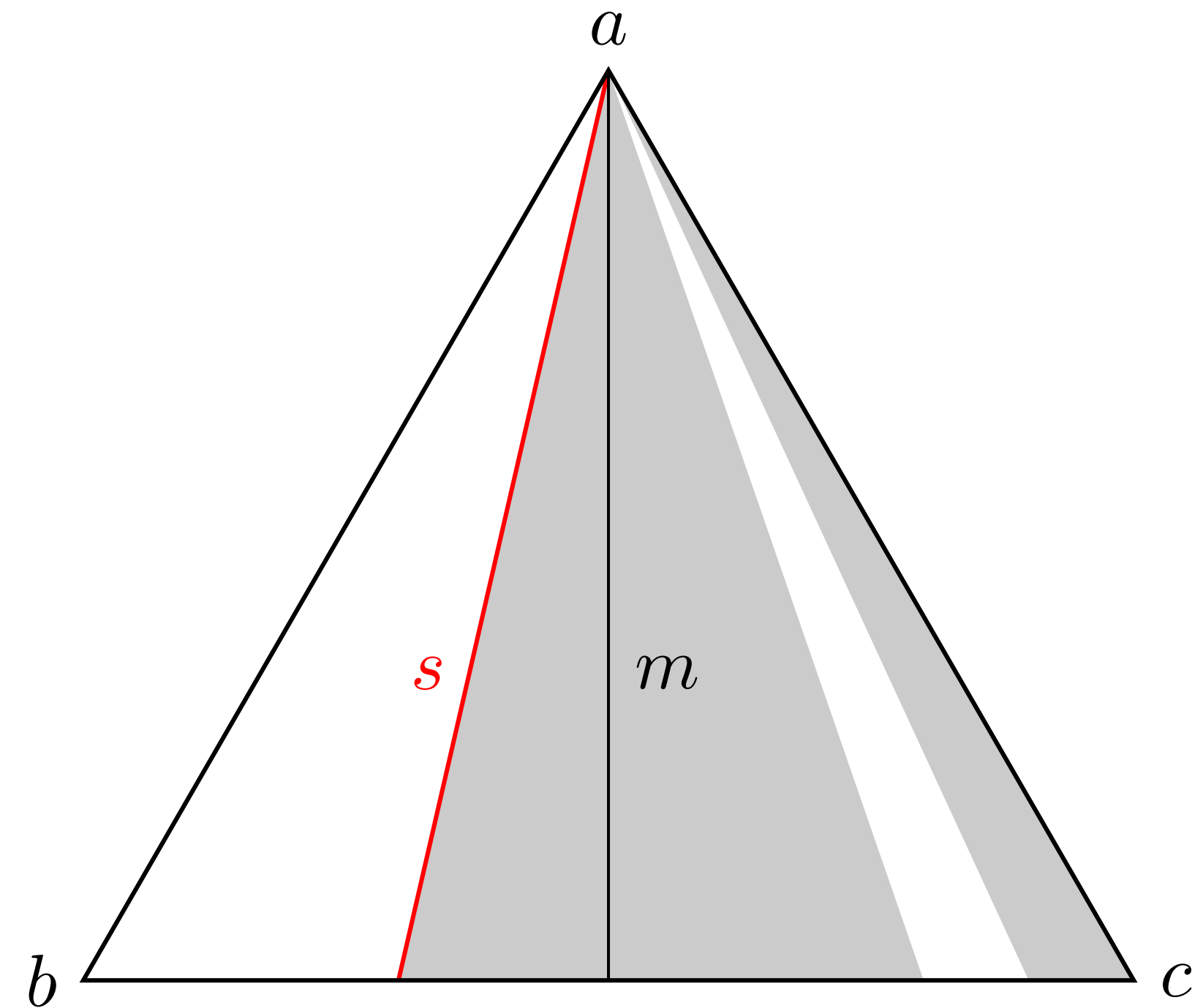


# Lower Bound for Equilateral Triangles

Case 1:  $m$  is not covered by  $a$

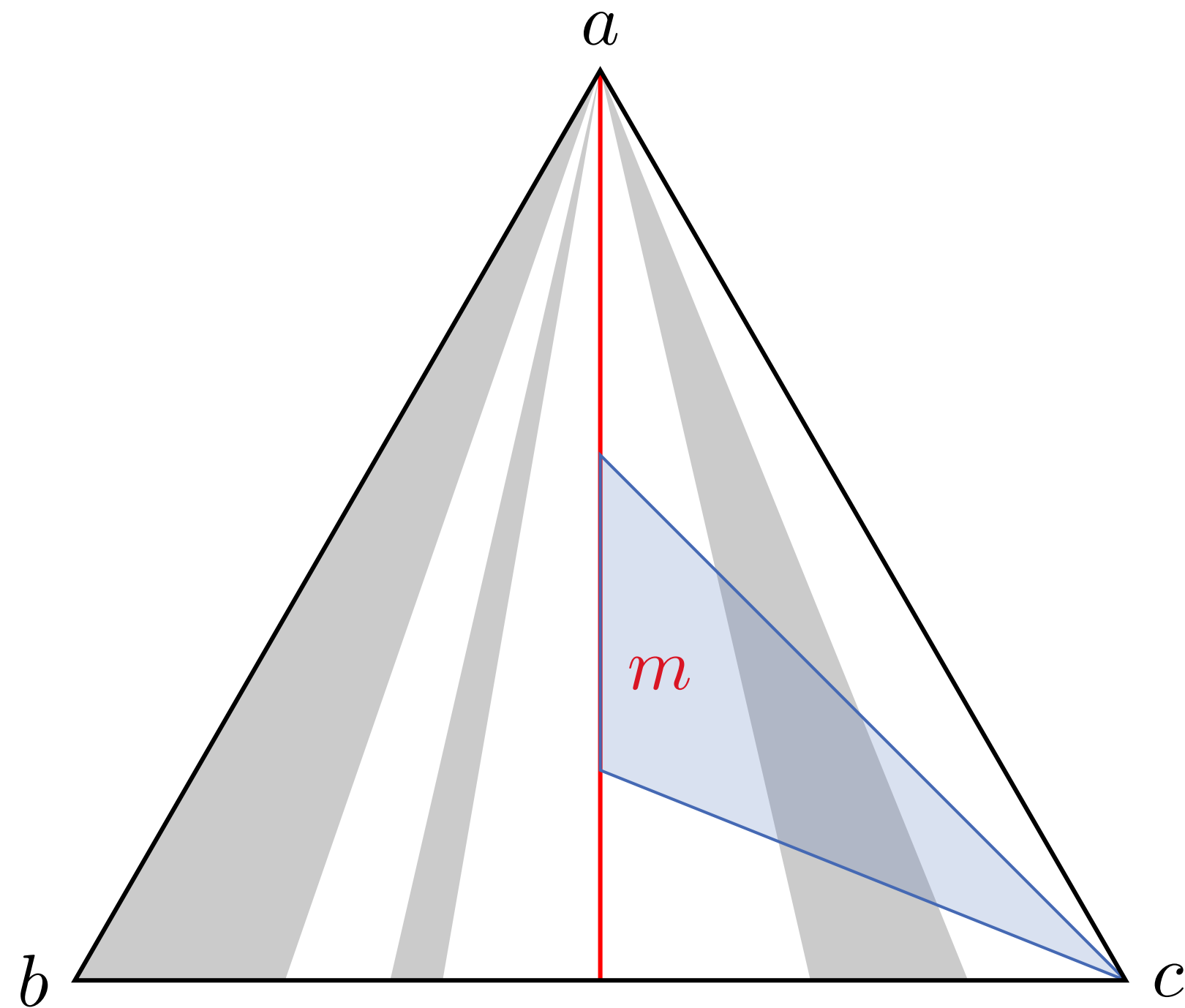


Case 2:  $m$  is covered by  $a$

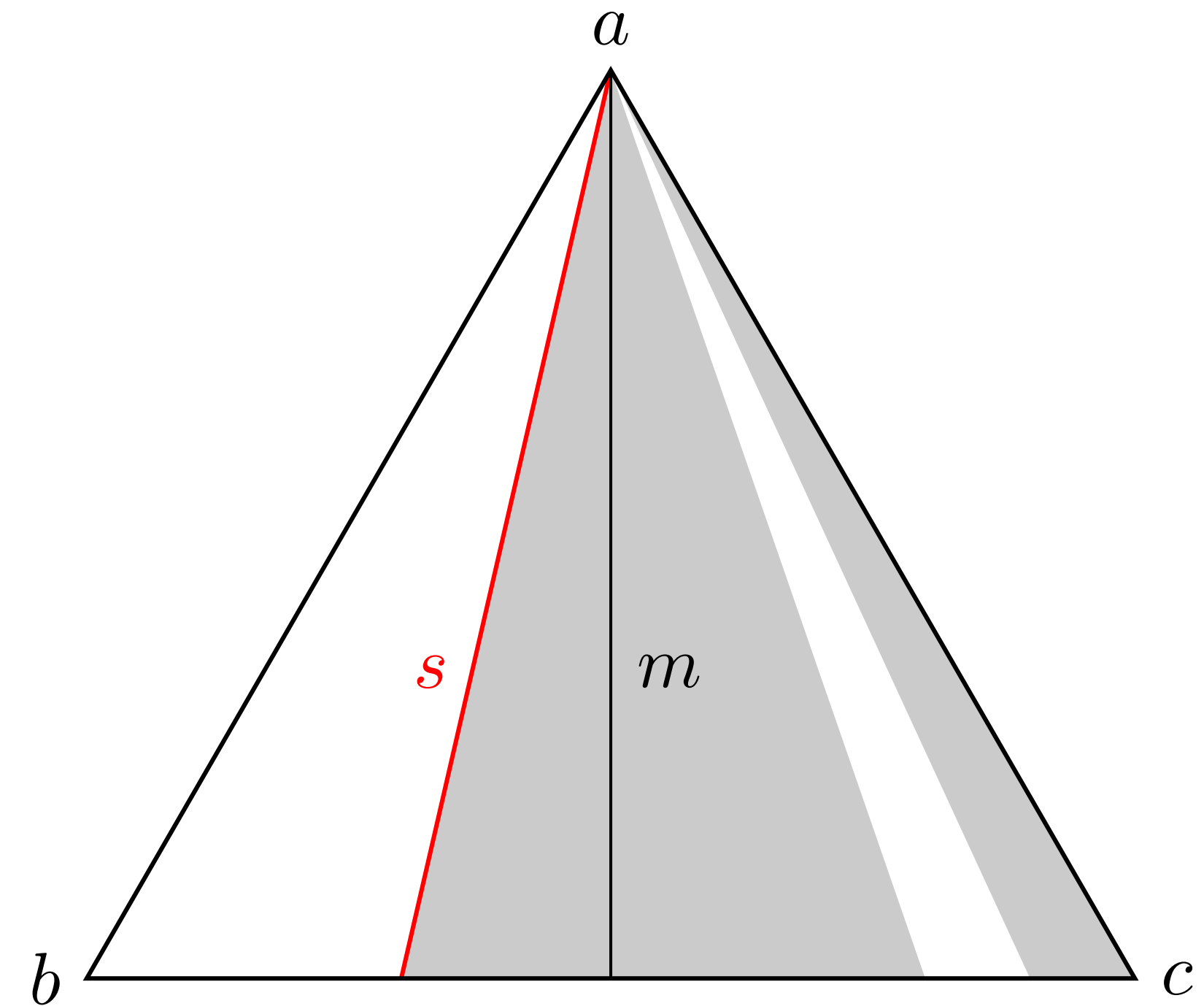


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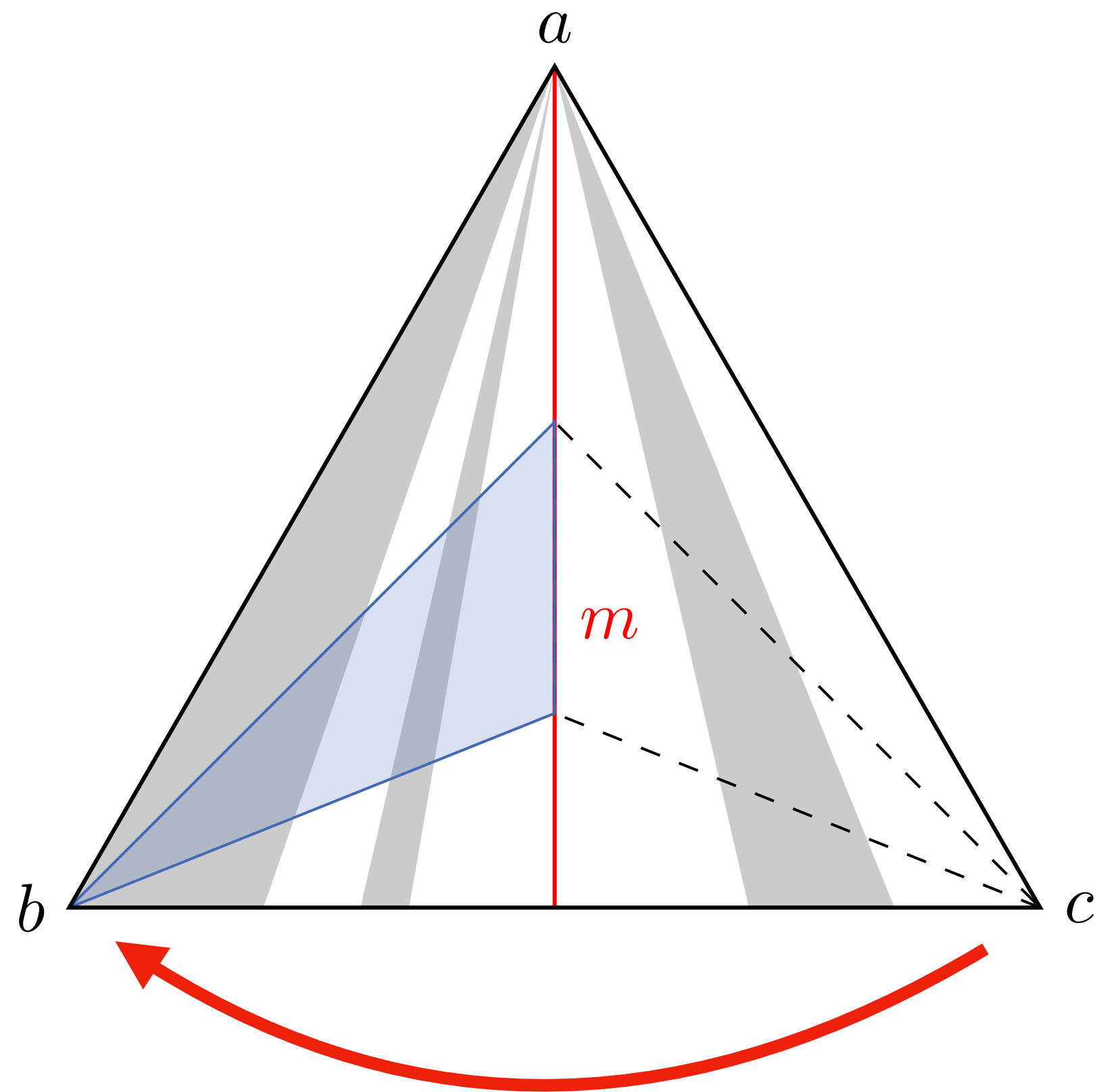


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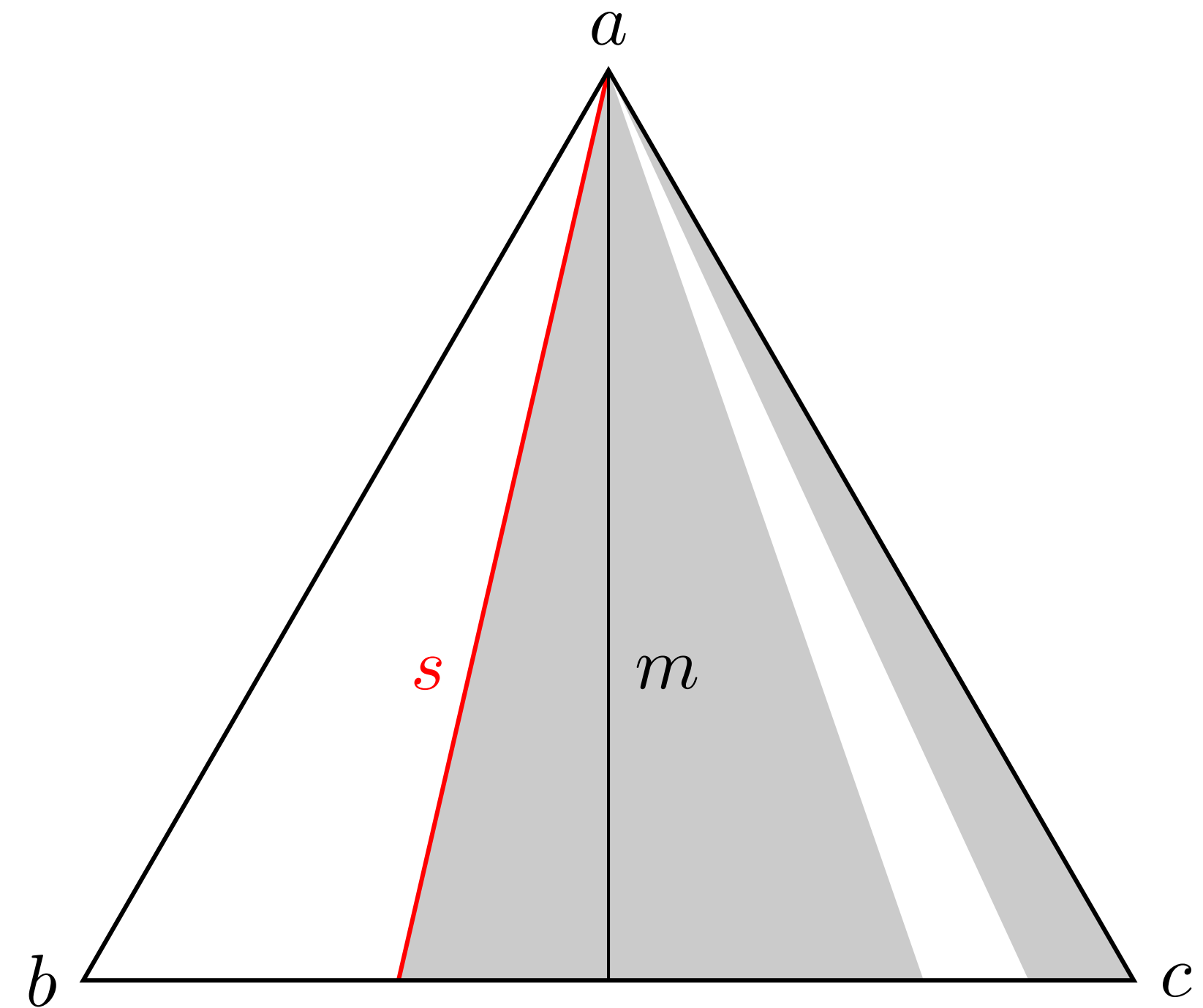


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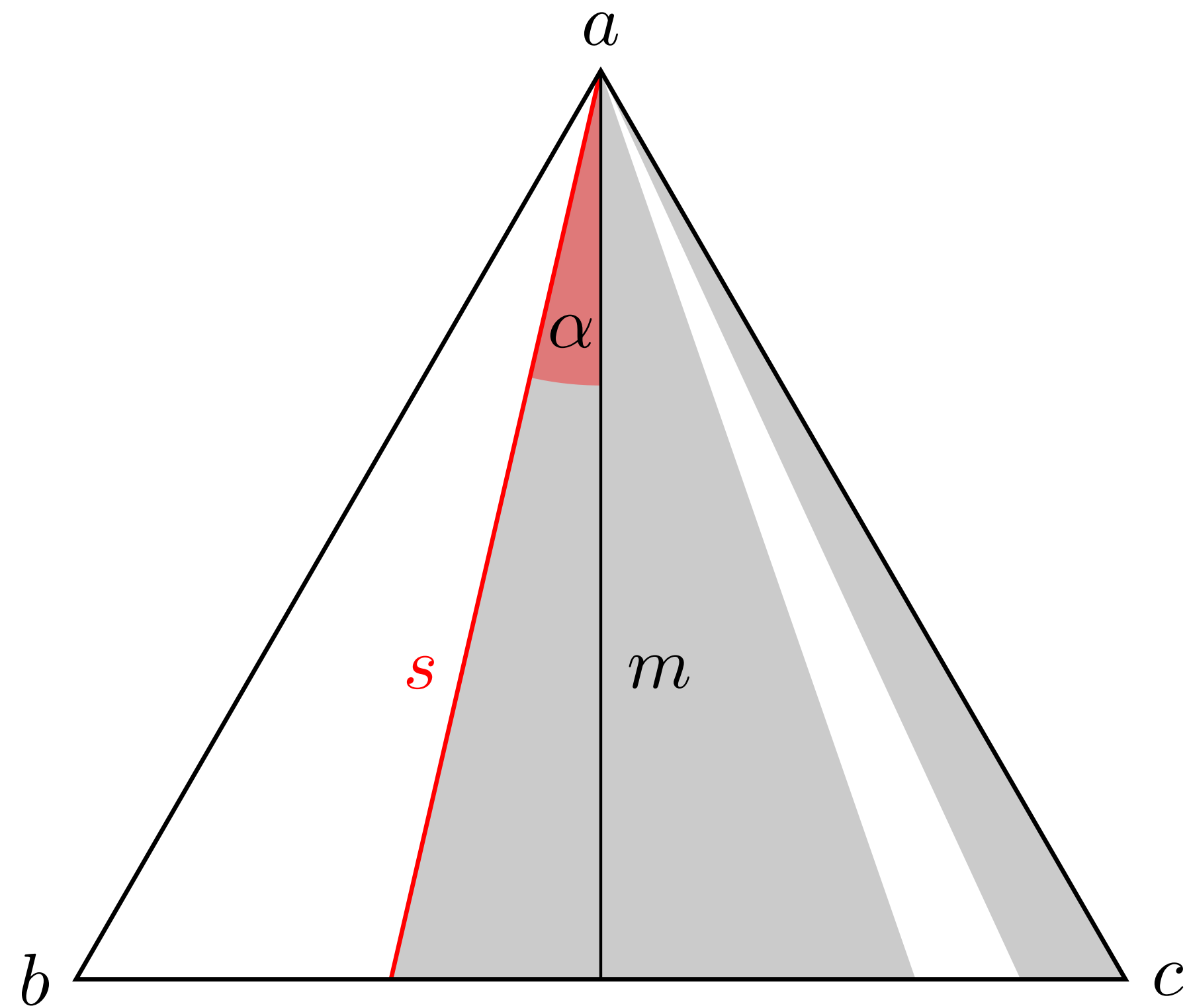
Case 1:  $m$  is not covered by  $a$



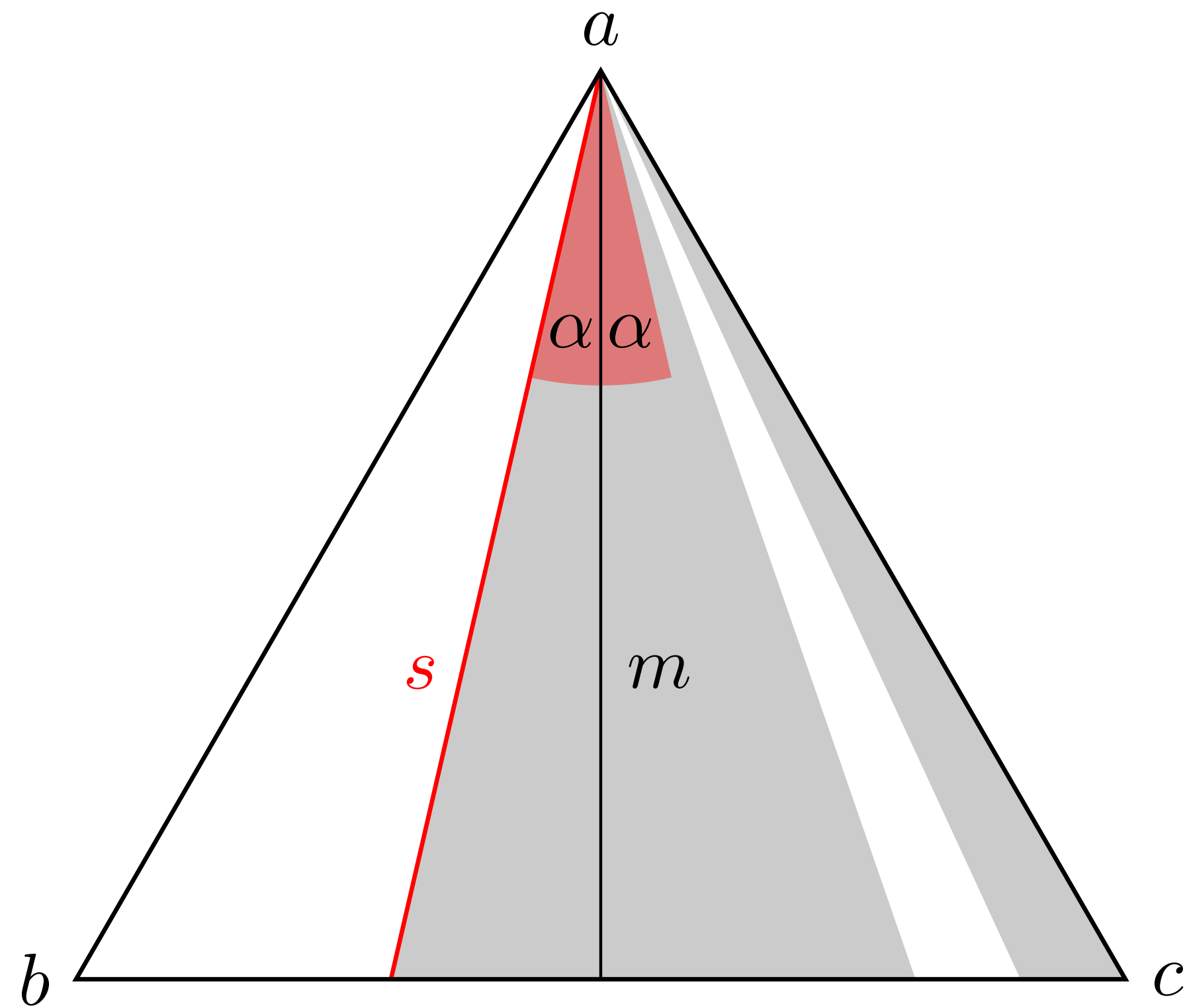
Case 2:  $m$  is covered by  $a$



# Lower Bound for Equilateral Triangles, case 2

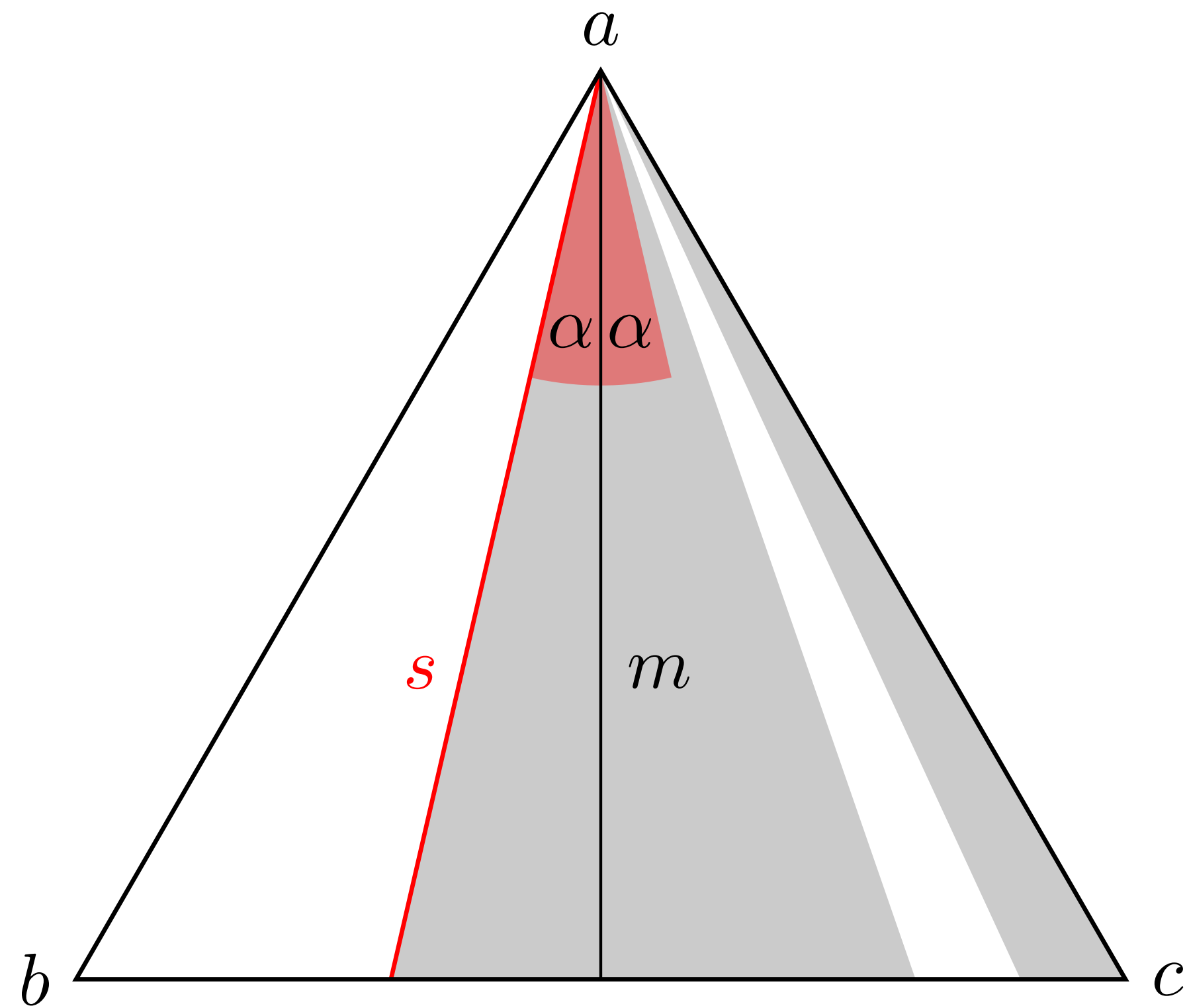


# Lower Bound for Equilateral Triangles, case 2





# Lower Bound for Equilateral Triangles, case 2

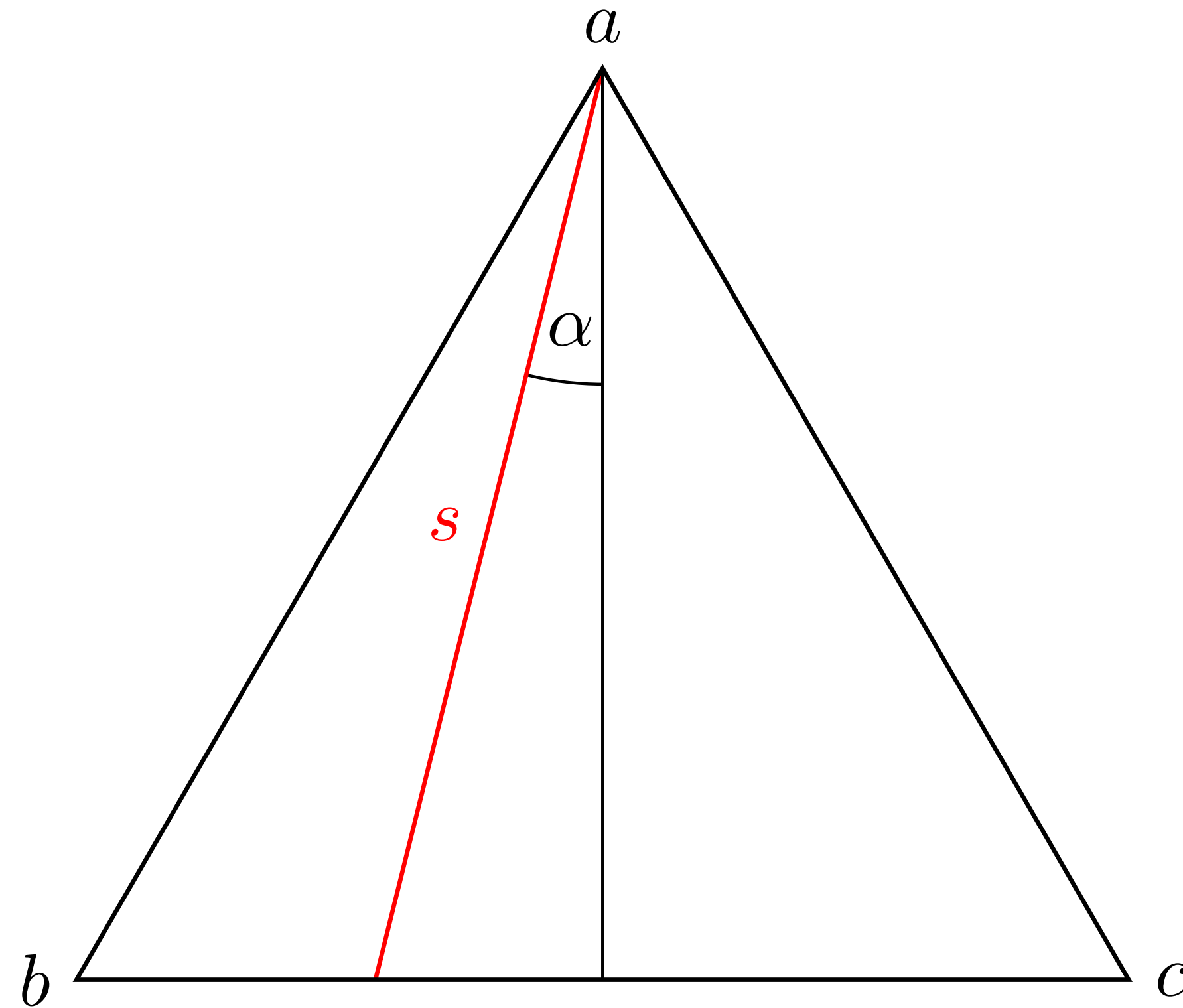


Idea: Showing that an angle of at least

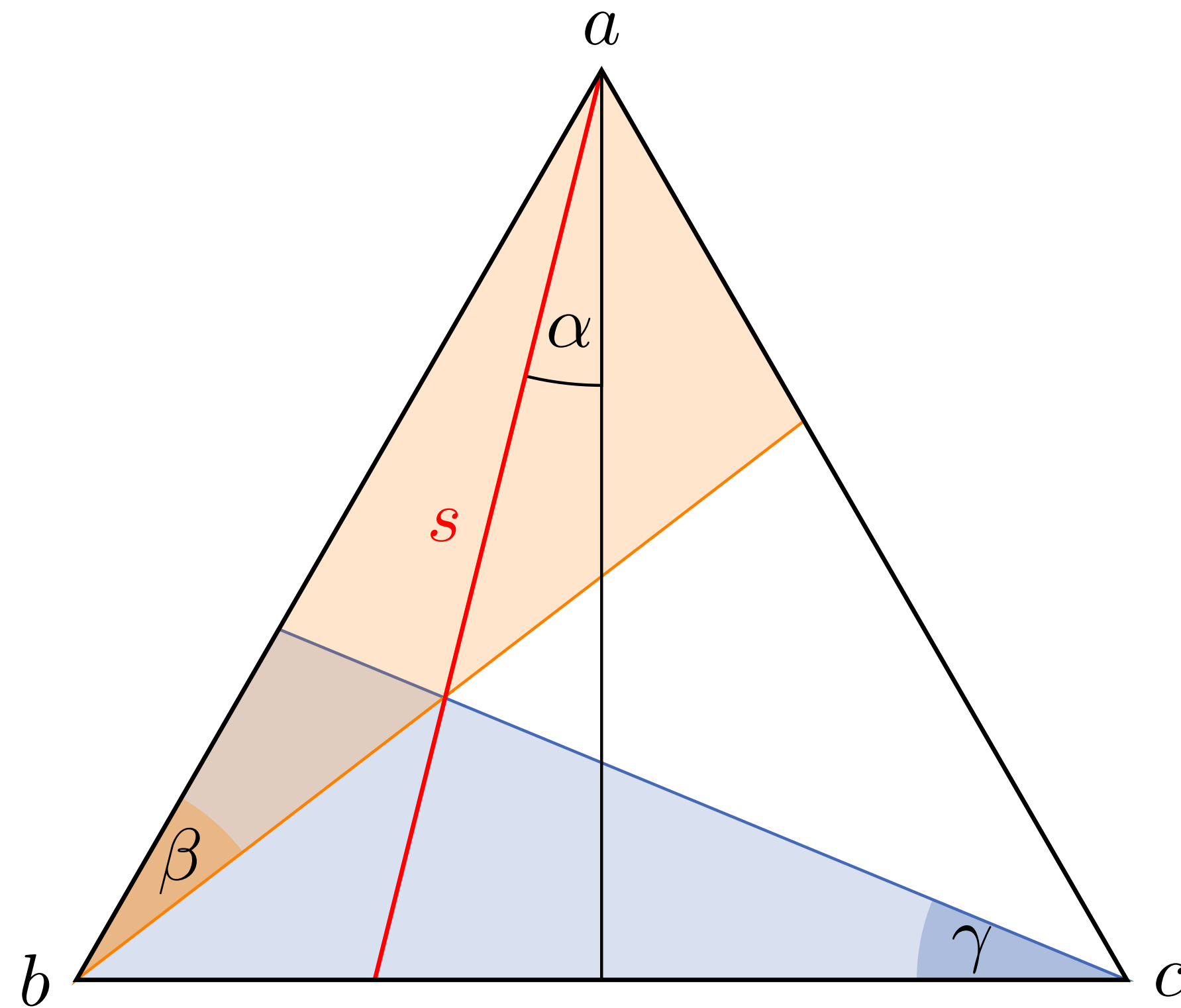
$$\frac{\pi}{3} - 2\alpha$$

is required to cover  $s$ .

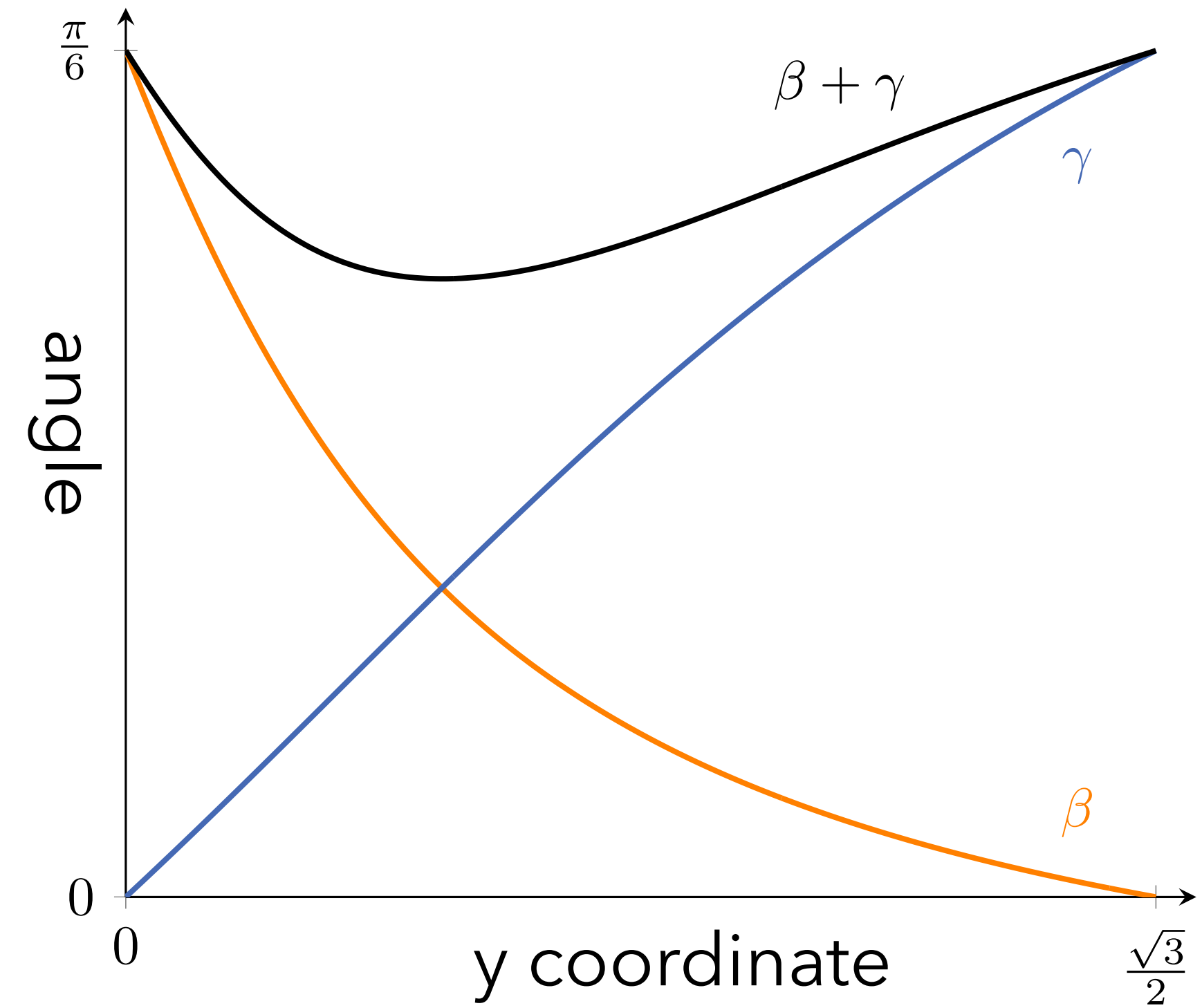
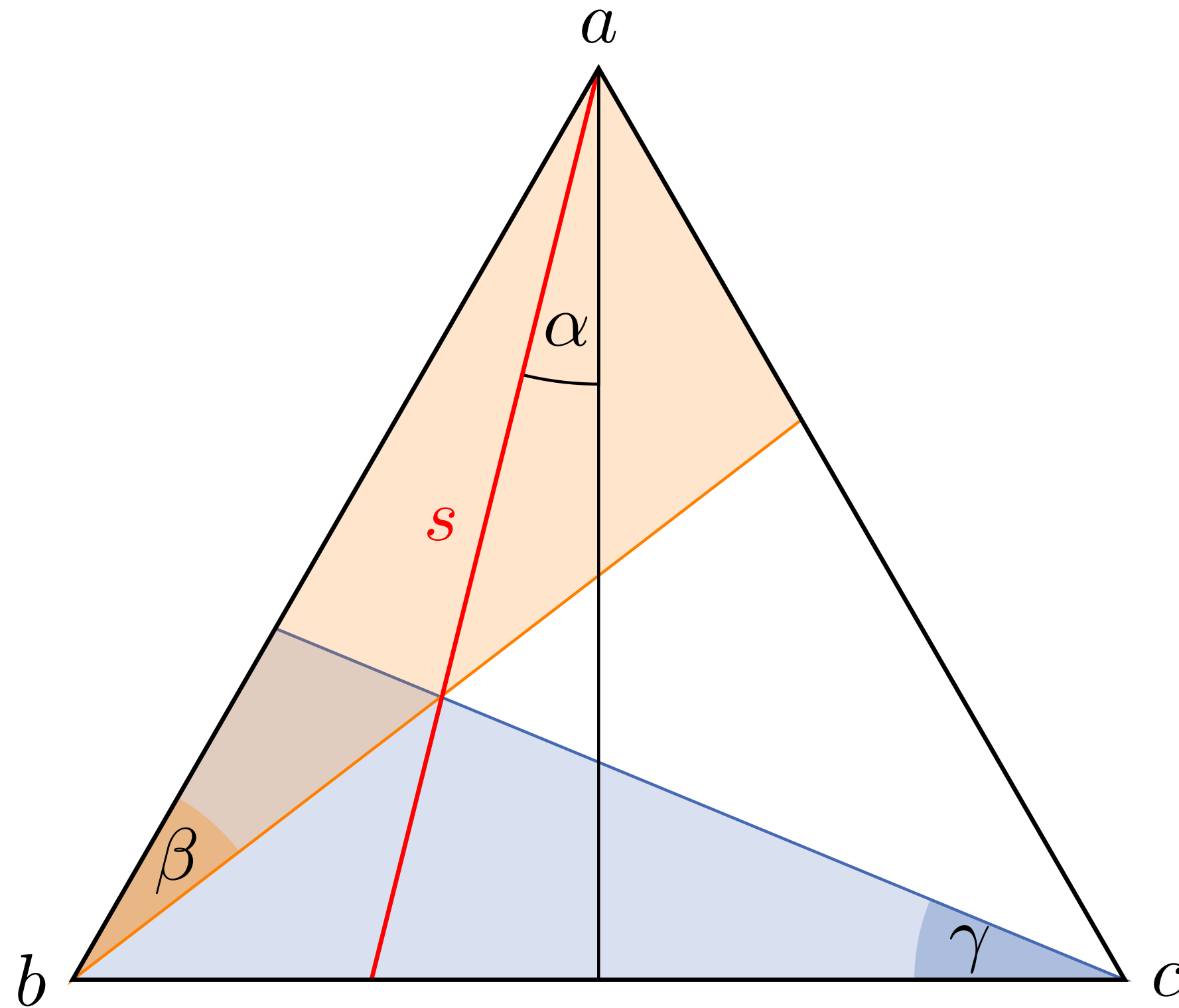
# Minimum Covering of $s$



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# Minimum Covering of $s$



# Minimum Covering of $S$

$$(\beta + \gamma)(y) = \arctan \left( \frac{\frac{1}{2} - \tan(\alpha) \left( \frac{\sqrt{3}}{2} - y \right)}{y} \right) - \frac{\pi}{6} + \arctan \left( \frac{y}{\frac{1}{2} + \tan(\alpha) \left( \frac{\sqrt{3}}{2} - y \right)} \right)$$



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$$y_{\min} = -\frac{1 - 3 \tan^2(\alpha)}{2\sqrt{3} (\tan^2(\alpha) + 1)} + \sqrt{\left( \frac{1 - 3 \tan^2(\alpha)}{2\sqrt{3} \tan^2(\alpha) + 2} \right)^2 - \frac{3 \tan^2(\alpha) - 1}{4 (\tan^2(\alpha) + 1)}}$$

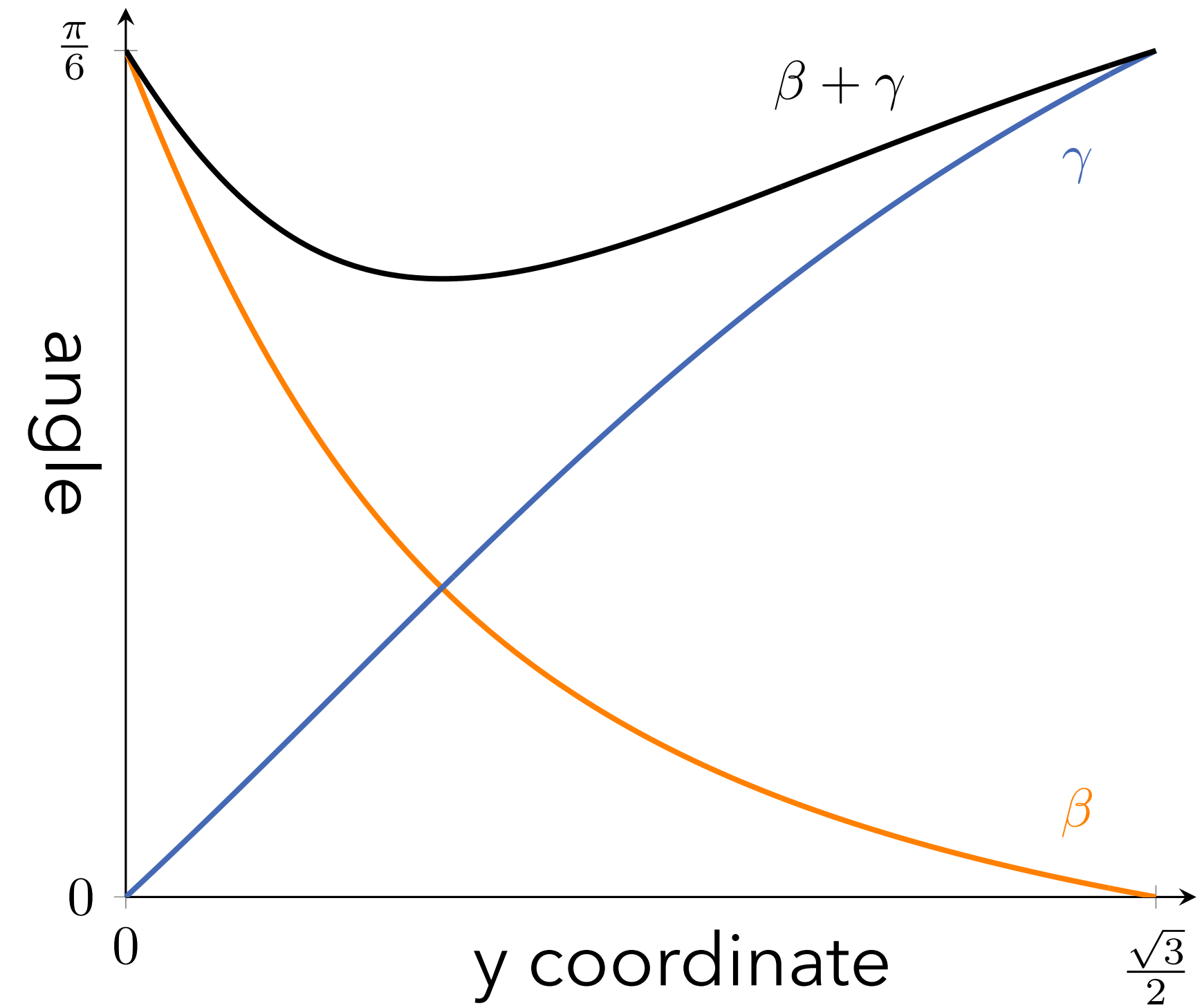
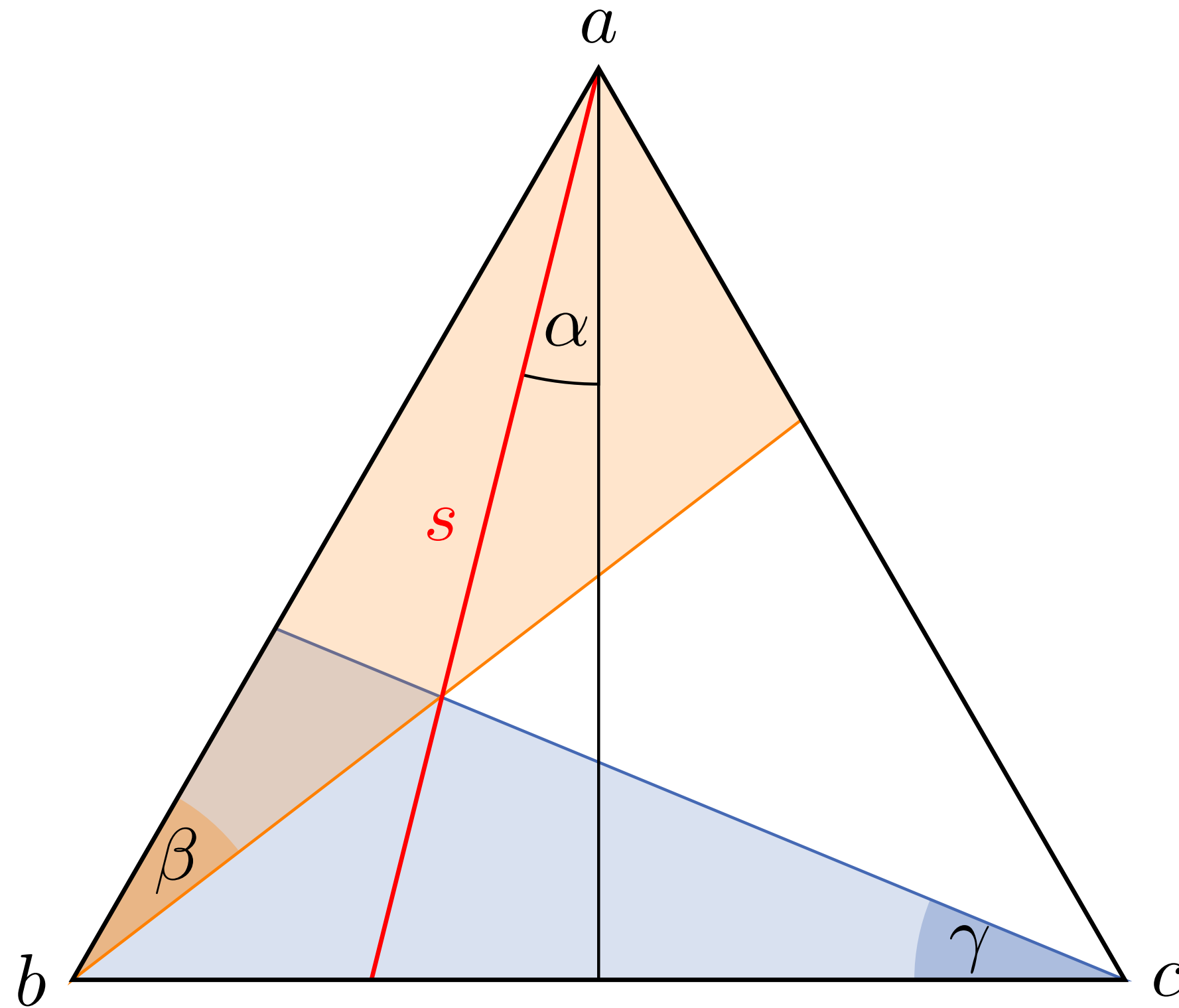
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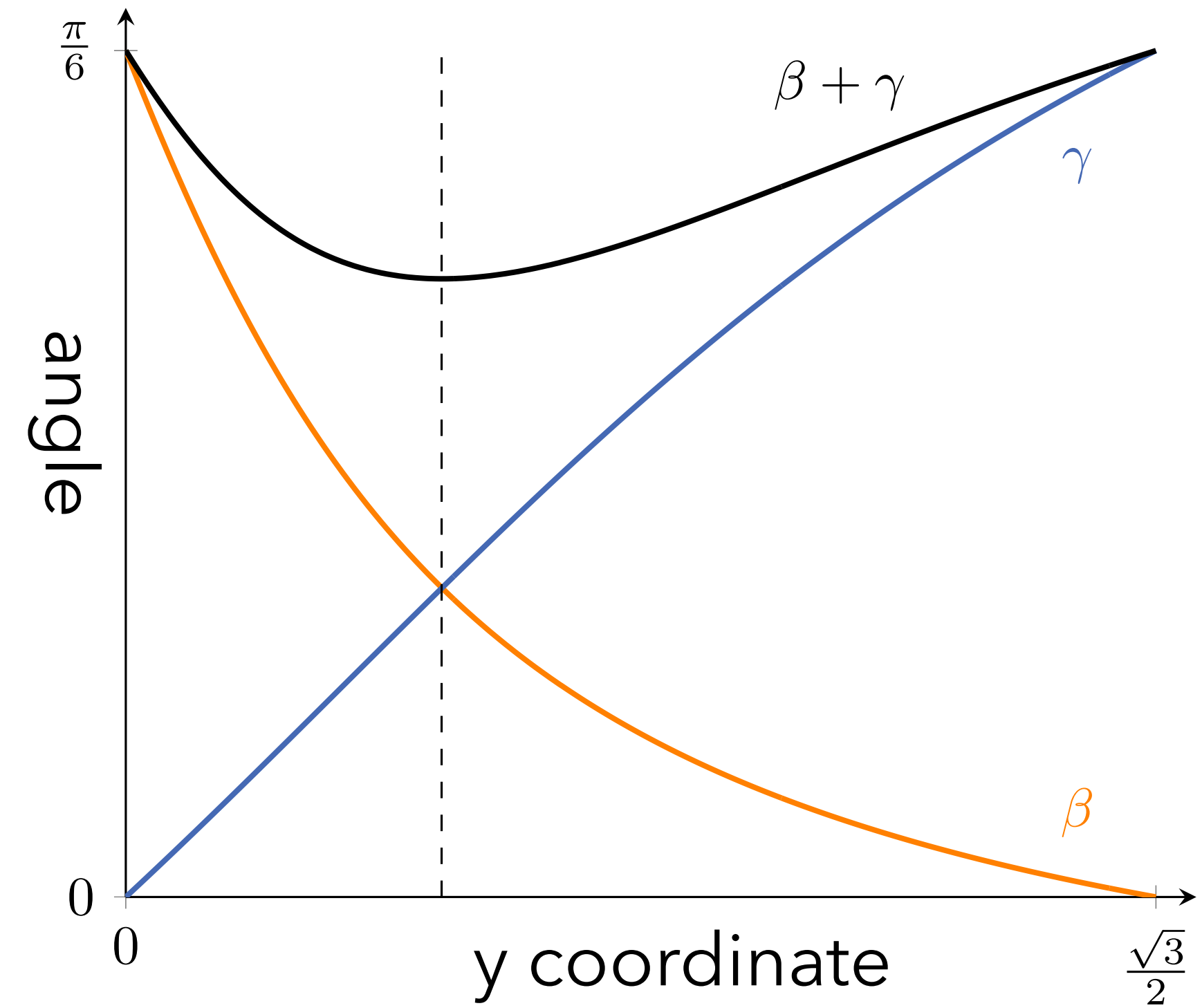
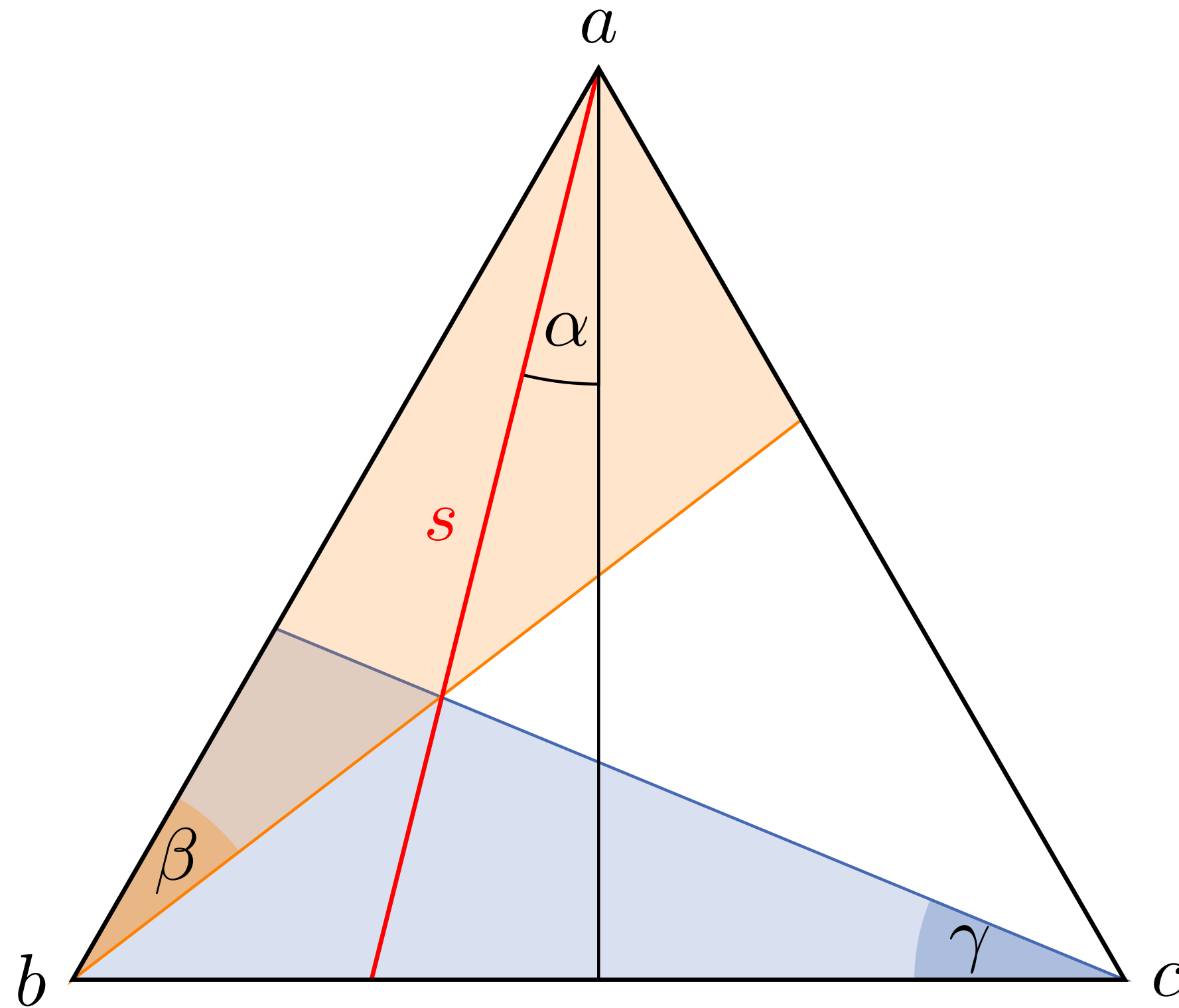
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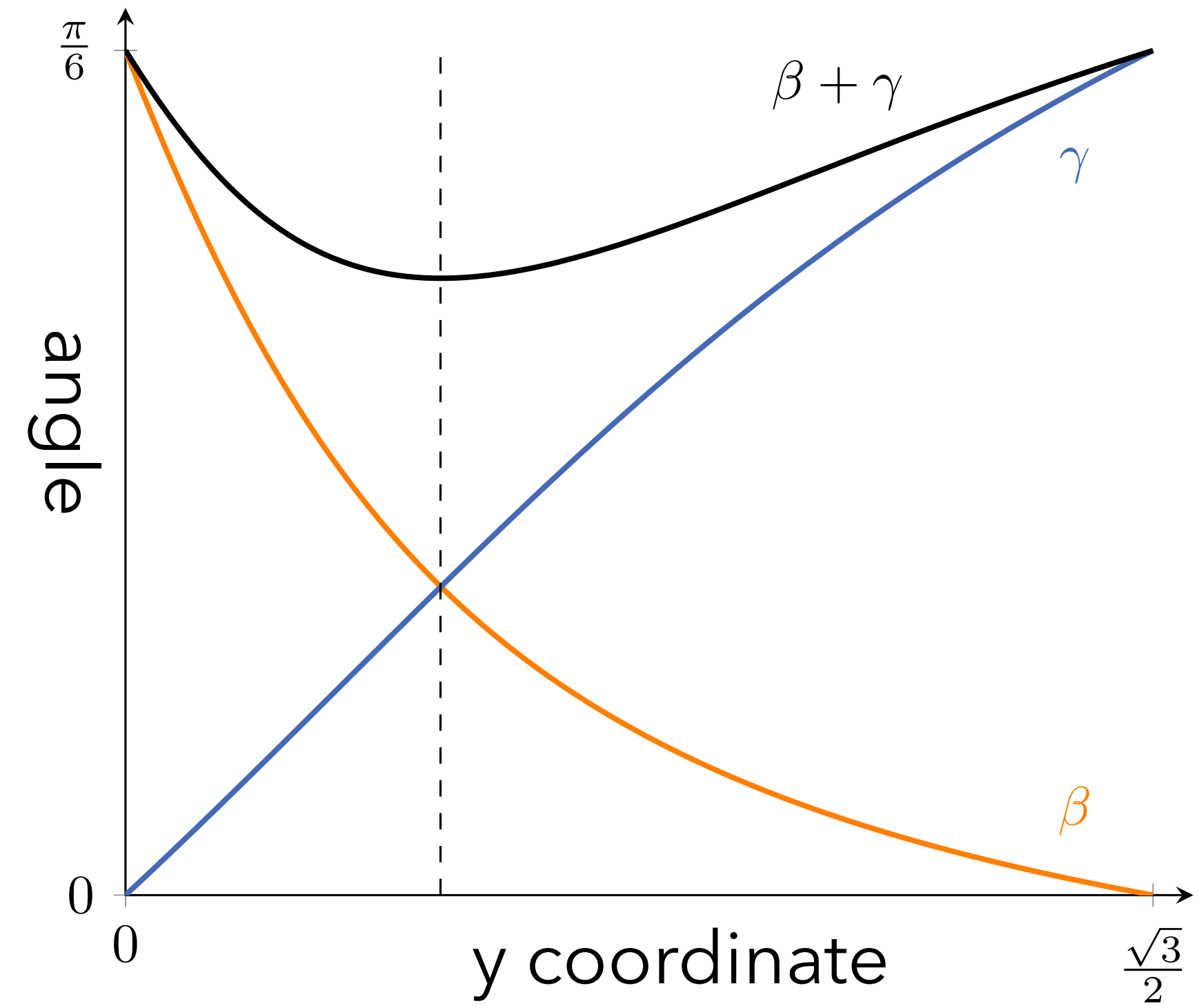
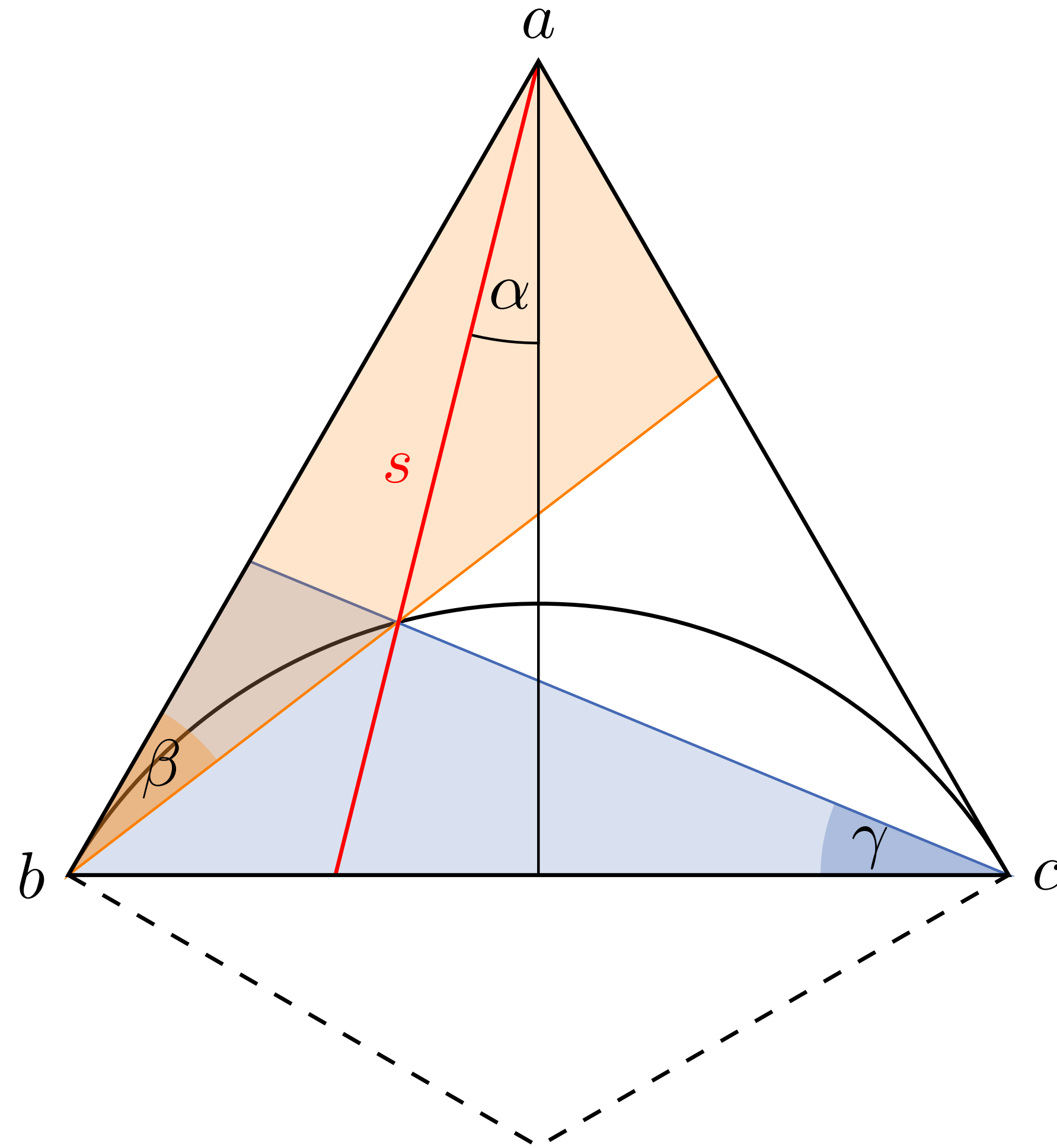
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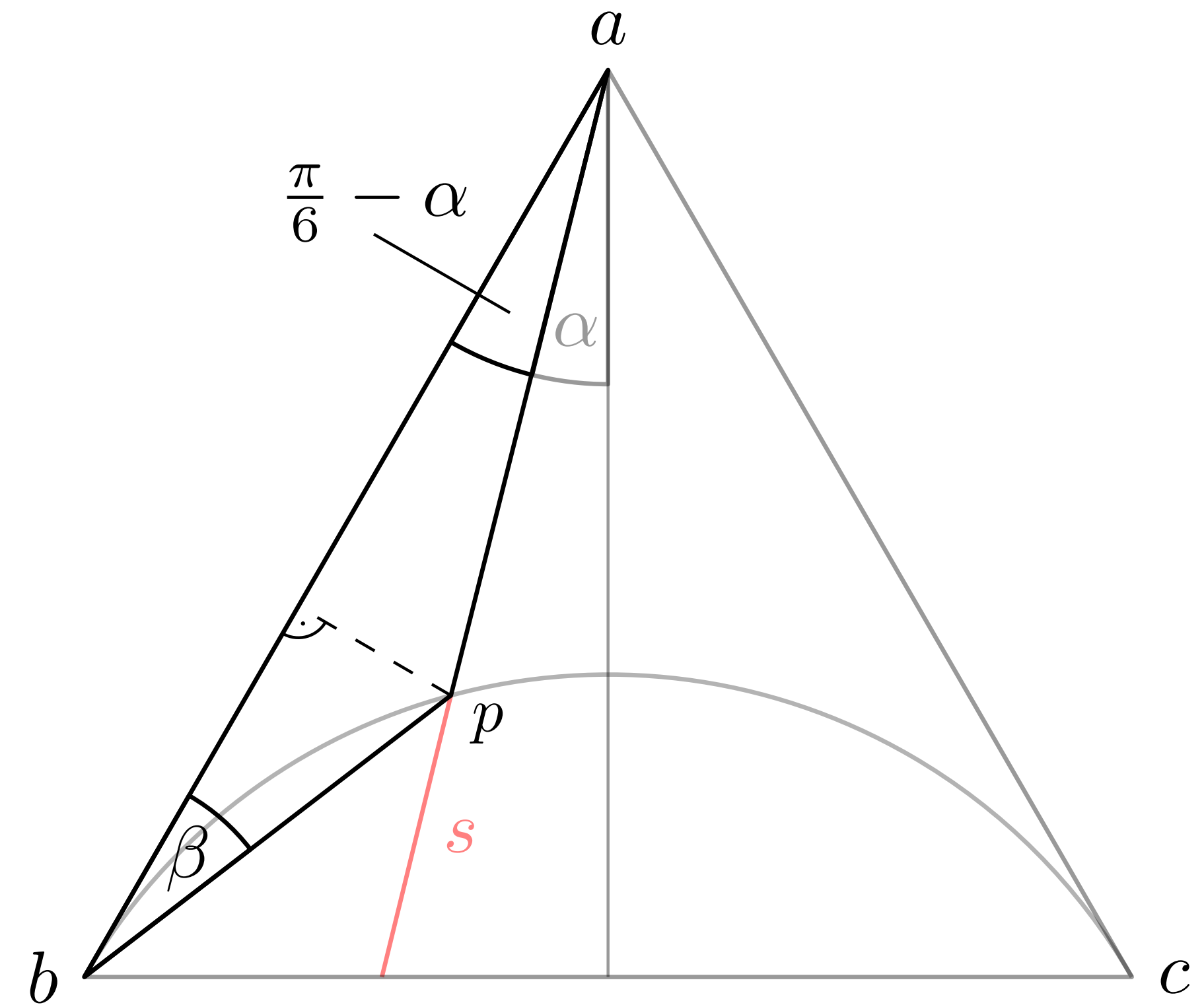
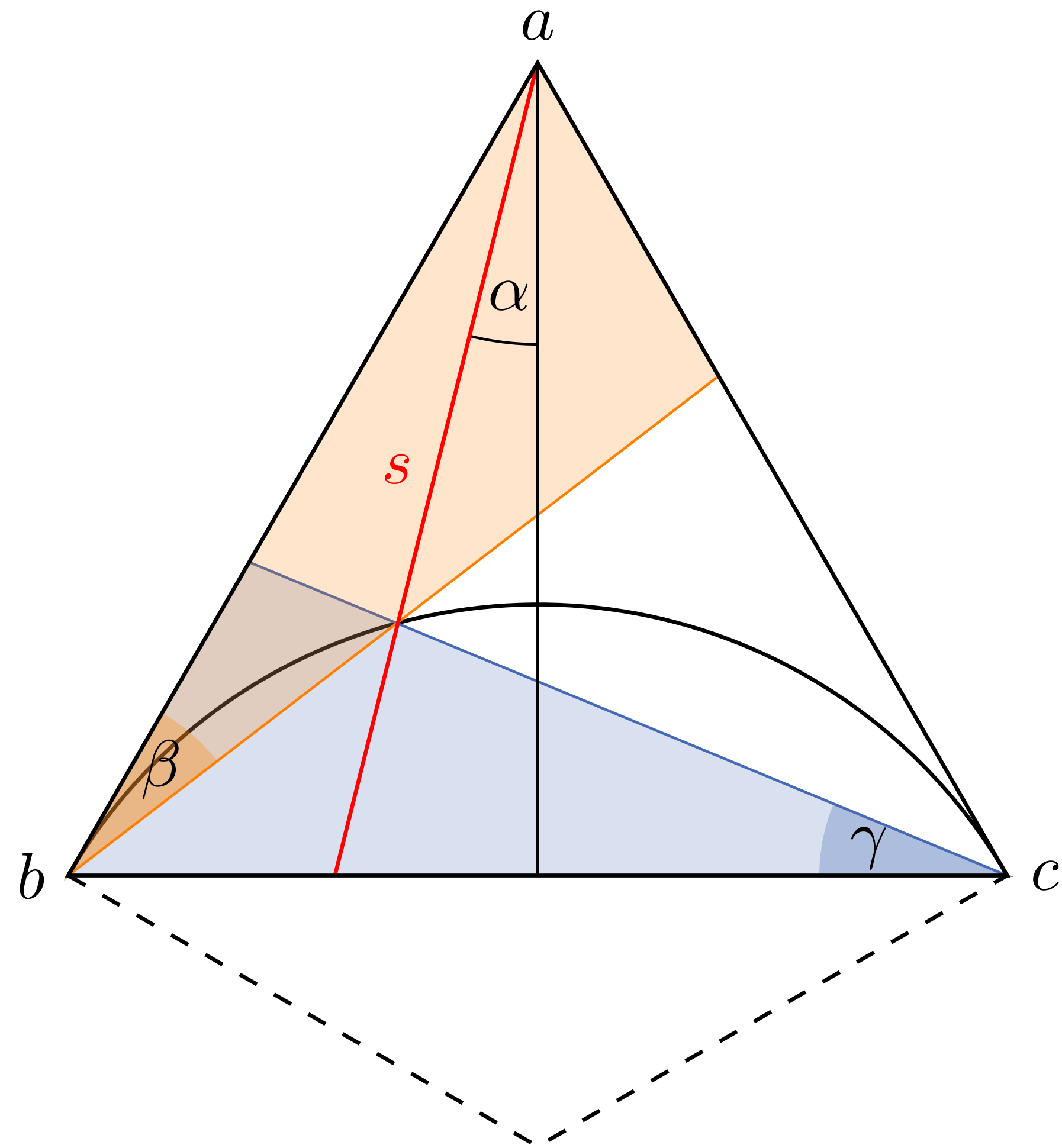
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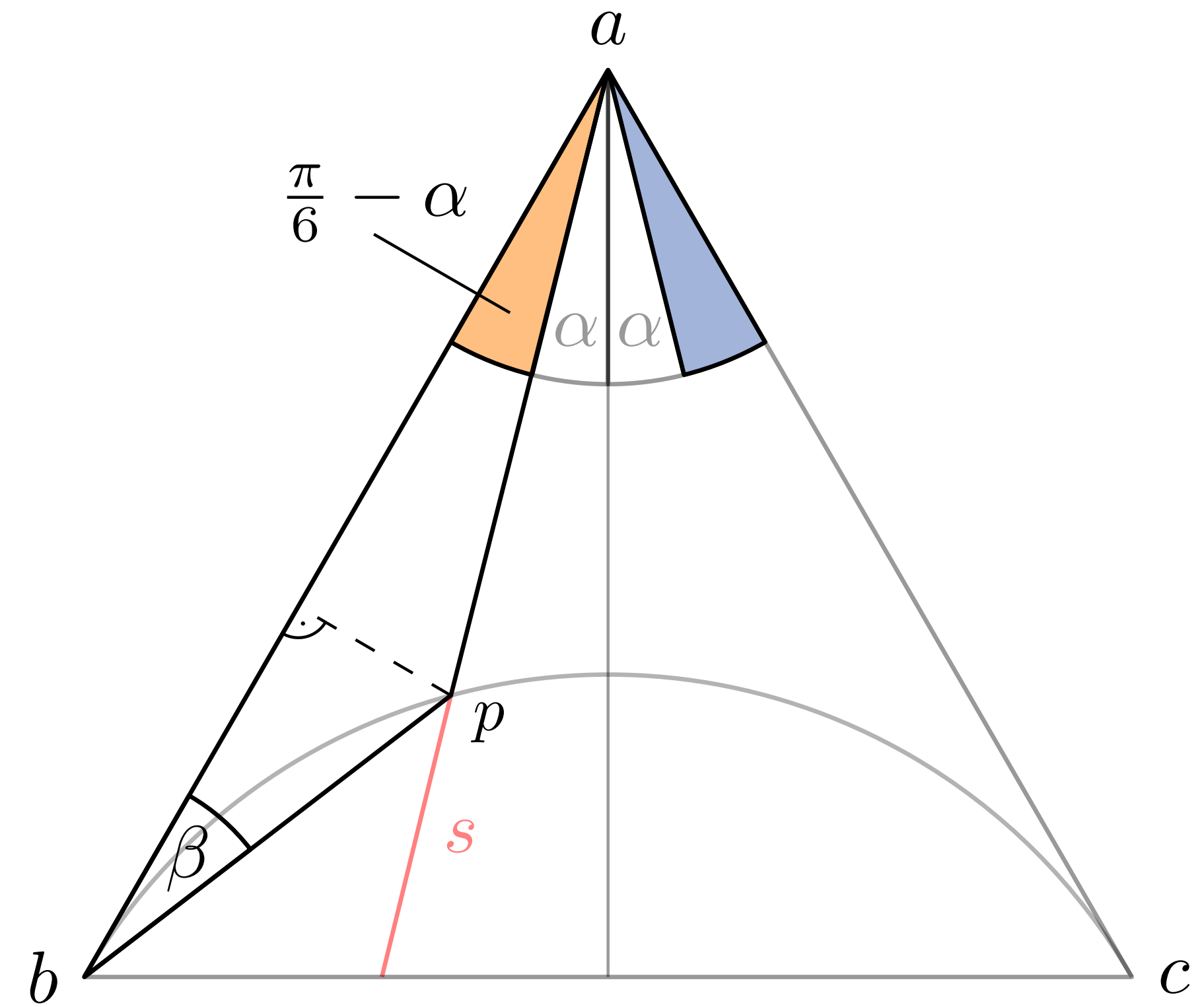
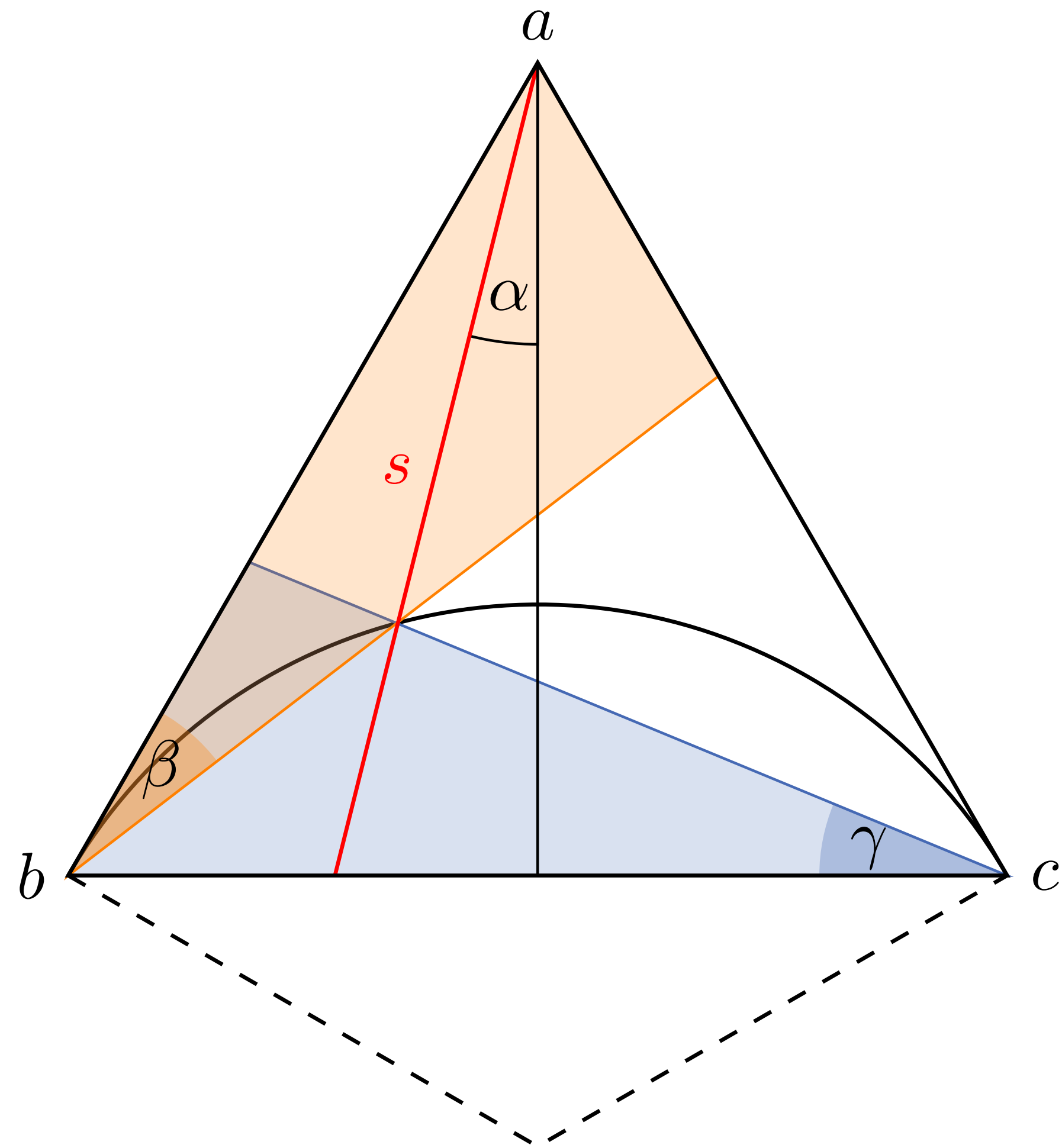
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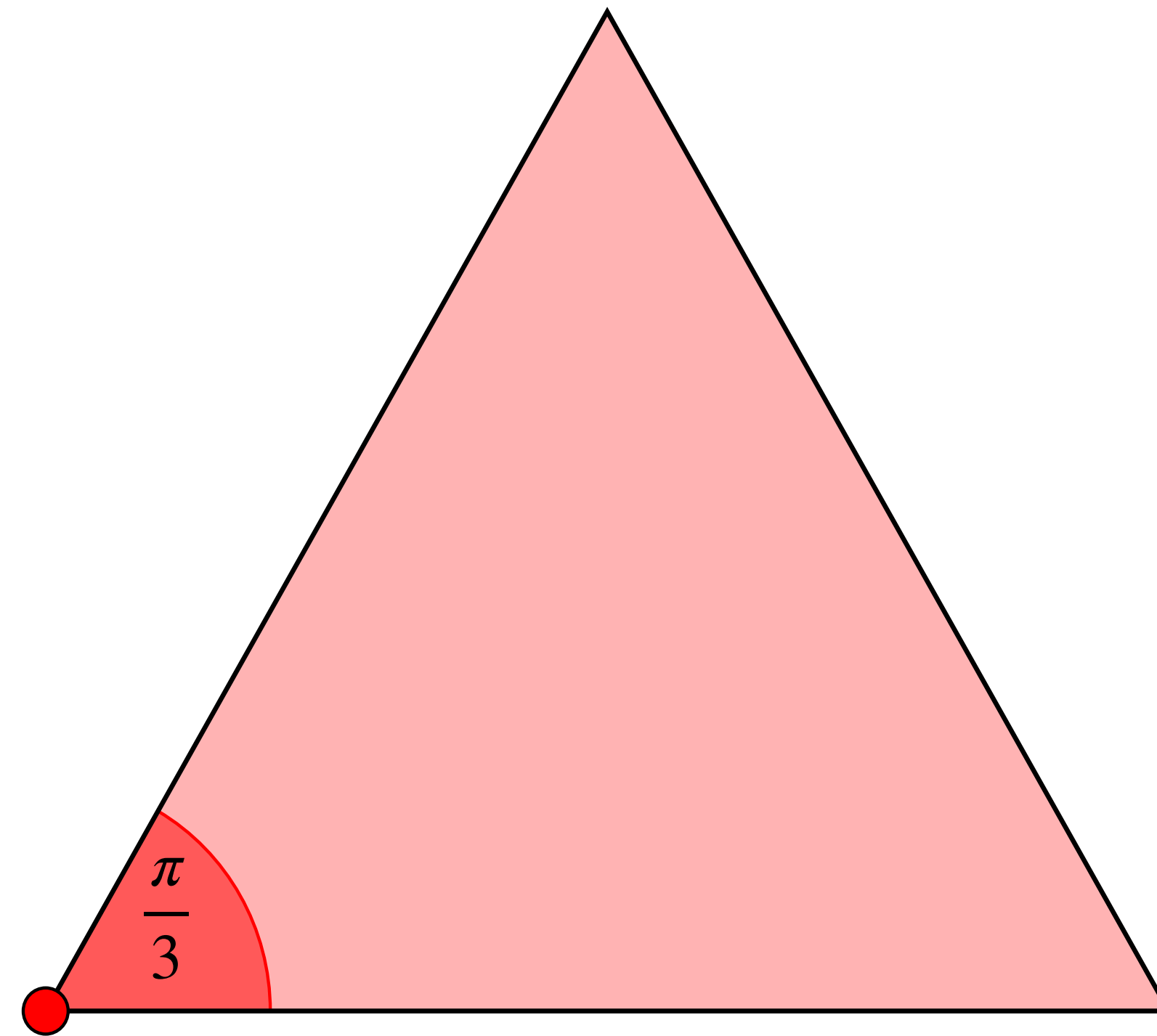
# Minimum Covering of $s$



# Minimum Covering of $s$



# Optimal Covering of Equilateral Triangles







Introduction

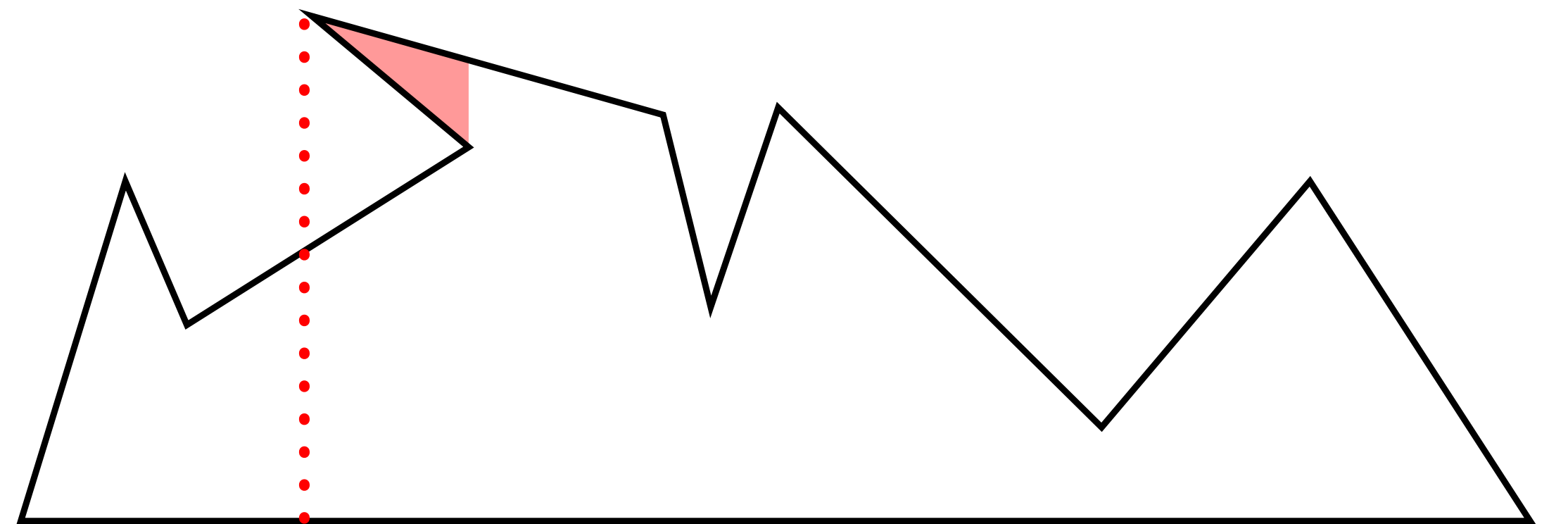
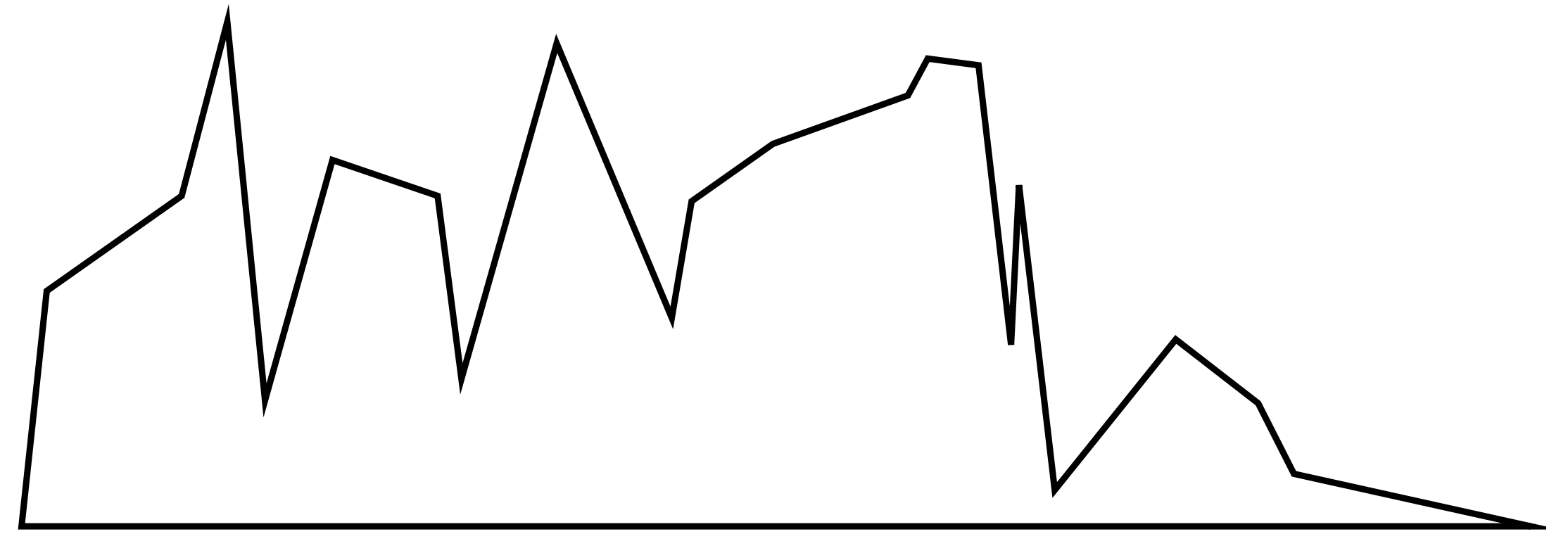
Equilateral Triangles

## Tight Upper Bound for **Histograms**

Simple Polygons

Duality

Conclusion



Introduction

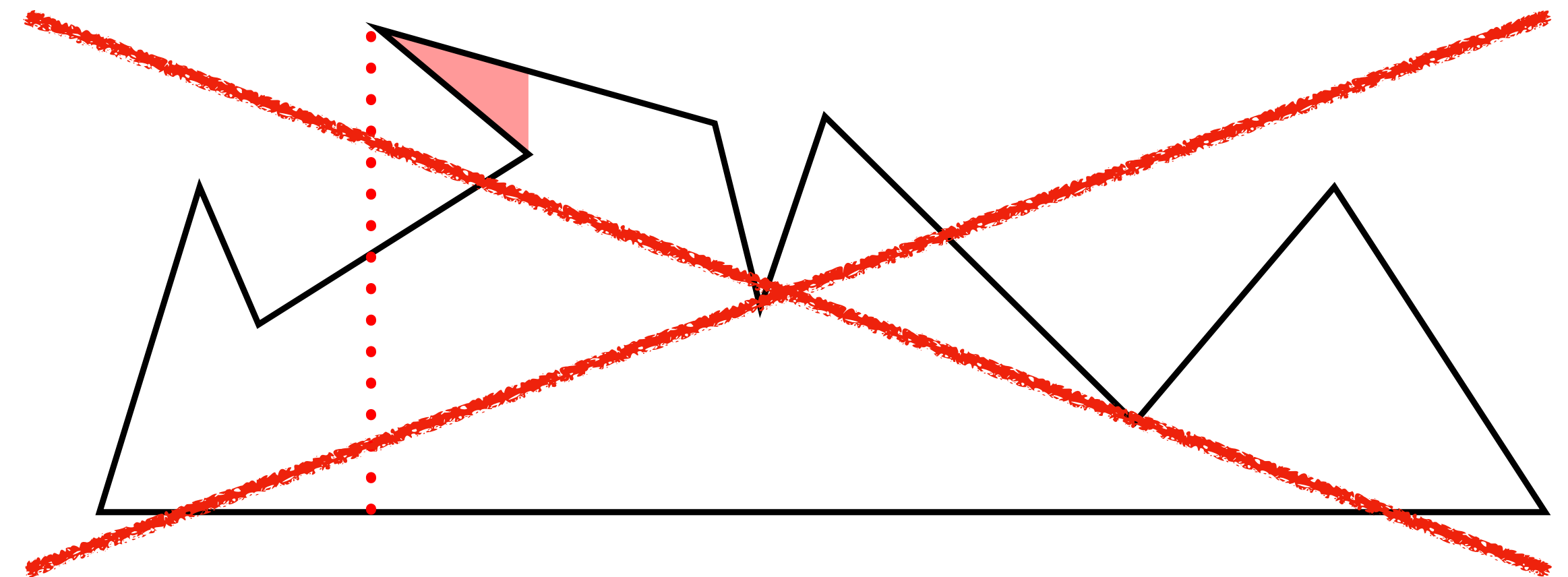
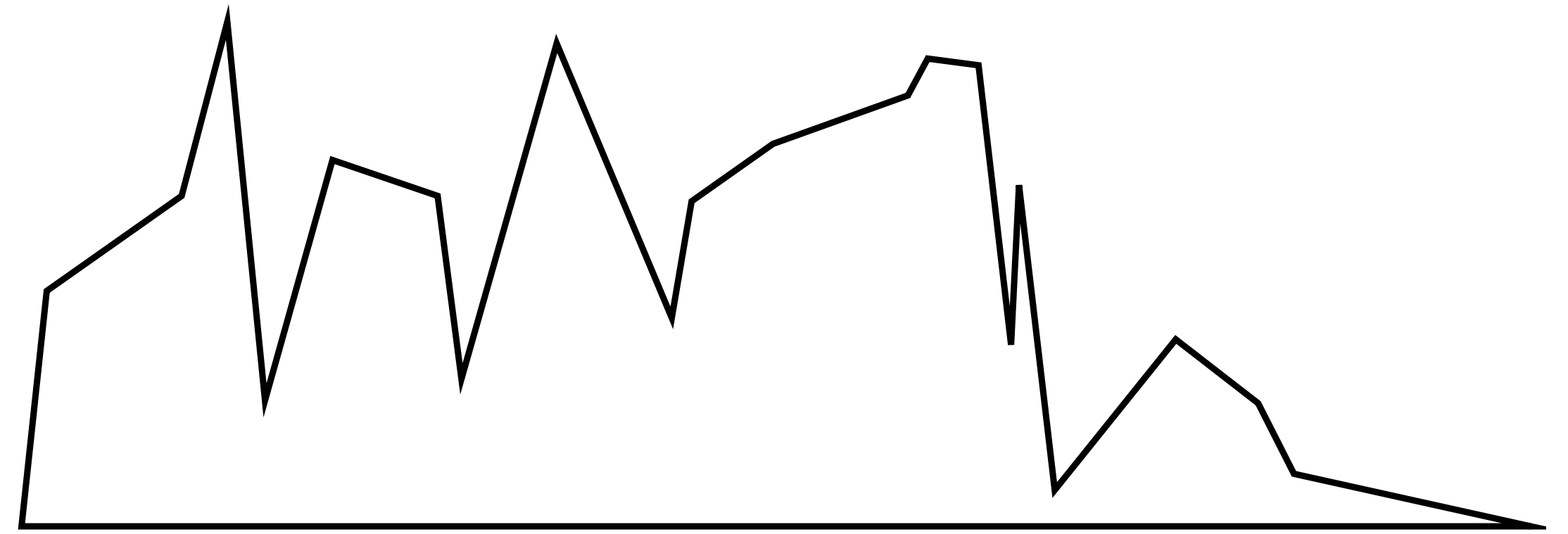
Equilateral Triangles

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Introduction

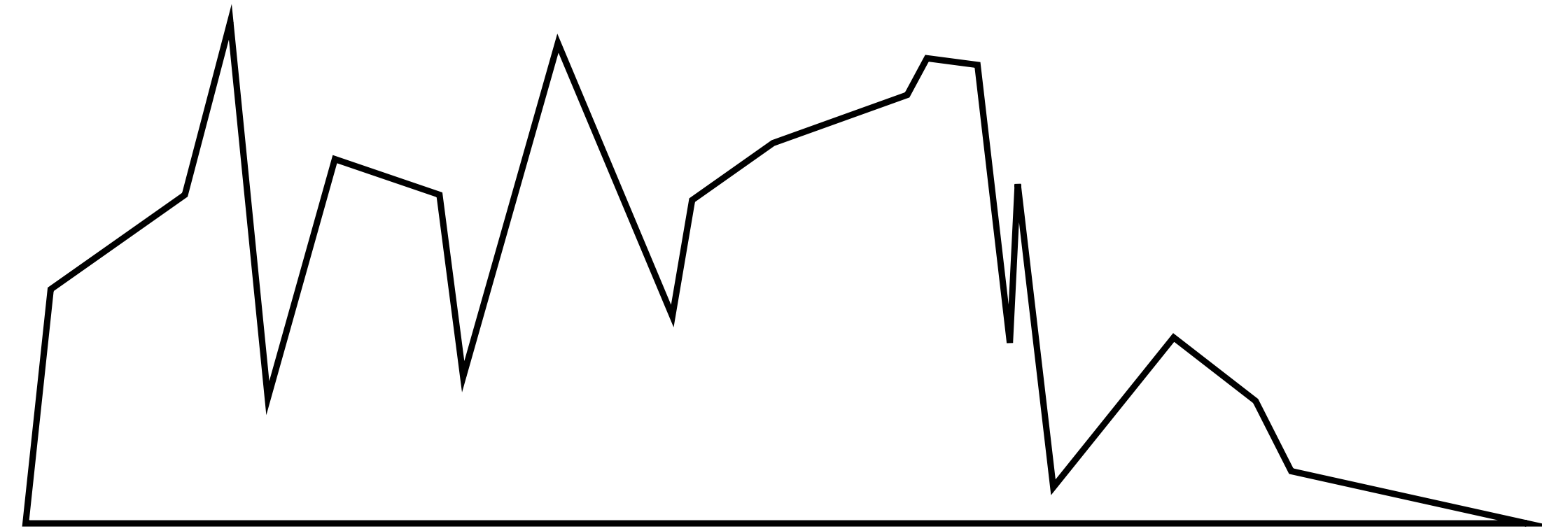
Equilateral Triangles

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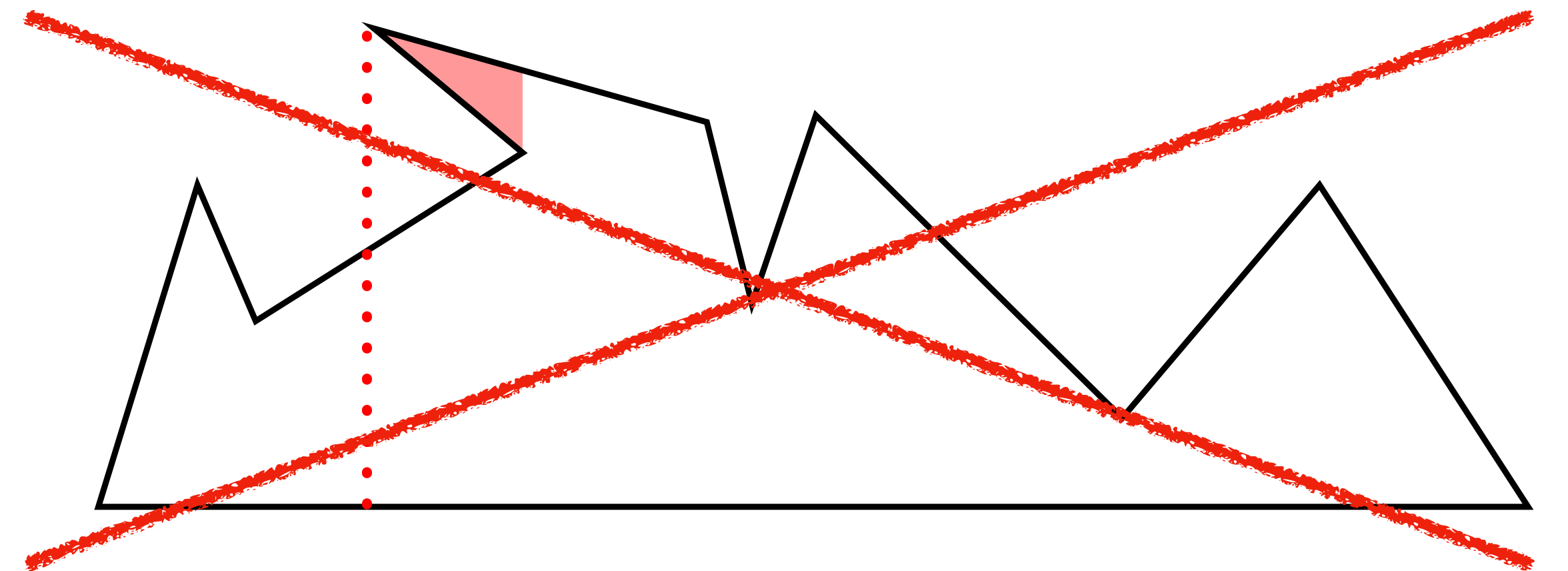
Simple Polygons

Duality

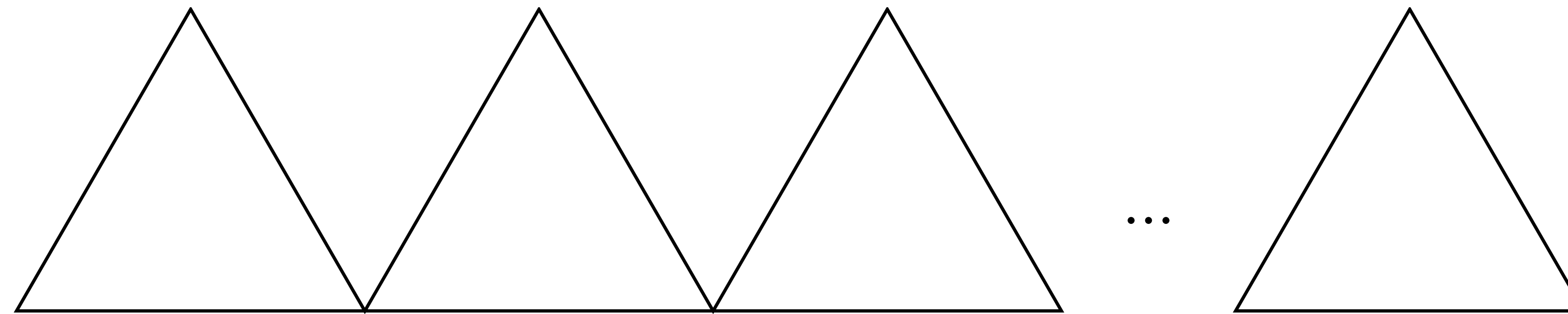
Conclusion



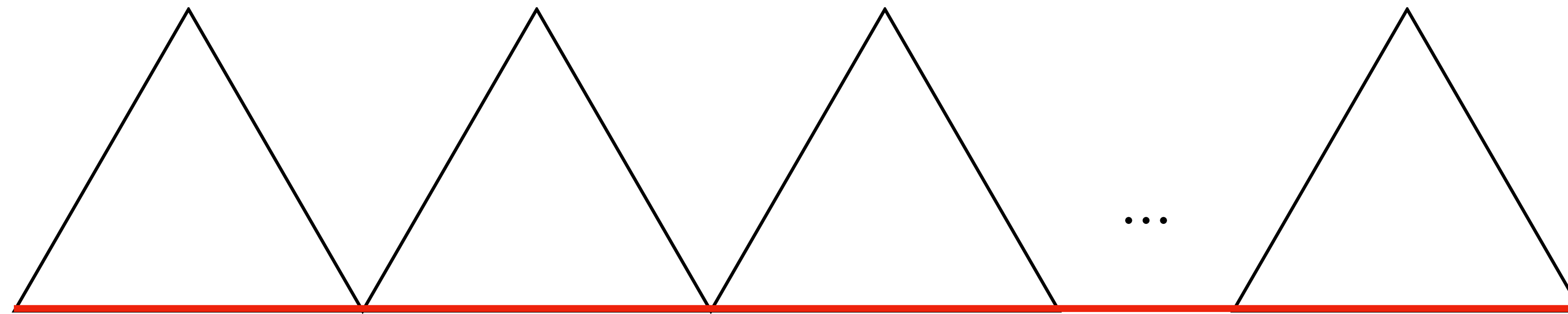
$$(n - 1) \frac{\pi}{6}$$



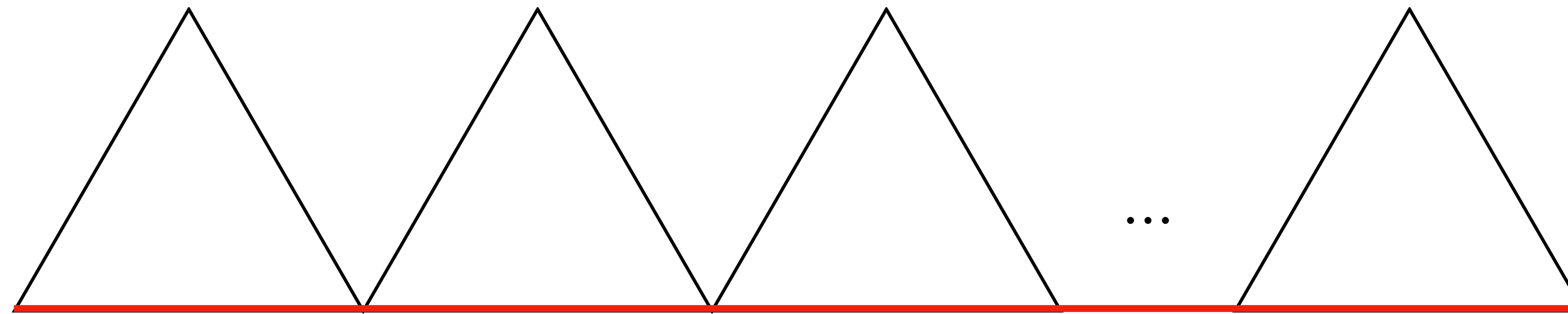
# Lower Bound for Histograms



# Lower Bound for Histograms

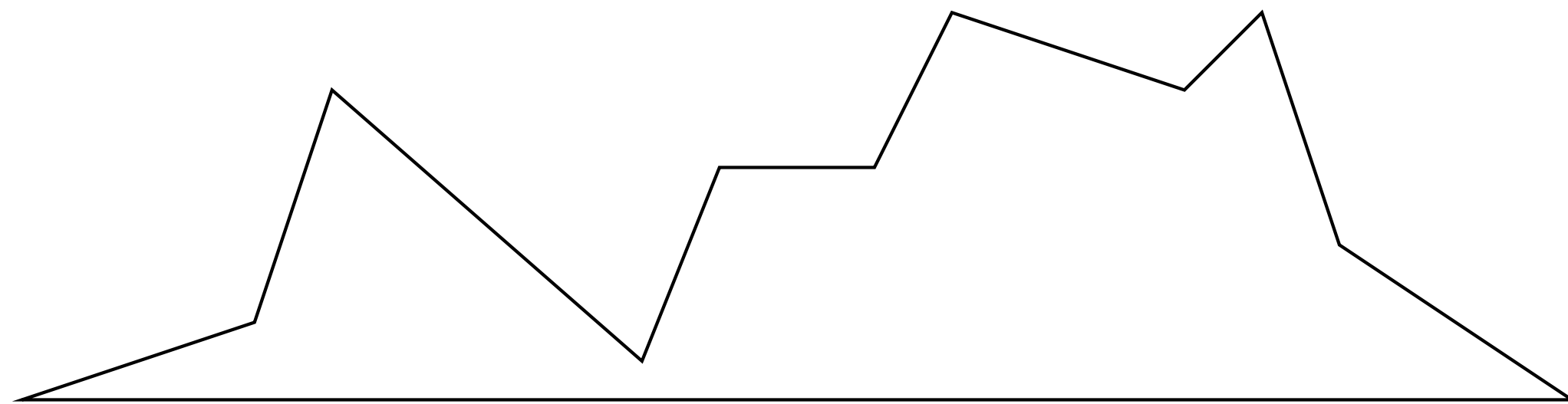


# Lower Bound for Histograms



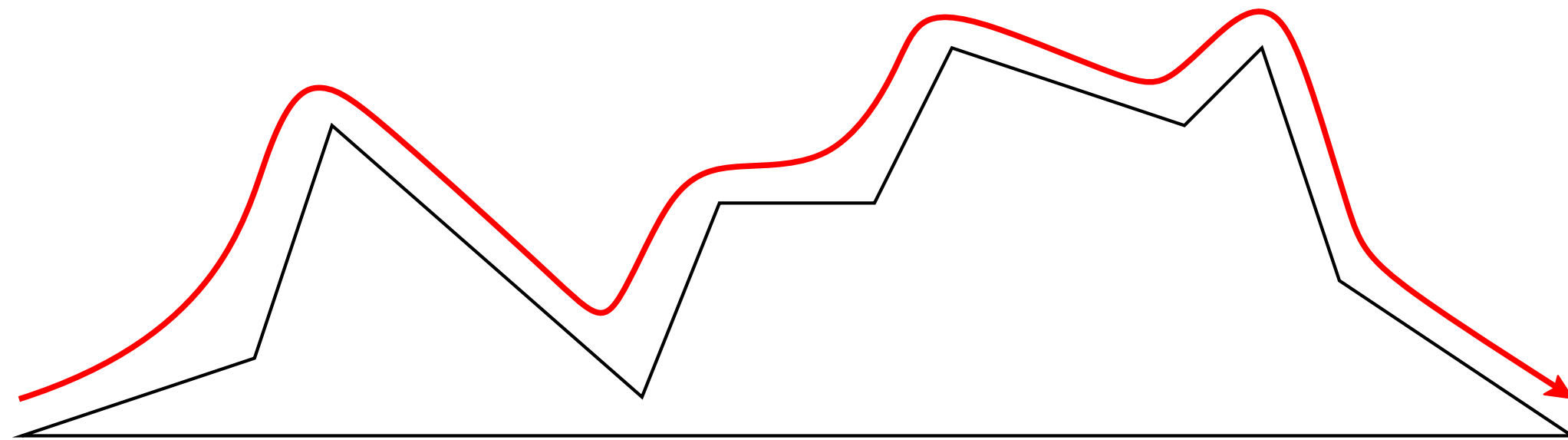
$$\text{Lower Bound: } \frac{n-1}{2} \frac{\pi}{3} = (n-1) \frac{\pi}{6}$$

# Upper Bound for Histograms

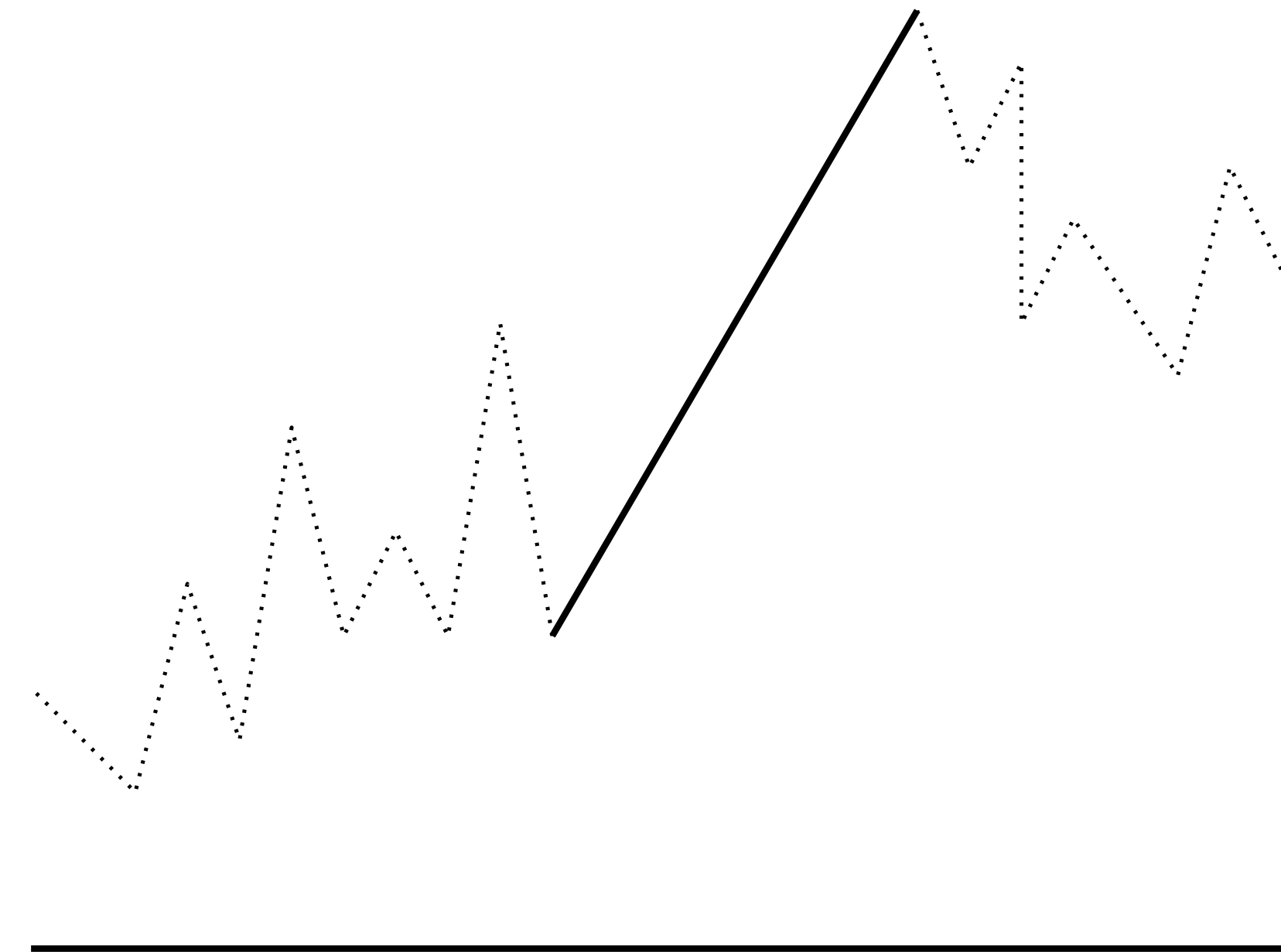
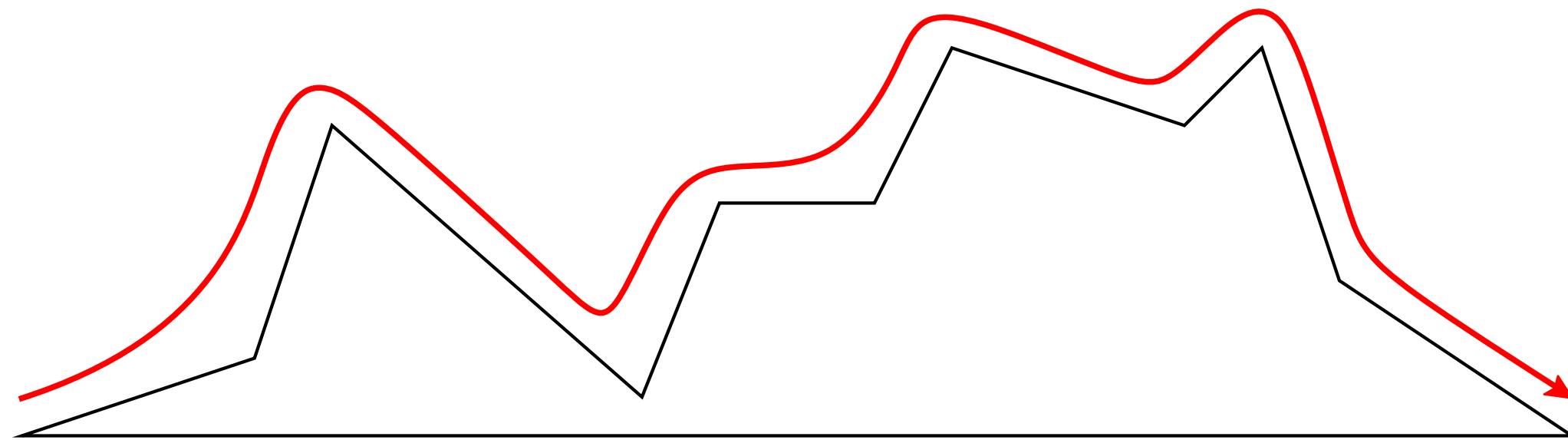




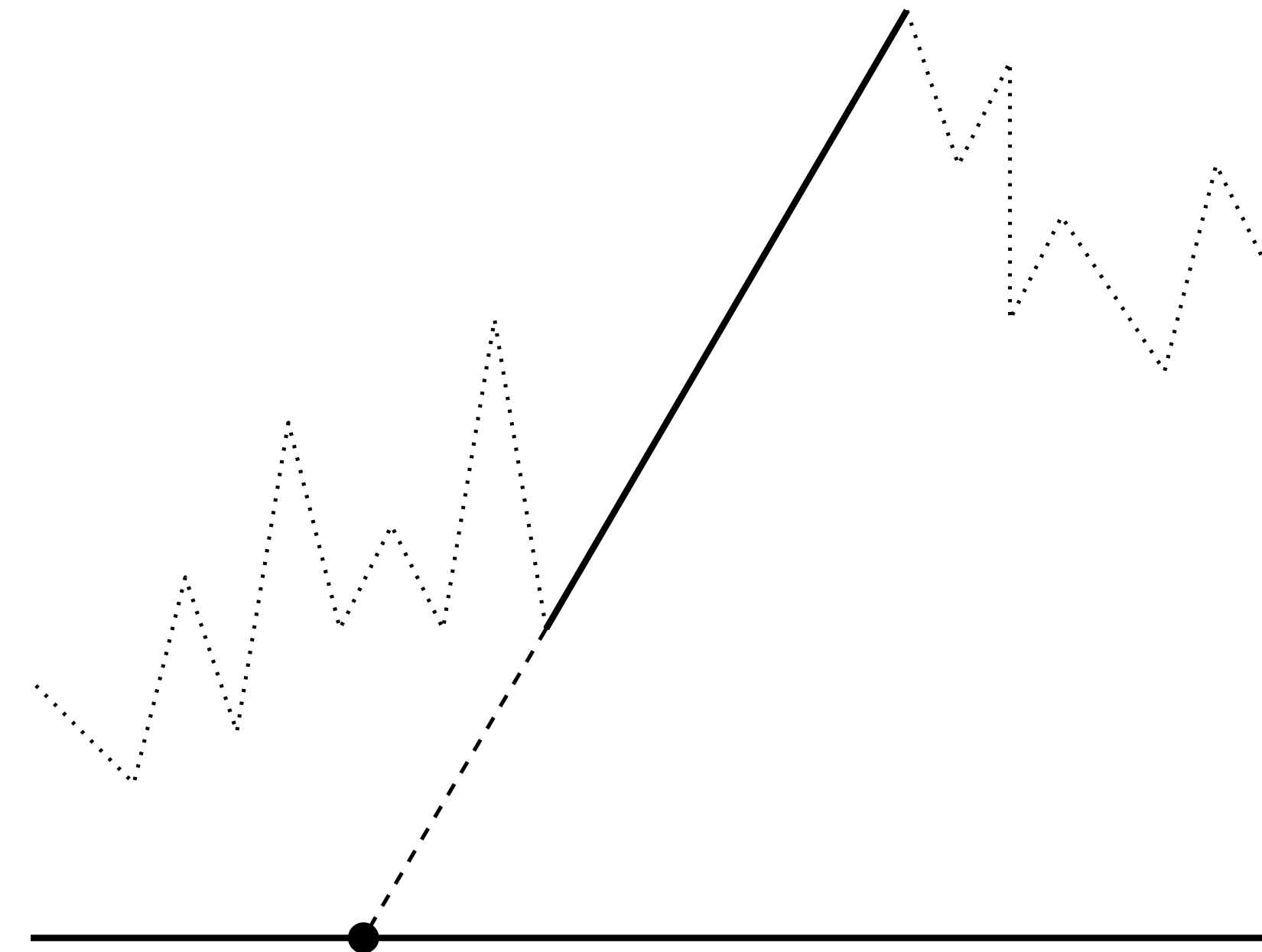
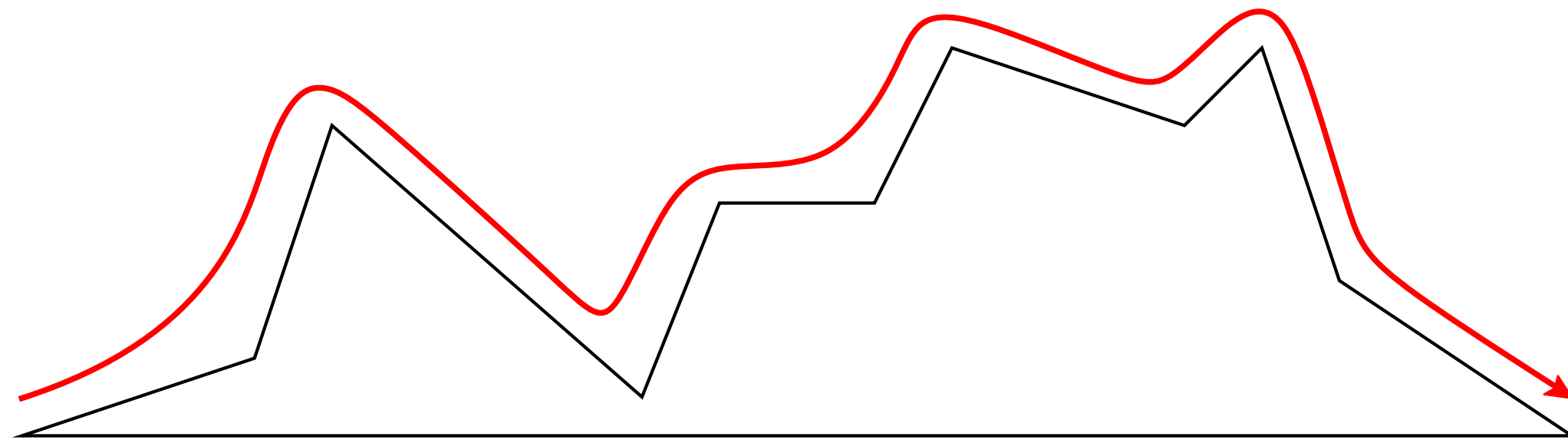
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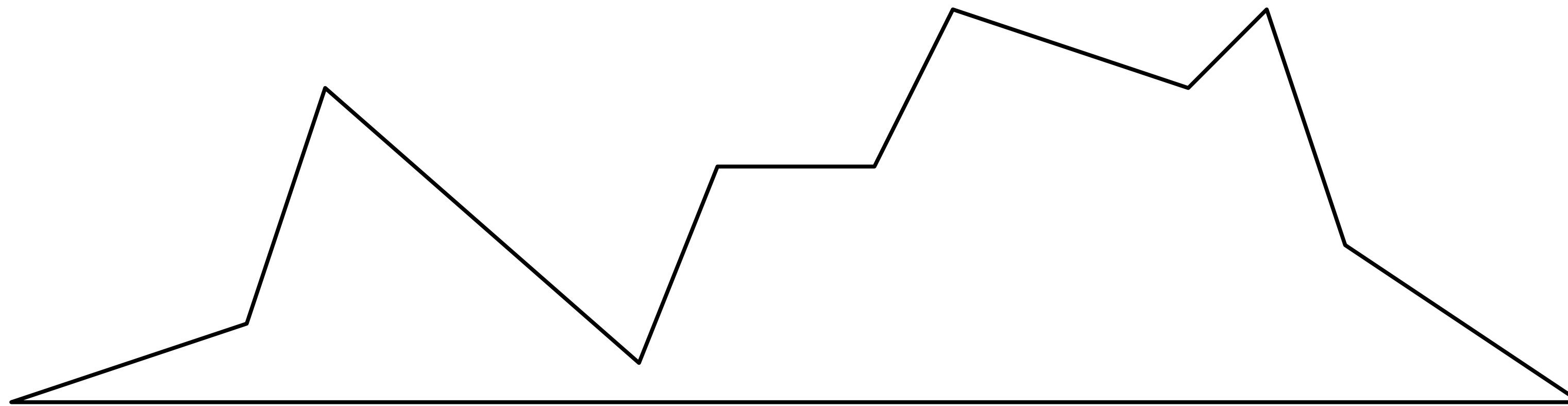
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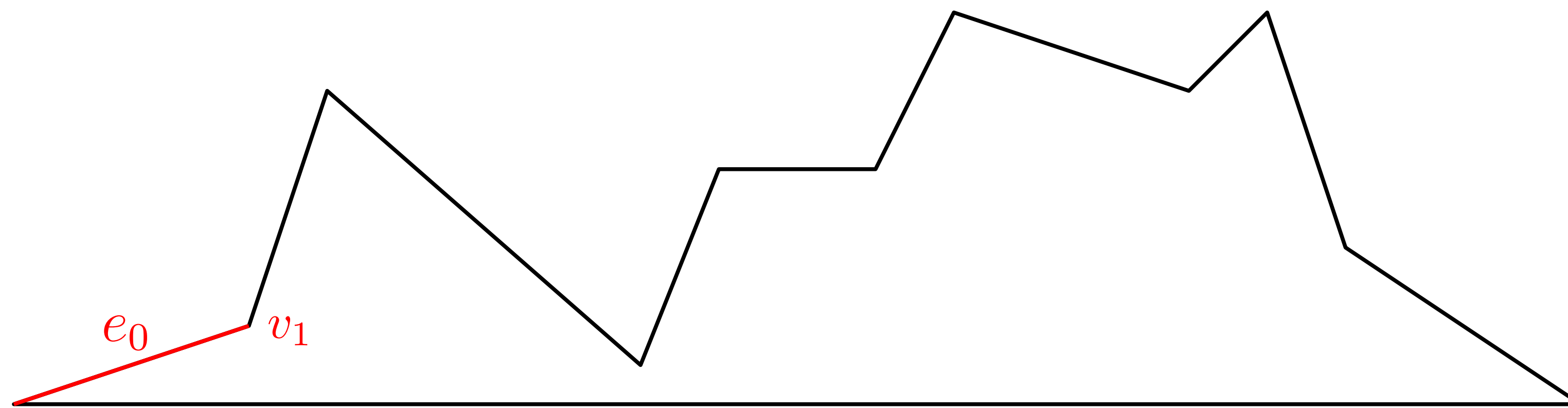


# Histogram Covering Algorithm – Example



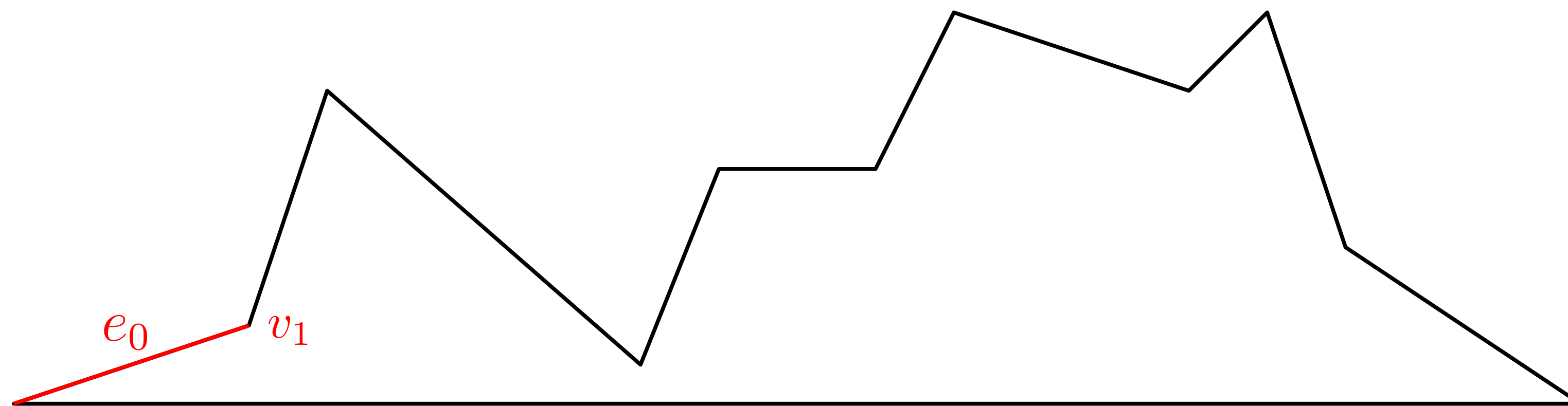
Case	Condition
1	$e_i$ is covered
2	$v_{i+1}$ is reflex
3	$v_{i+2}$ is convex
4	there is another uncovered, completely visible edge
5	visibility extension of $e_{i+1}$ intersects with baseline
6	none of the above applies

# Histogram Covering Algorithm – Example



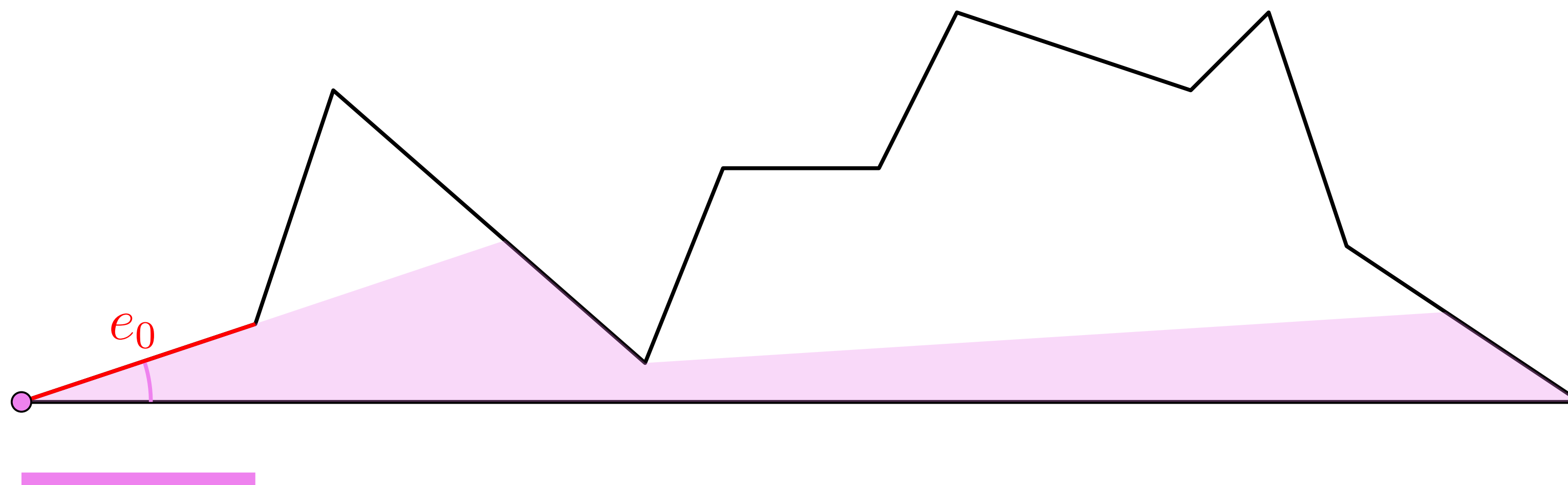
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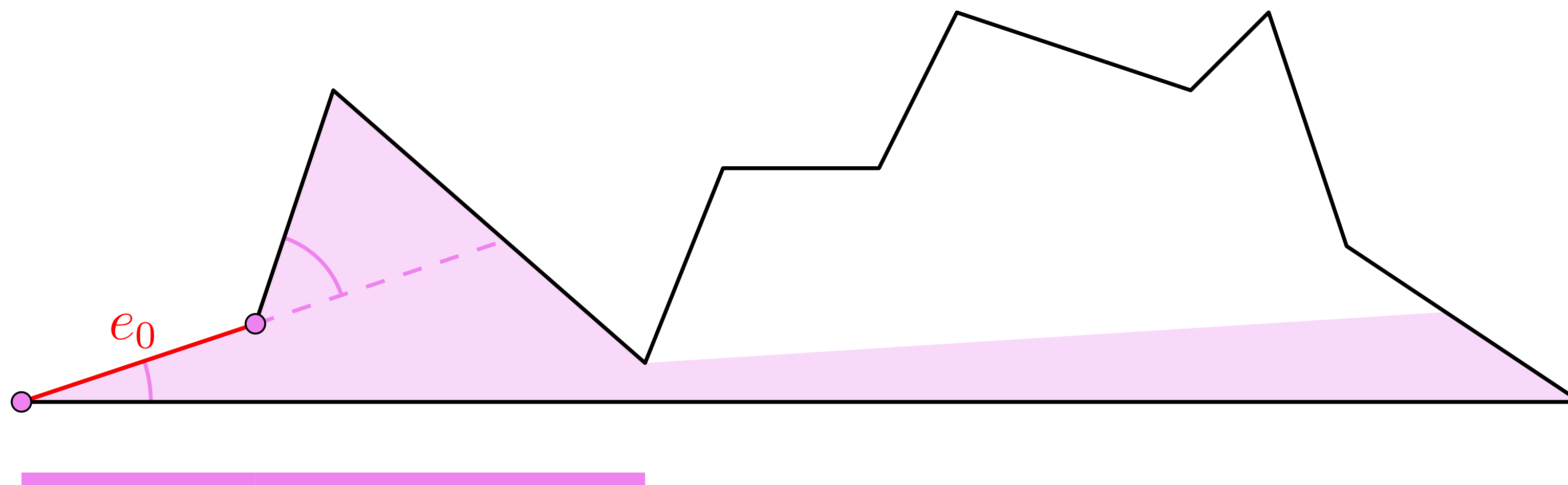
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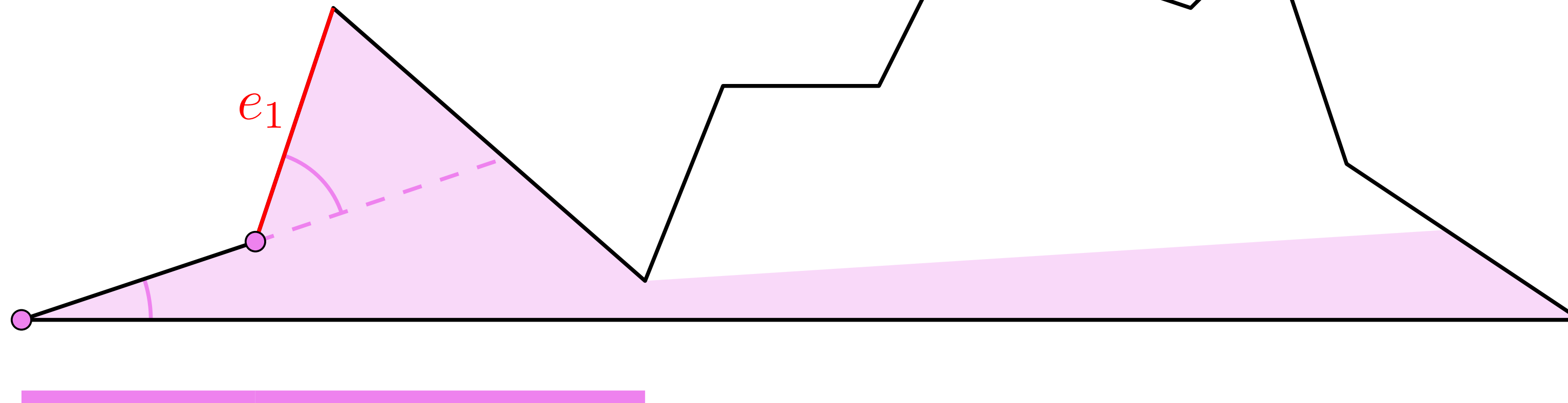
# Histogram Covering Algorithm – Example



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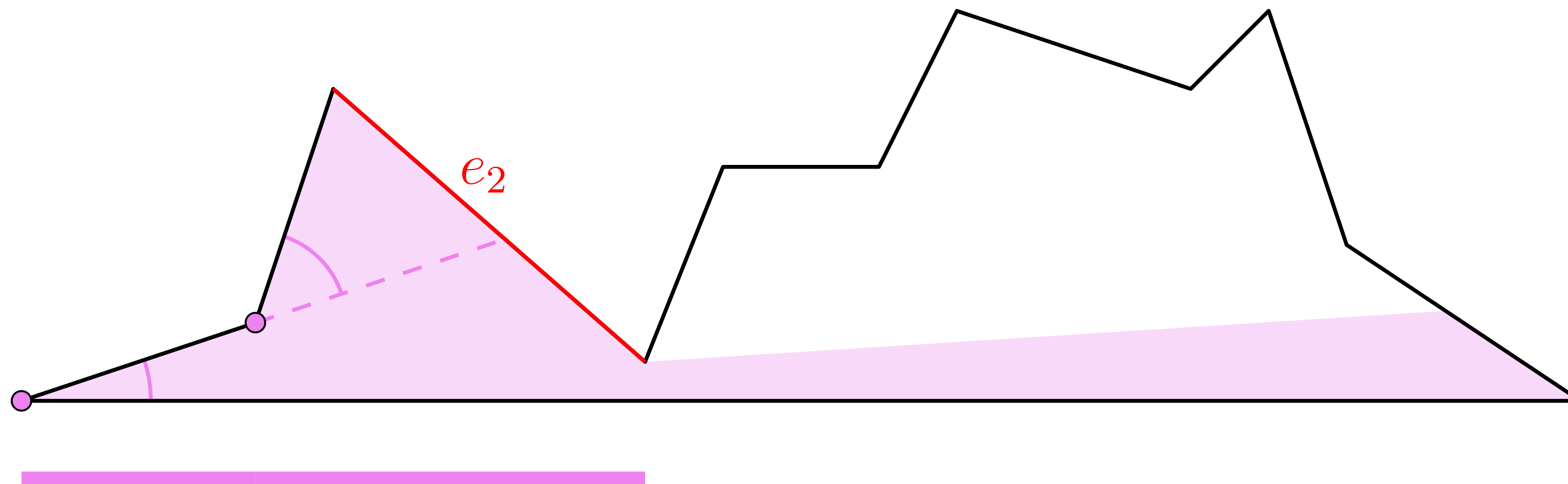


# Histogram Covering Algorithm – Example



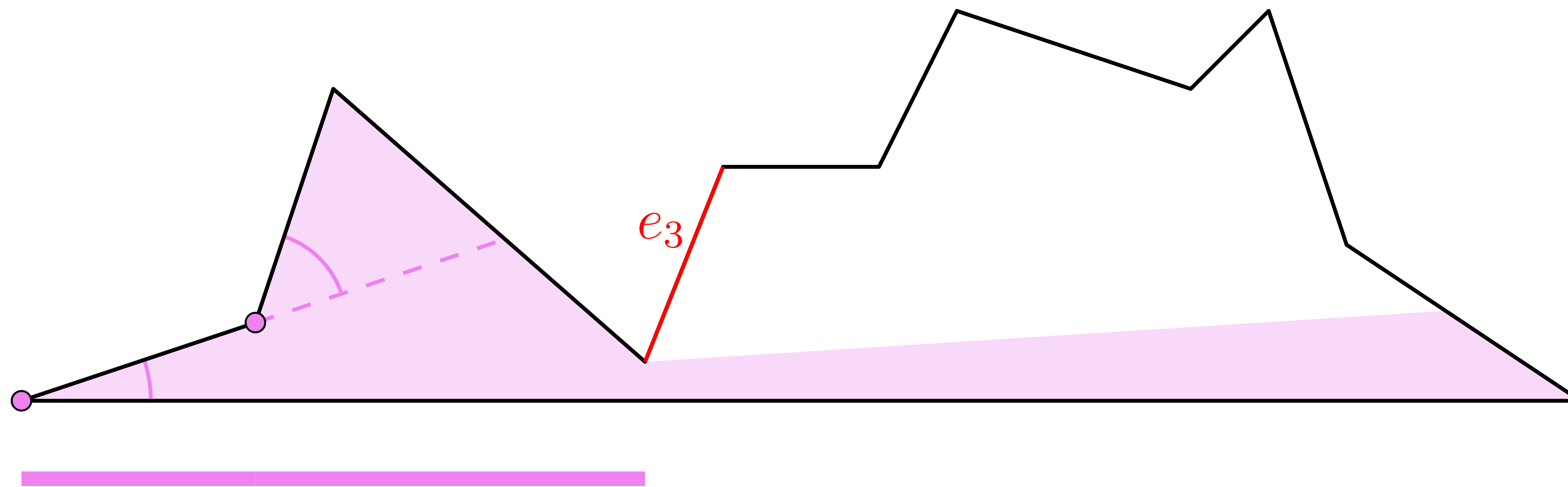
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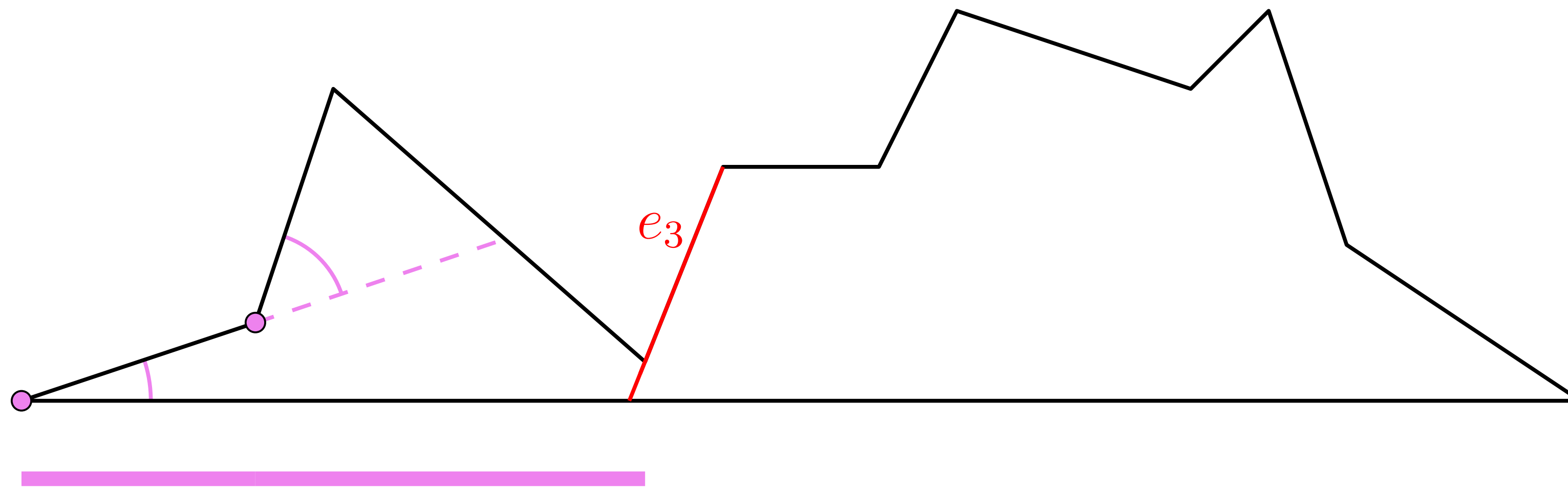
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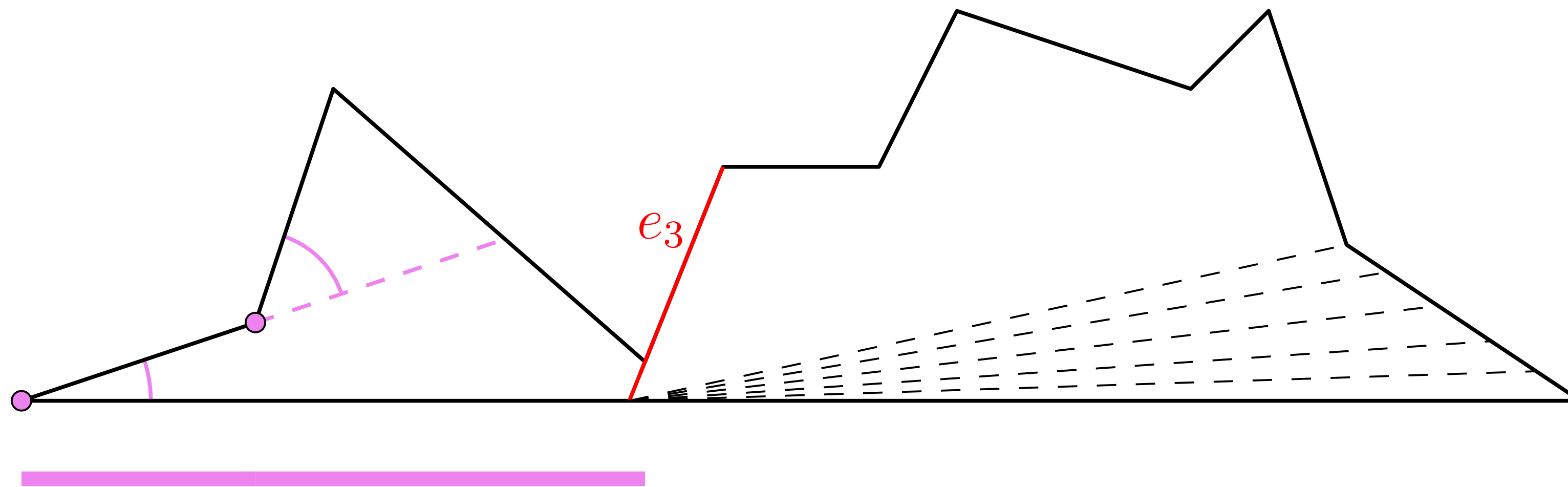
Case	Condition
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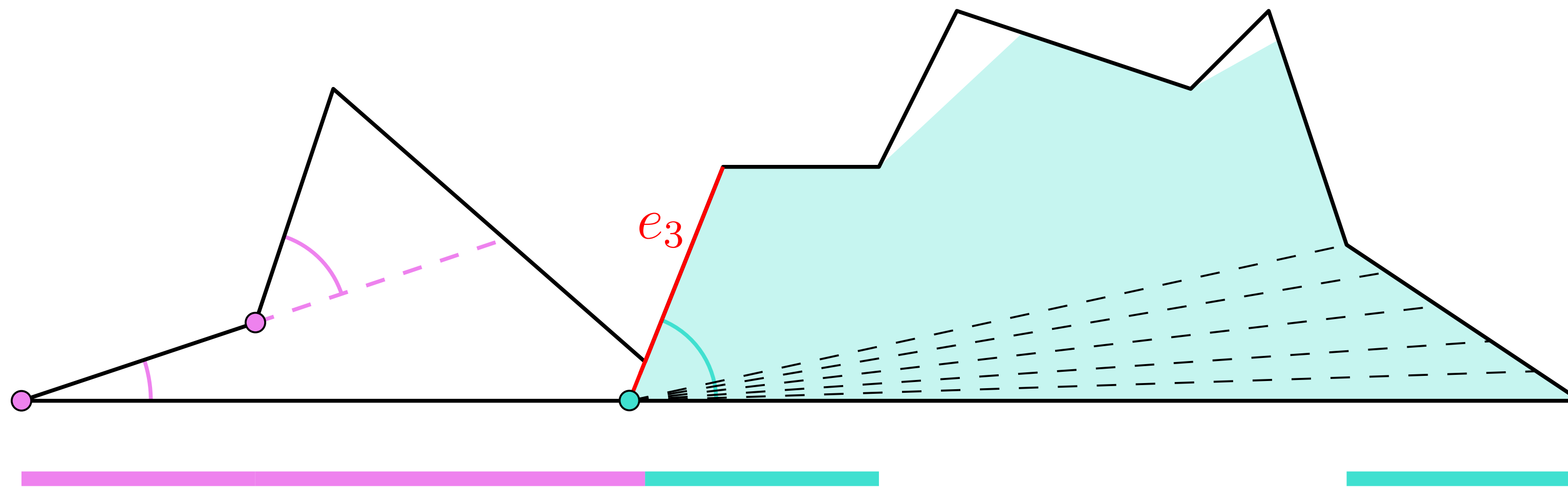
Case	Condition
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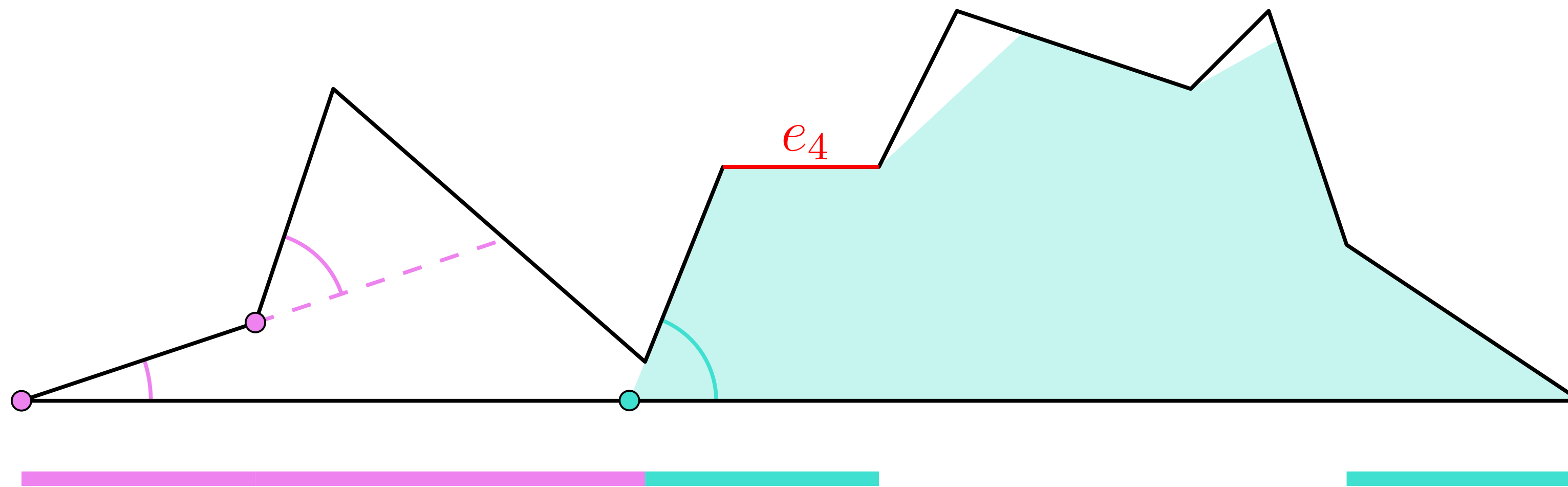
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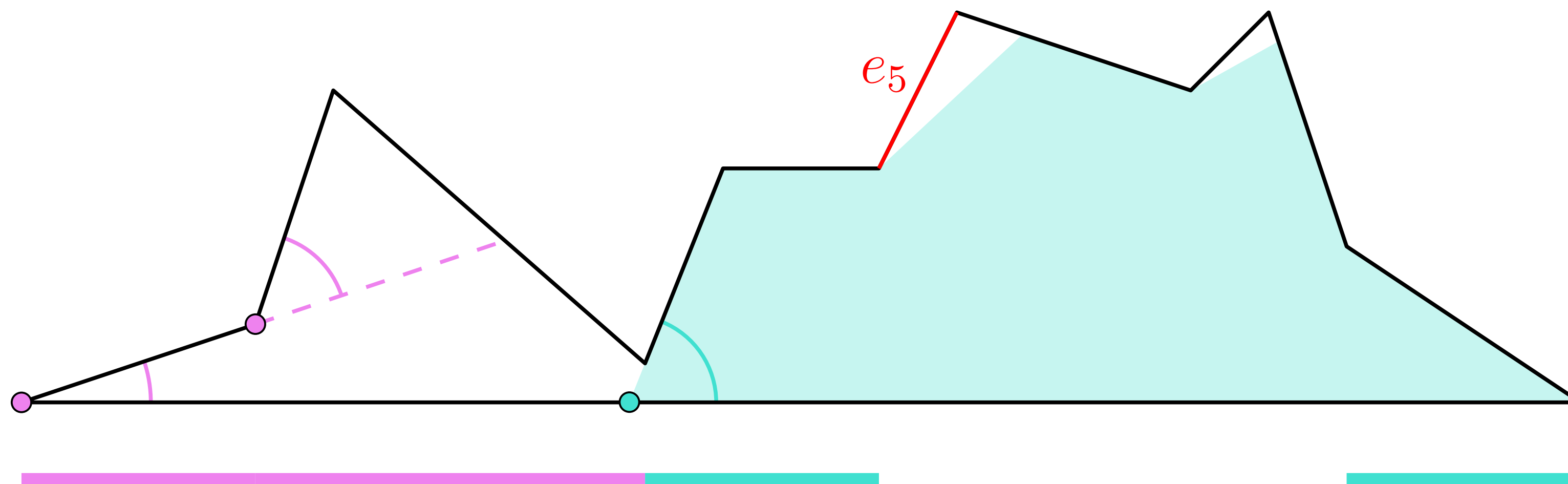
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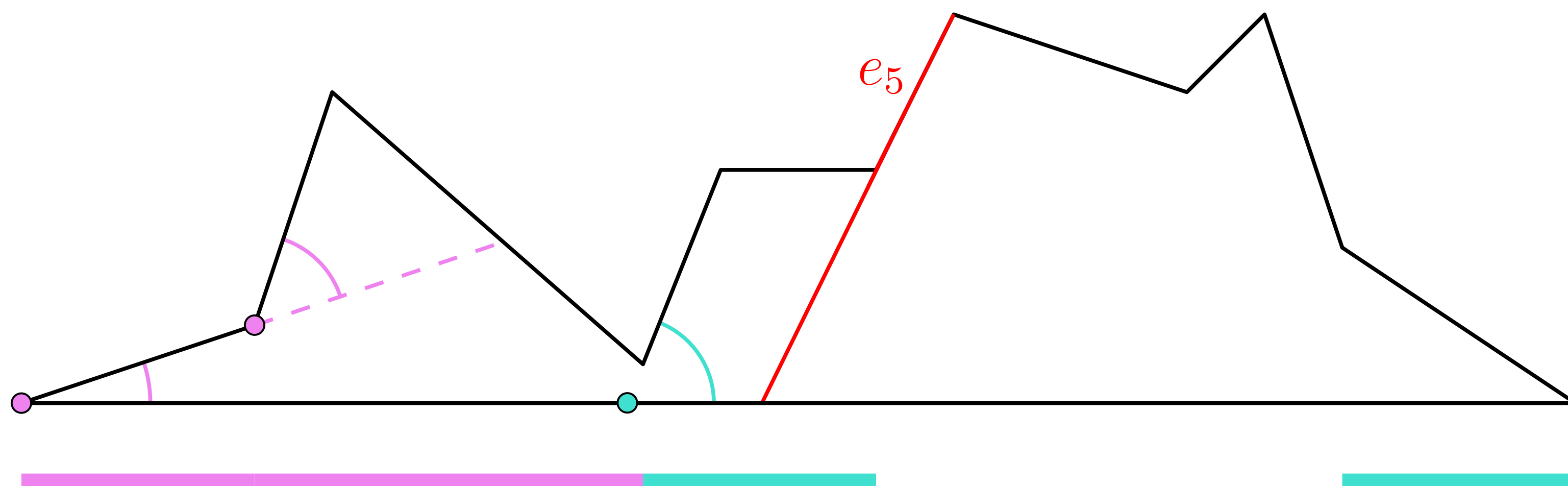
# Histogram Covering Algorithm – Example



Case	Condition
1	$e_i$ is covered
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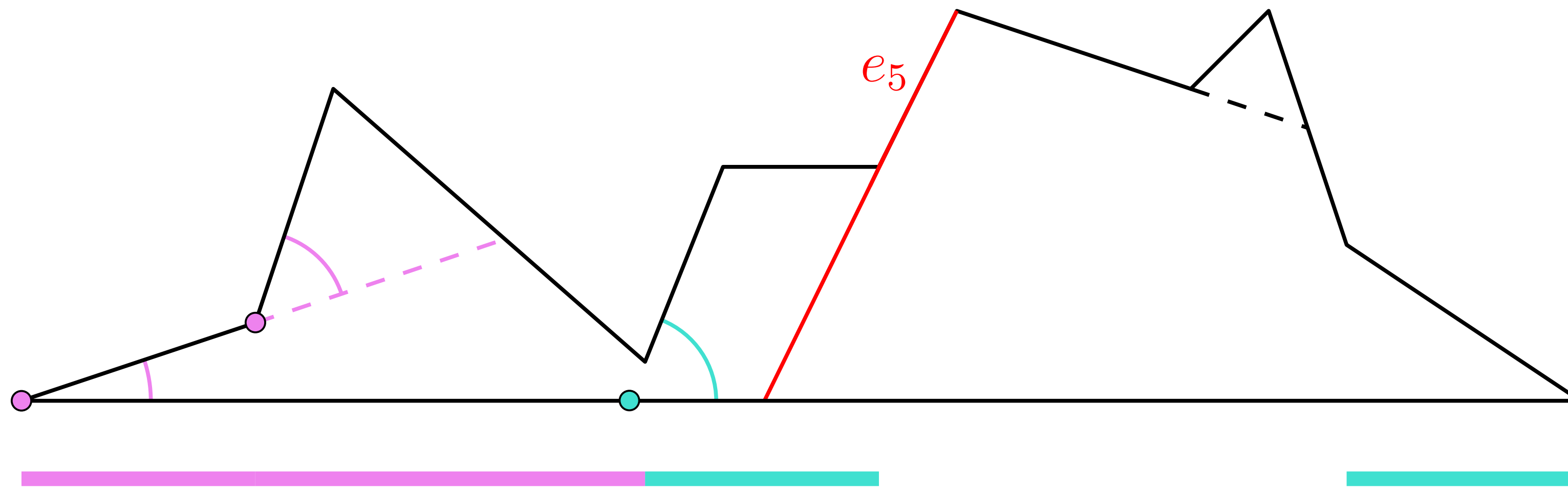


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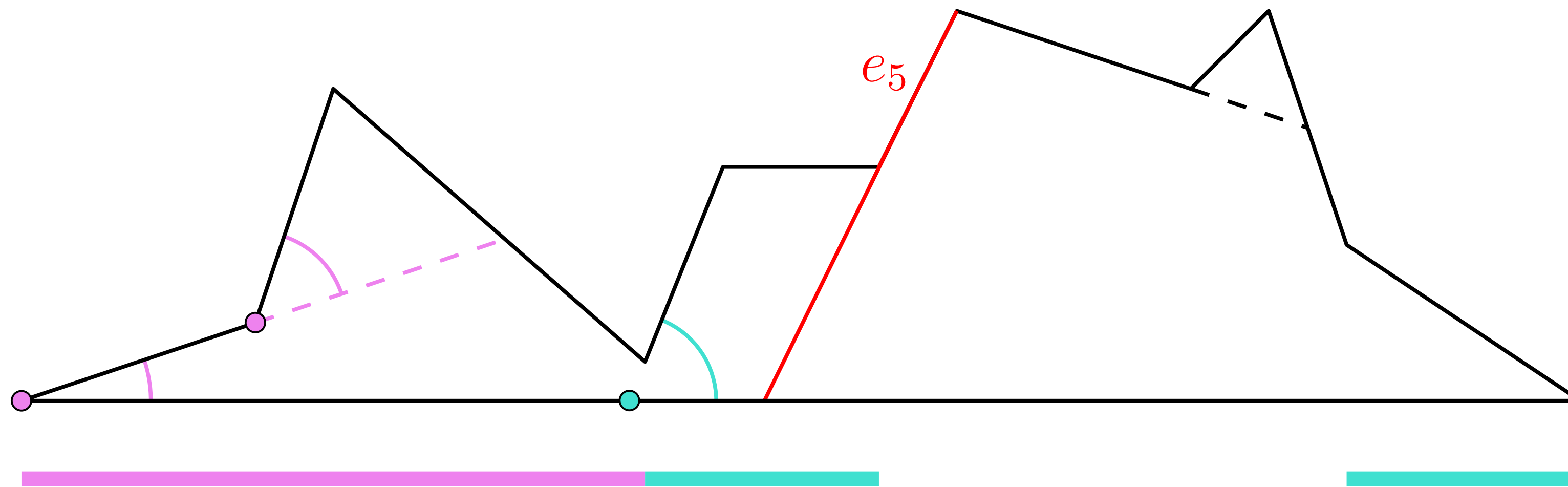
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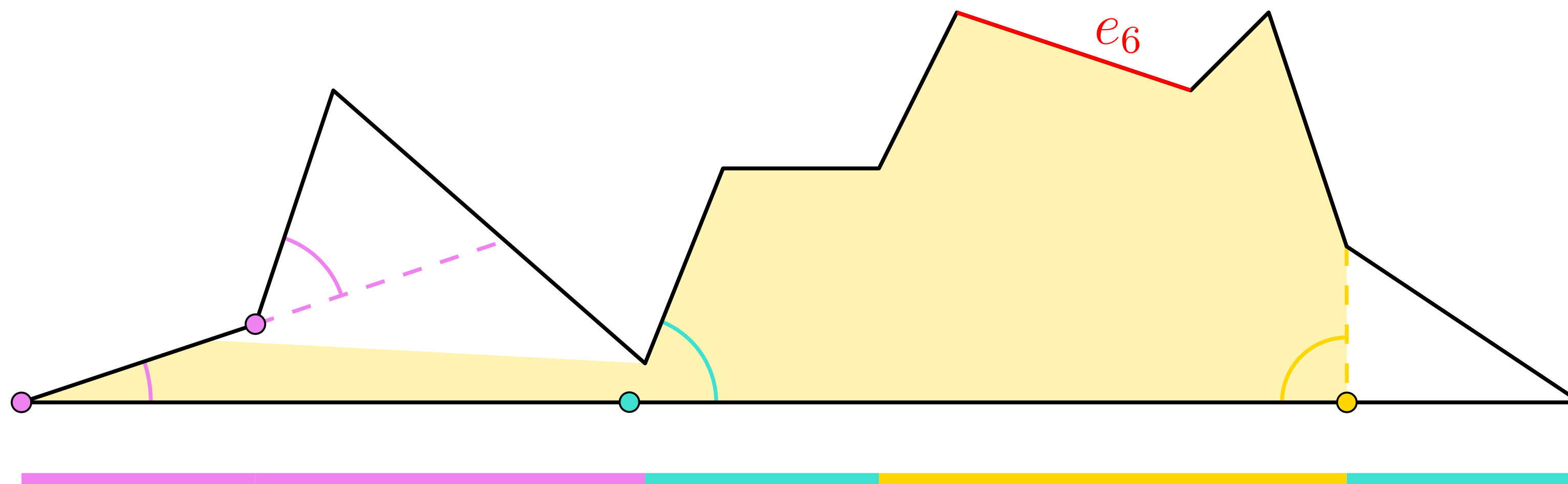
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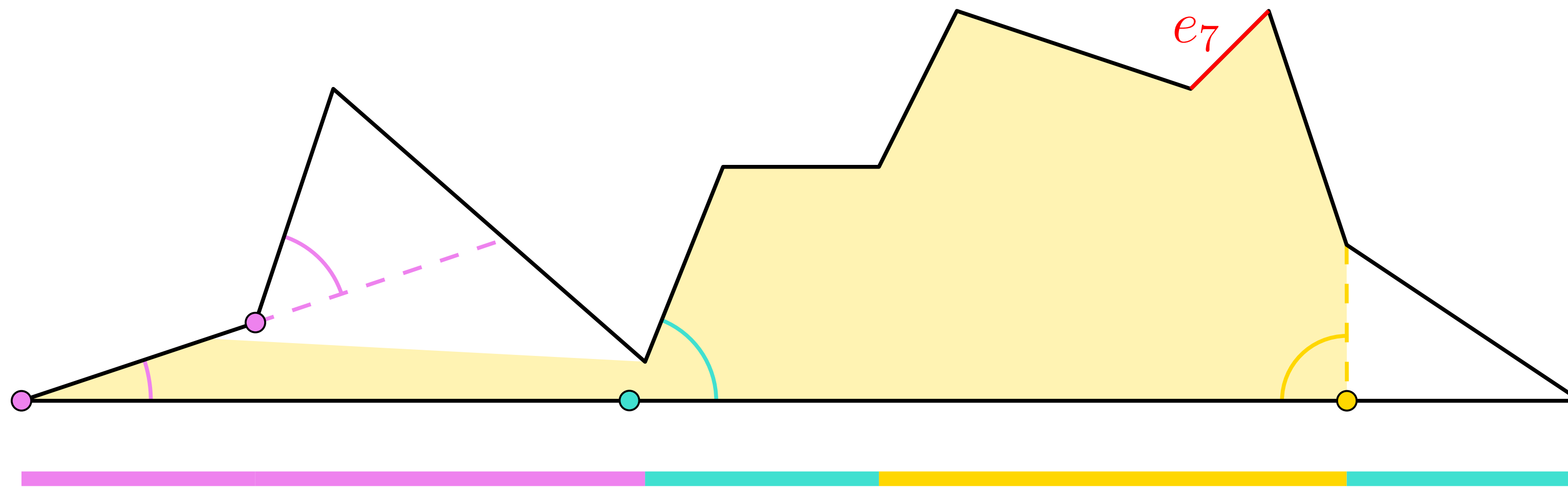


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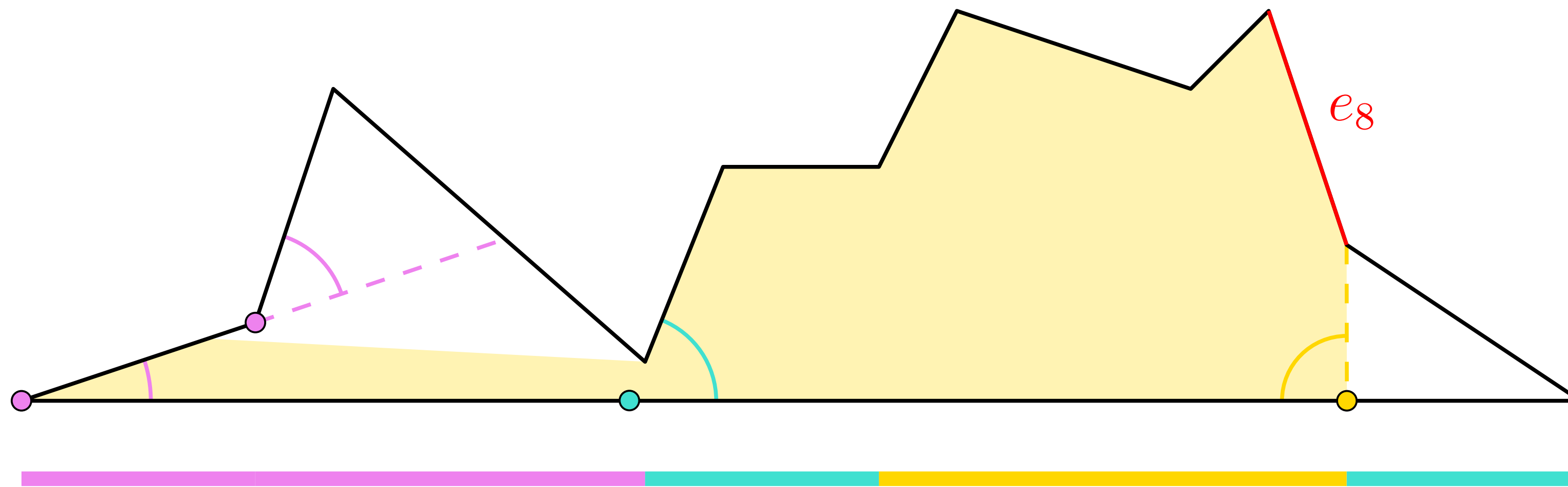
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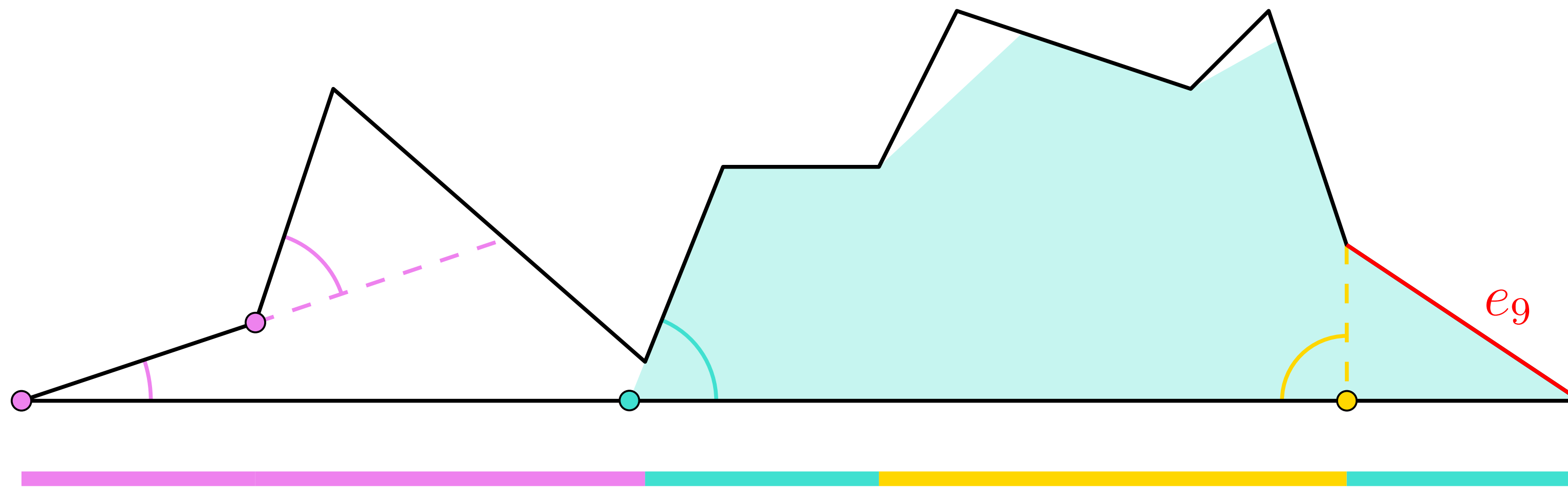
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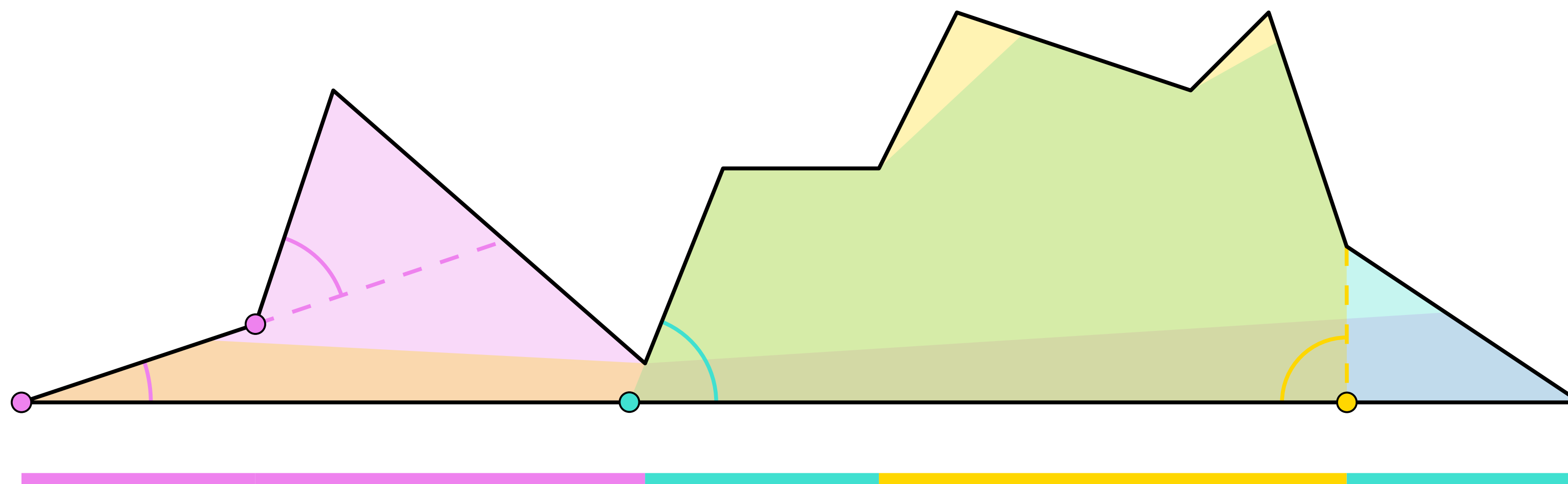
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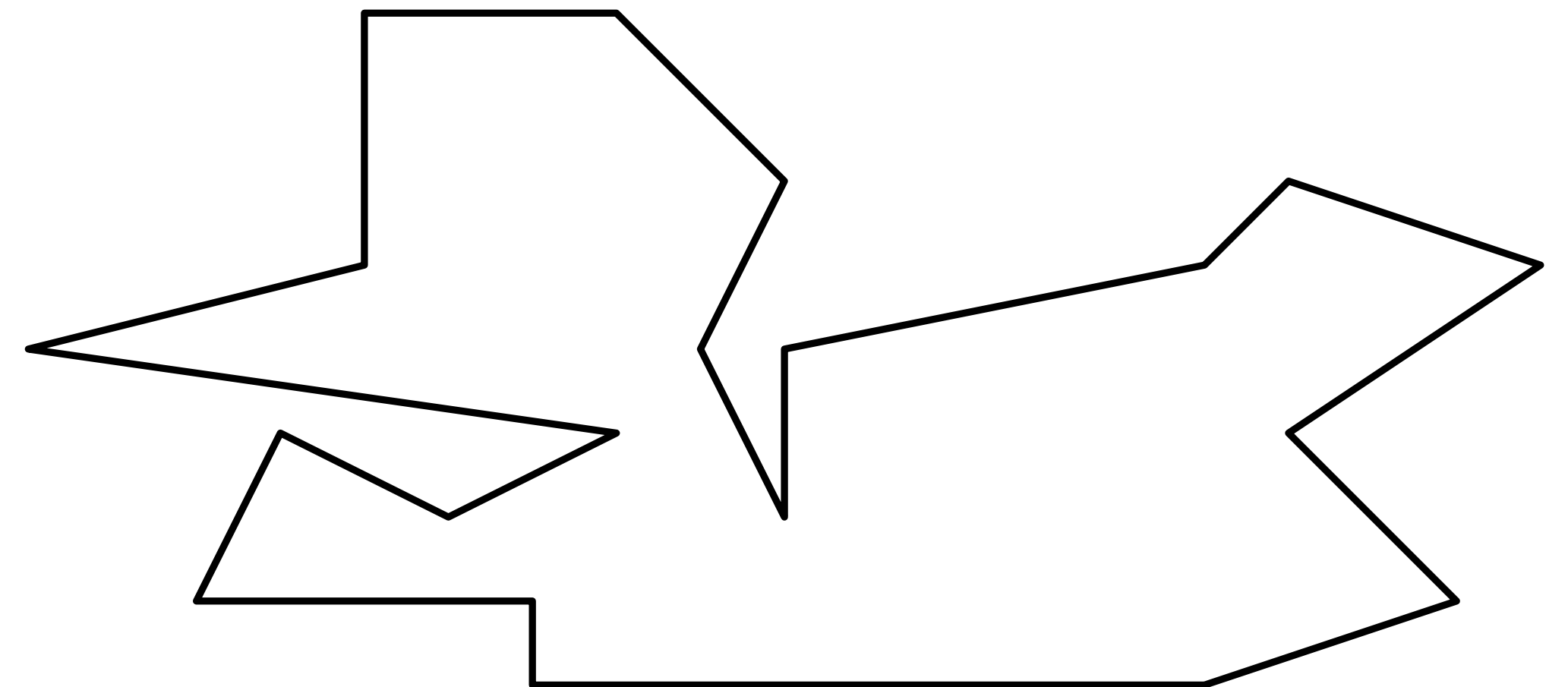
Equilateral Triangles

Histograms

Upper Bound for **Simple Polygons**

Duality

Conclusion



Introduction

Equilateral Triangles

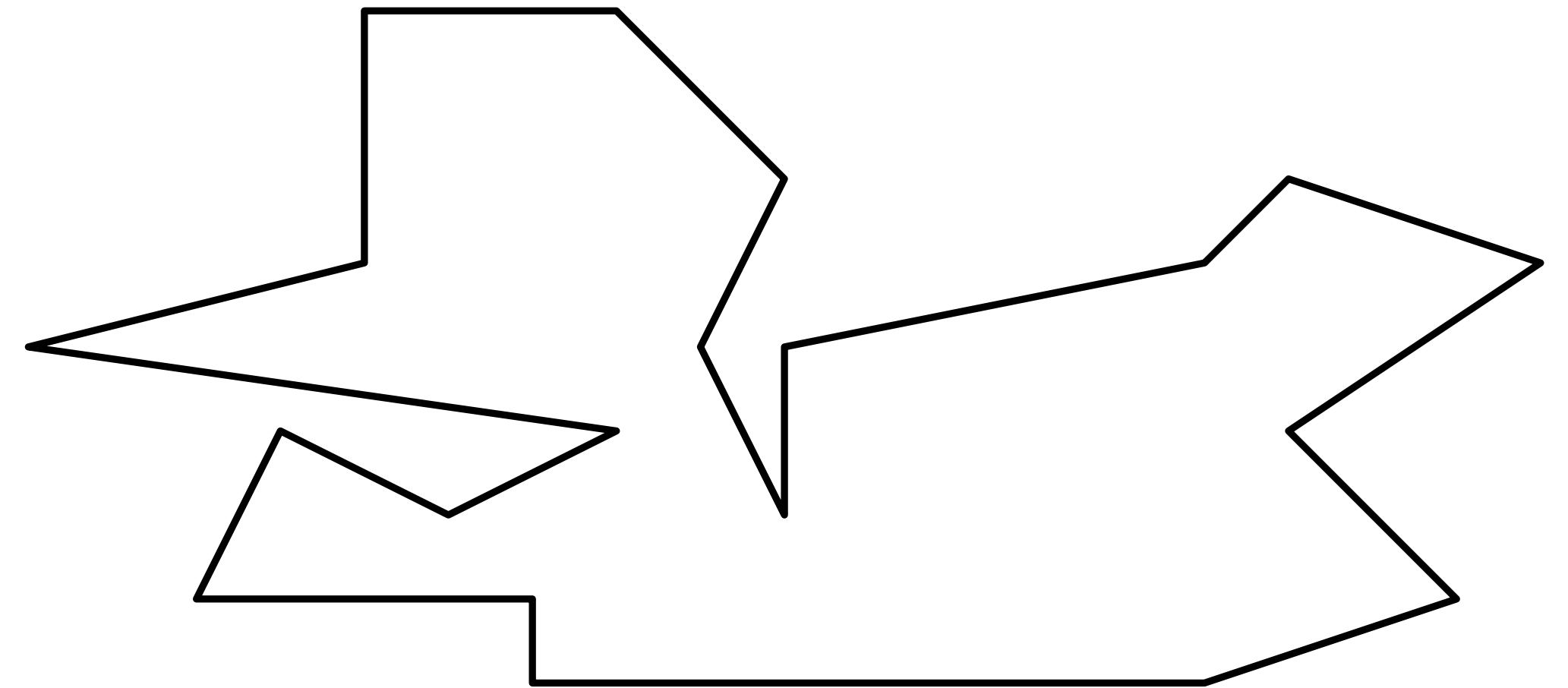
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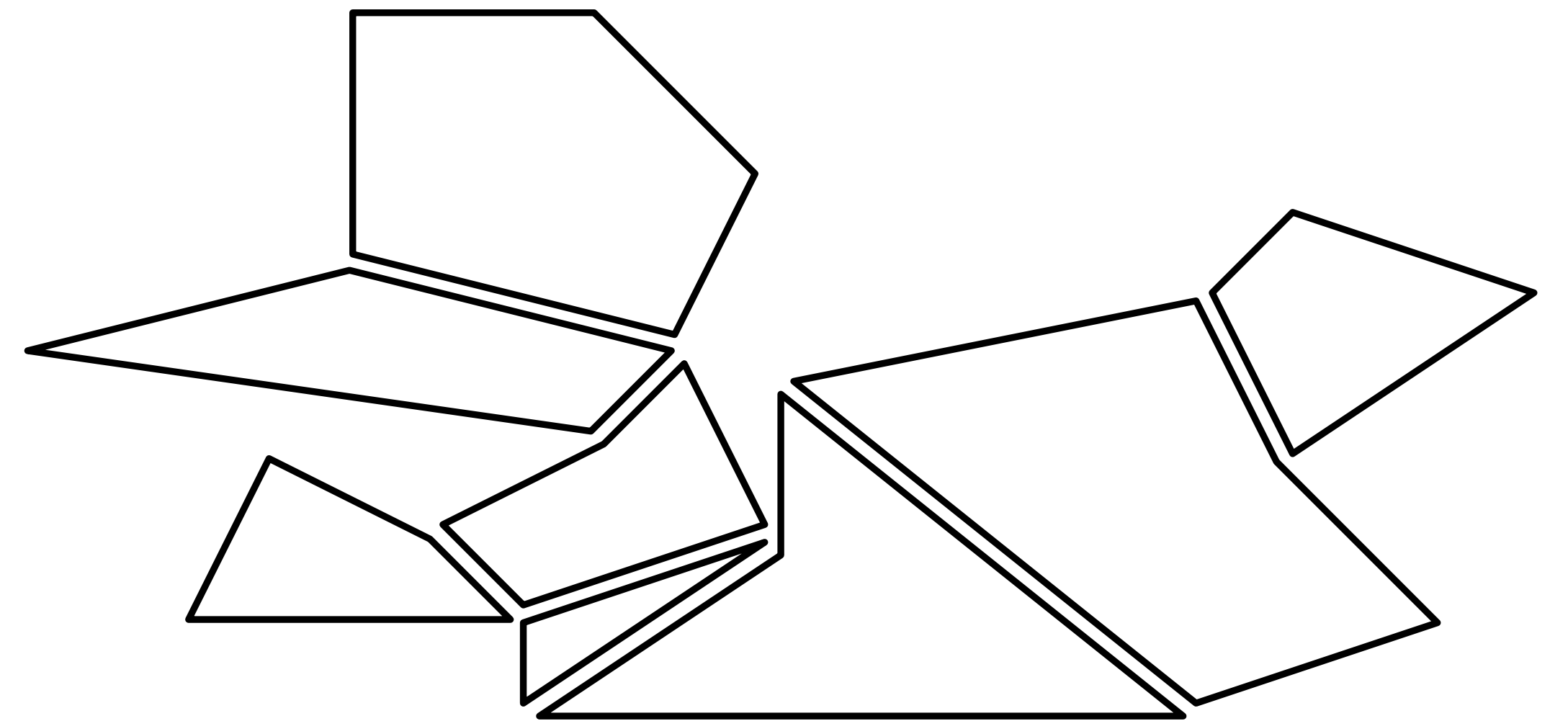
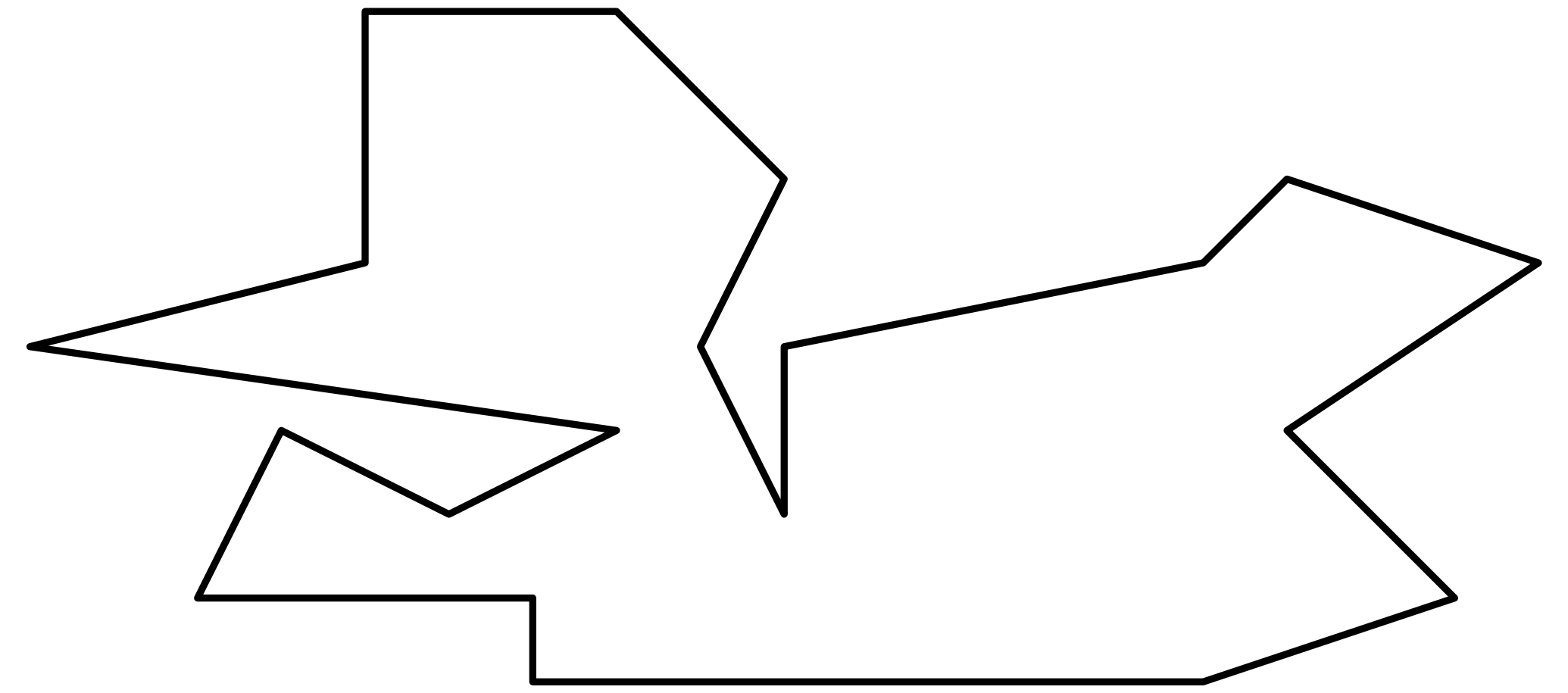
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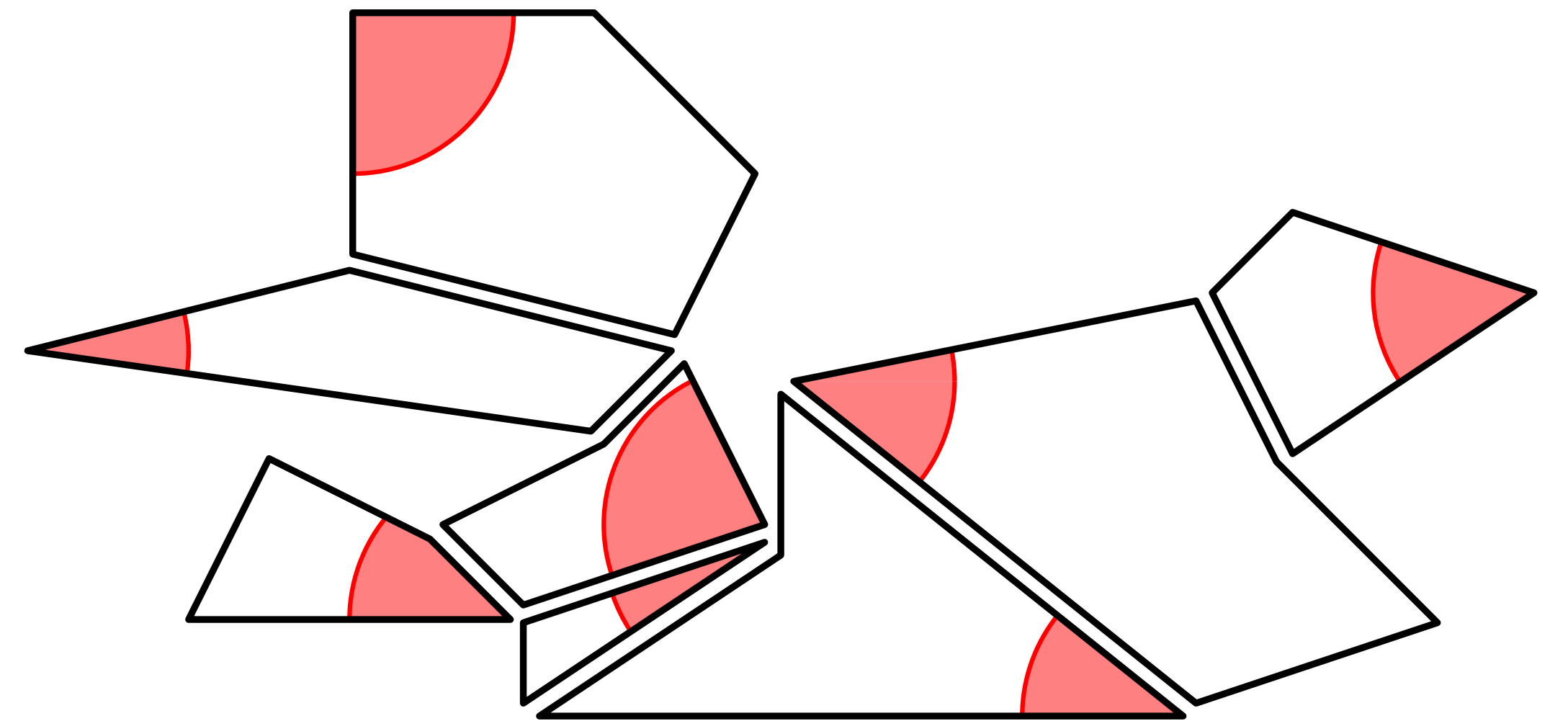
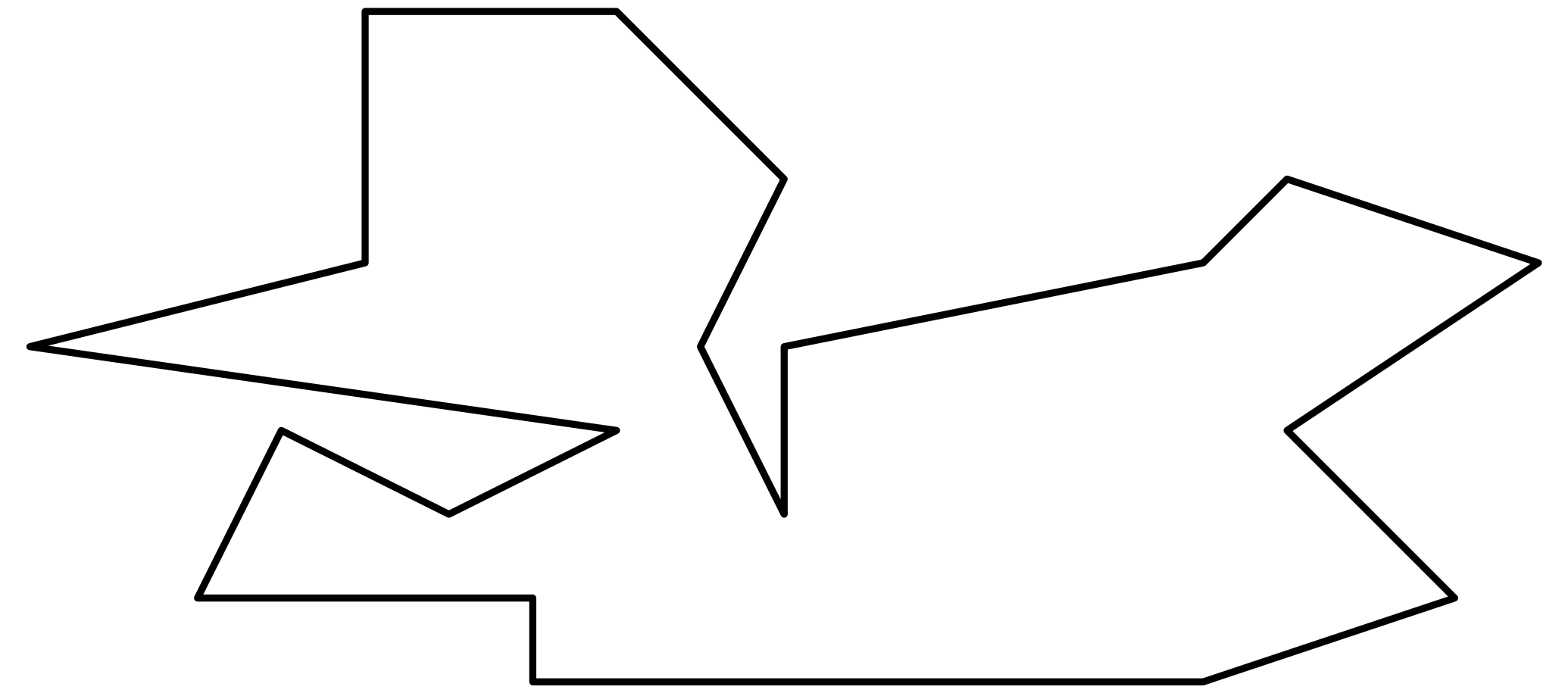


Introduction  
Equilateral Triangles  
Histograms

## Upper Bound for **Simple Polygons**

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# Upper bound for simple polygons

Upper bound for  $n \geq k + 2$  of

$$\left(1 + \frac{1}{k}\right) (n - 2) \frac{\pi}{6}$$

- Proven for  $k = 2$
- Prove for  $k = 10 \rightarrow UB \leq 1.1 LB$

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**Duality** to independent circle packing

Conclusion

# Duality to independent circle packing

## Independent Circle Packing Problem

Instance: A polygon  $P$

Wanted: A set of independent circles in  $P$

Maximize: The minimum required angle to cover all circles



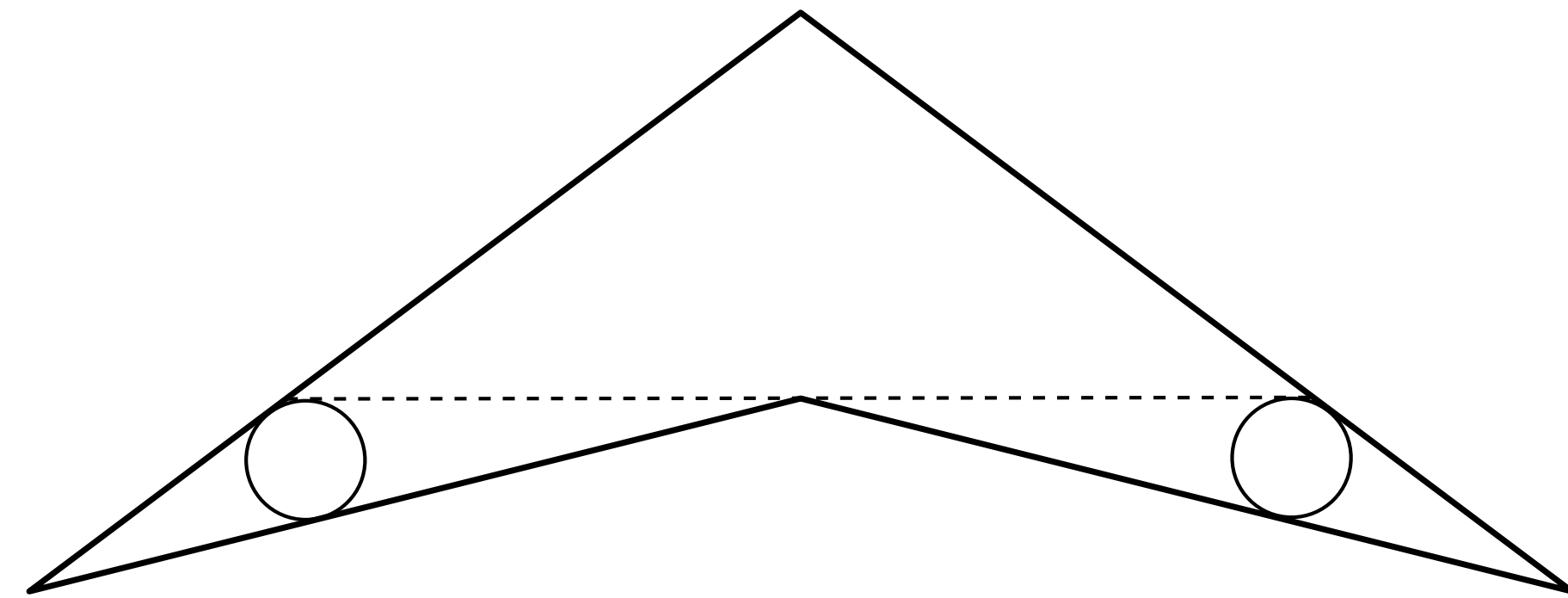
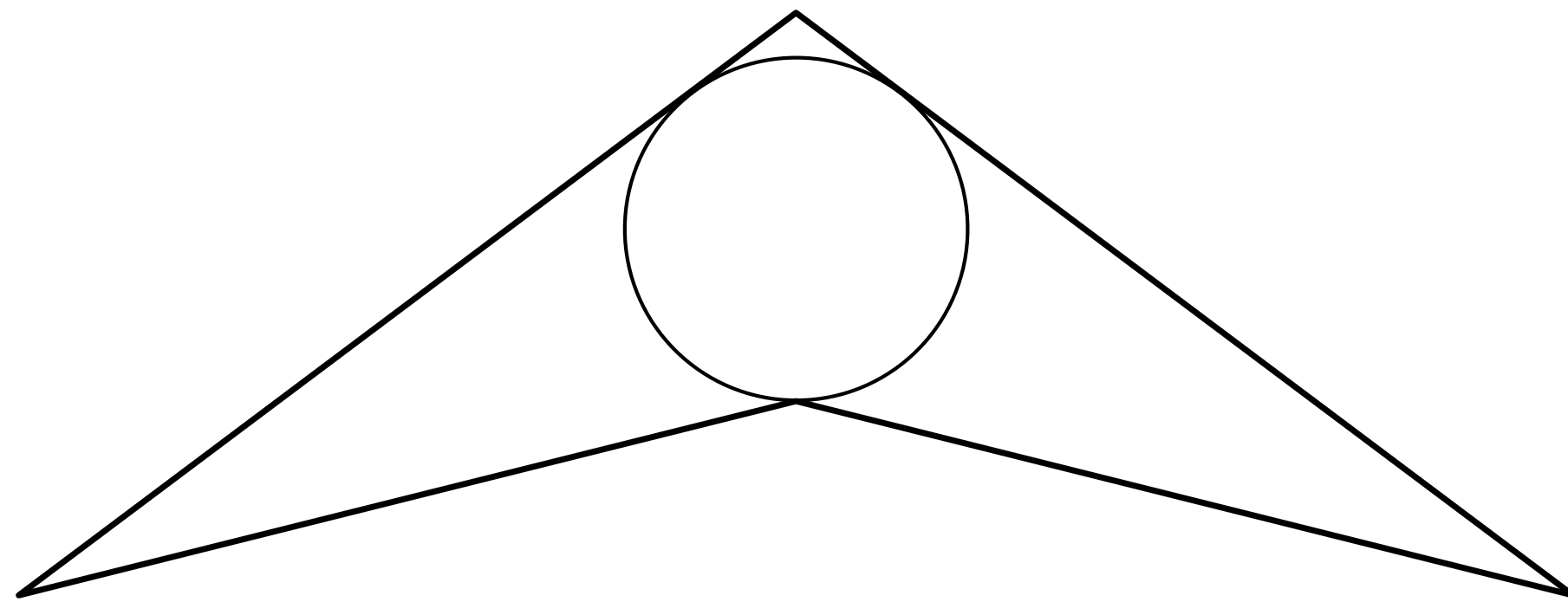
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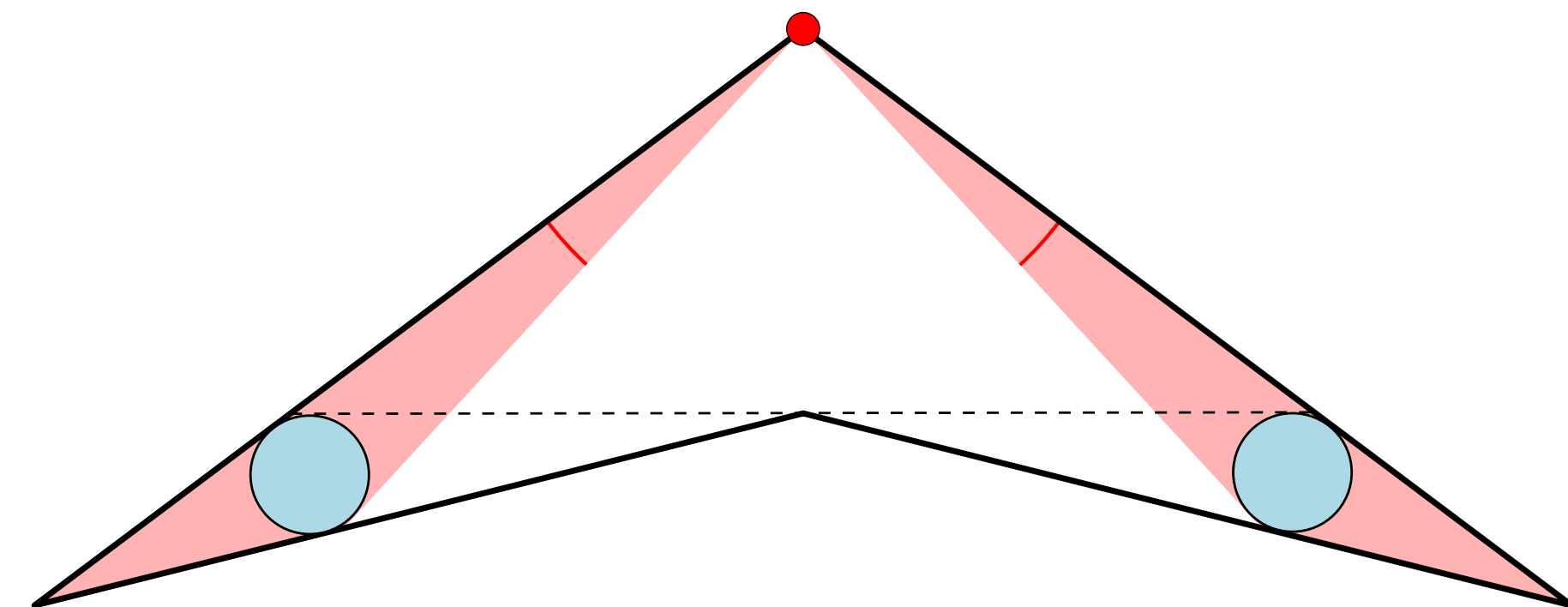
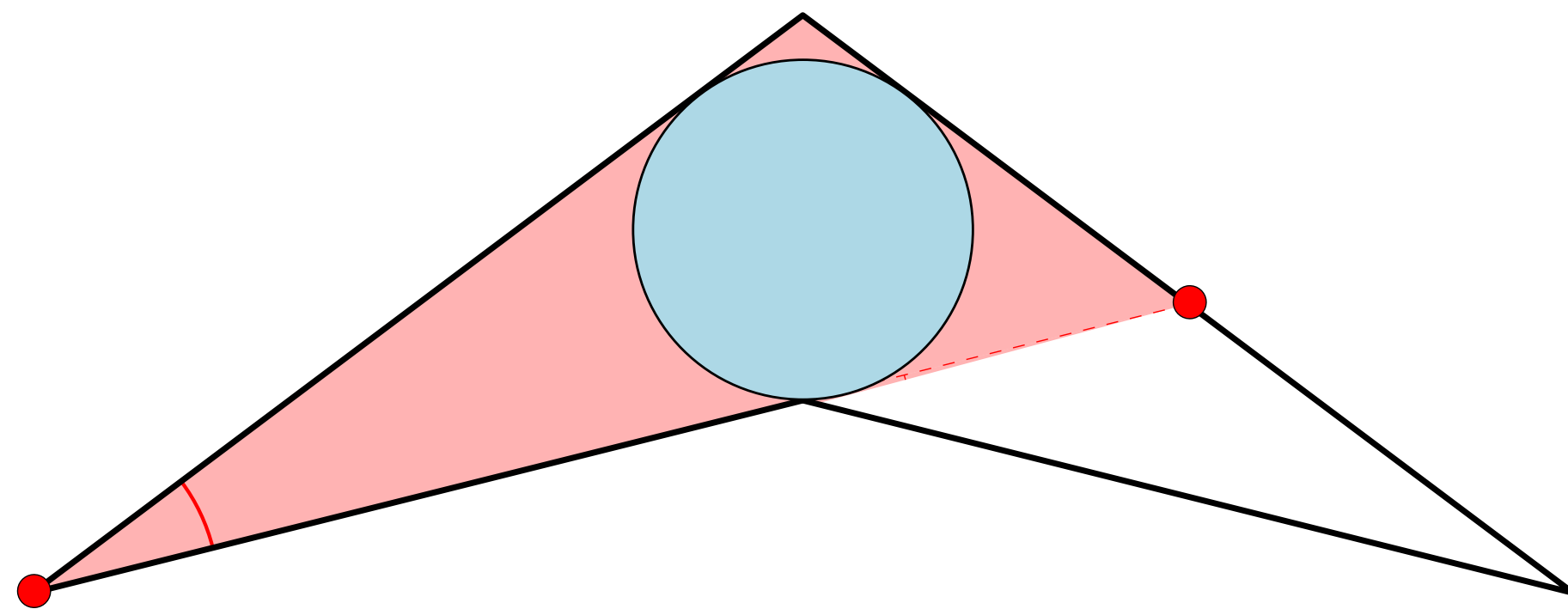
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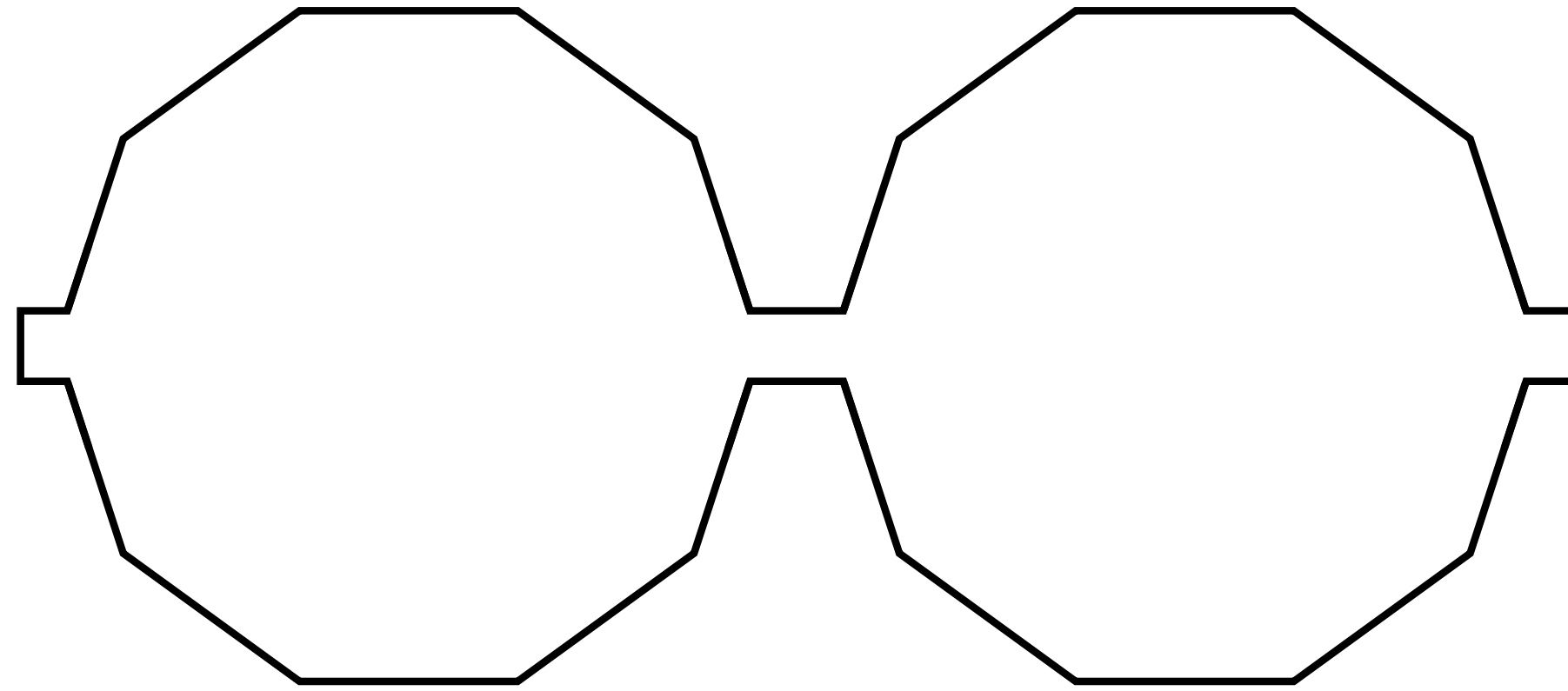
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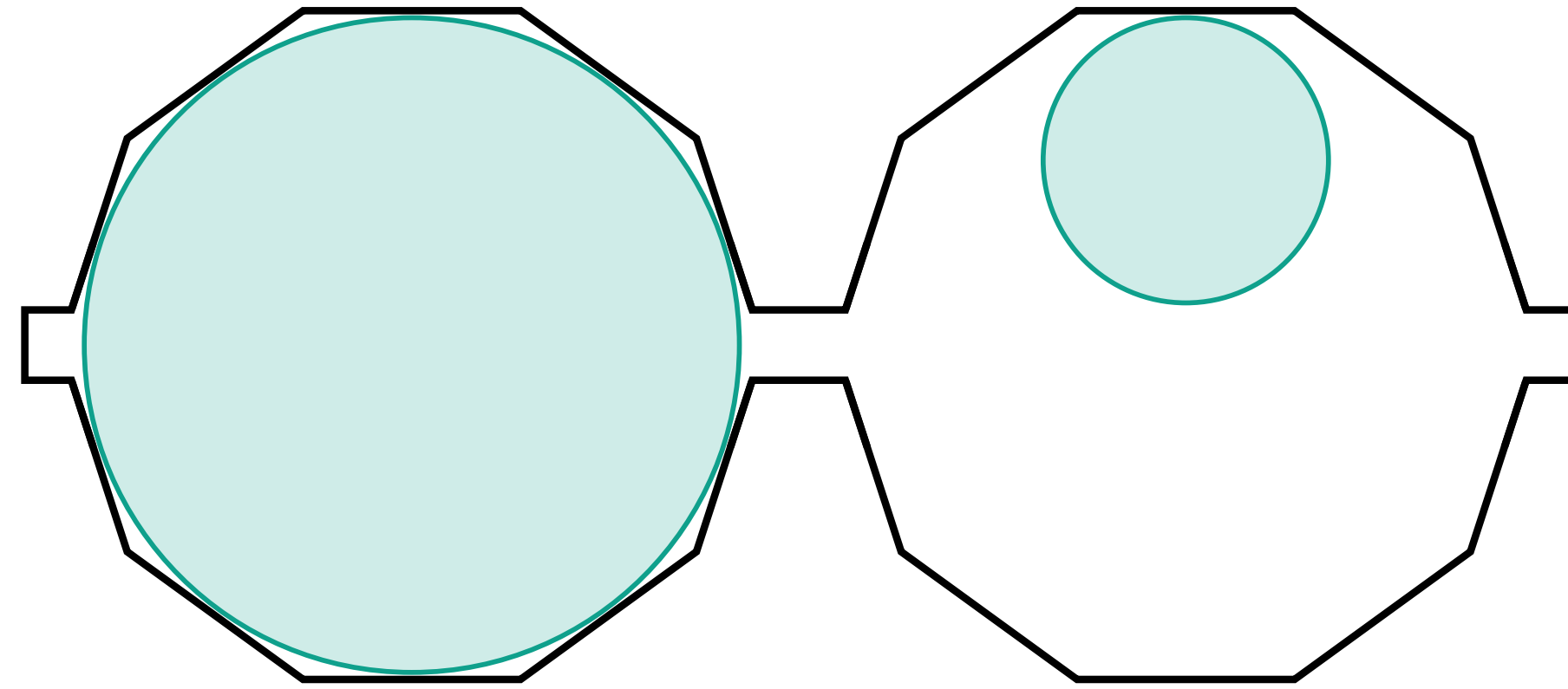
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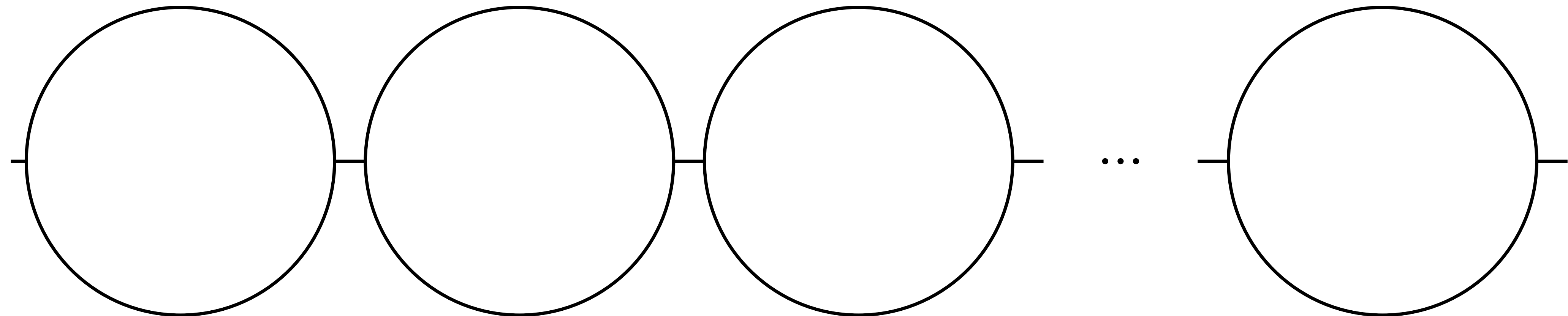
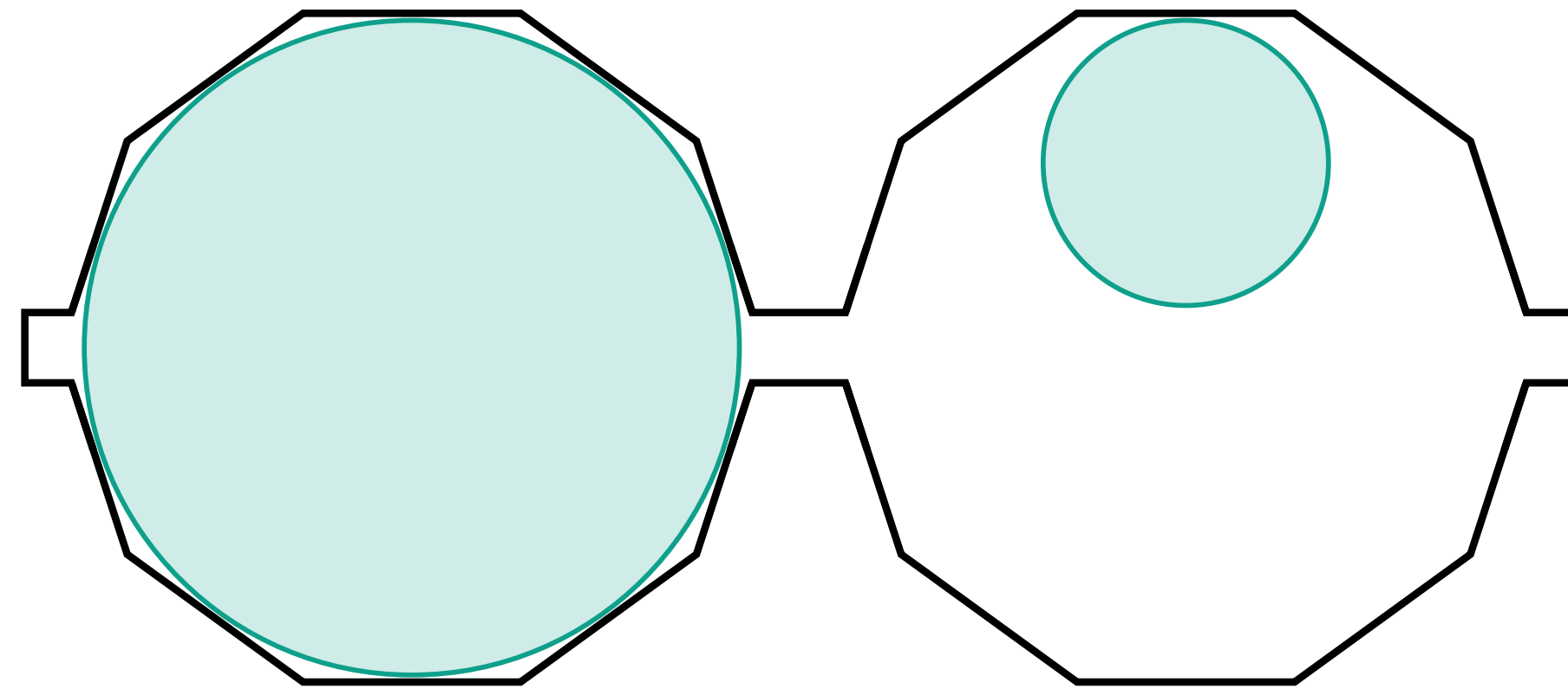
# Duality Gap



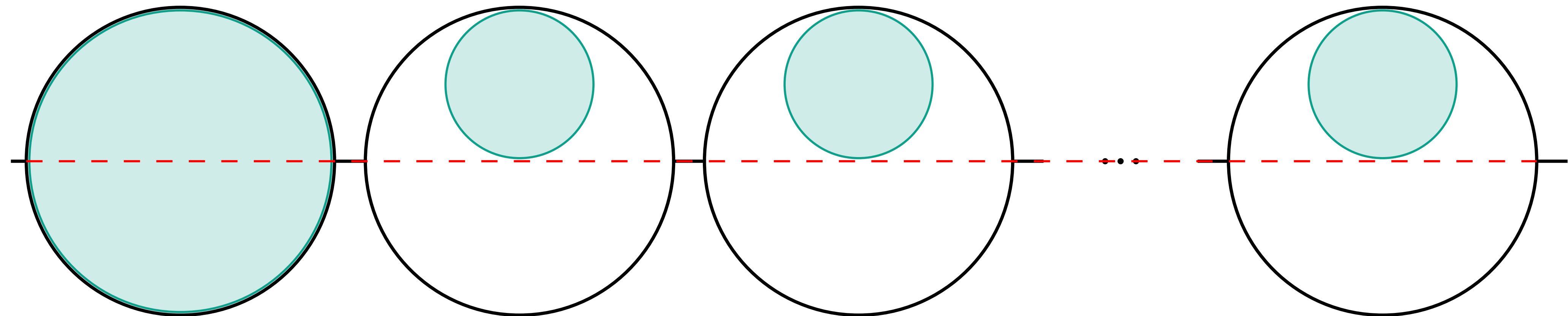
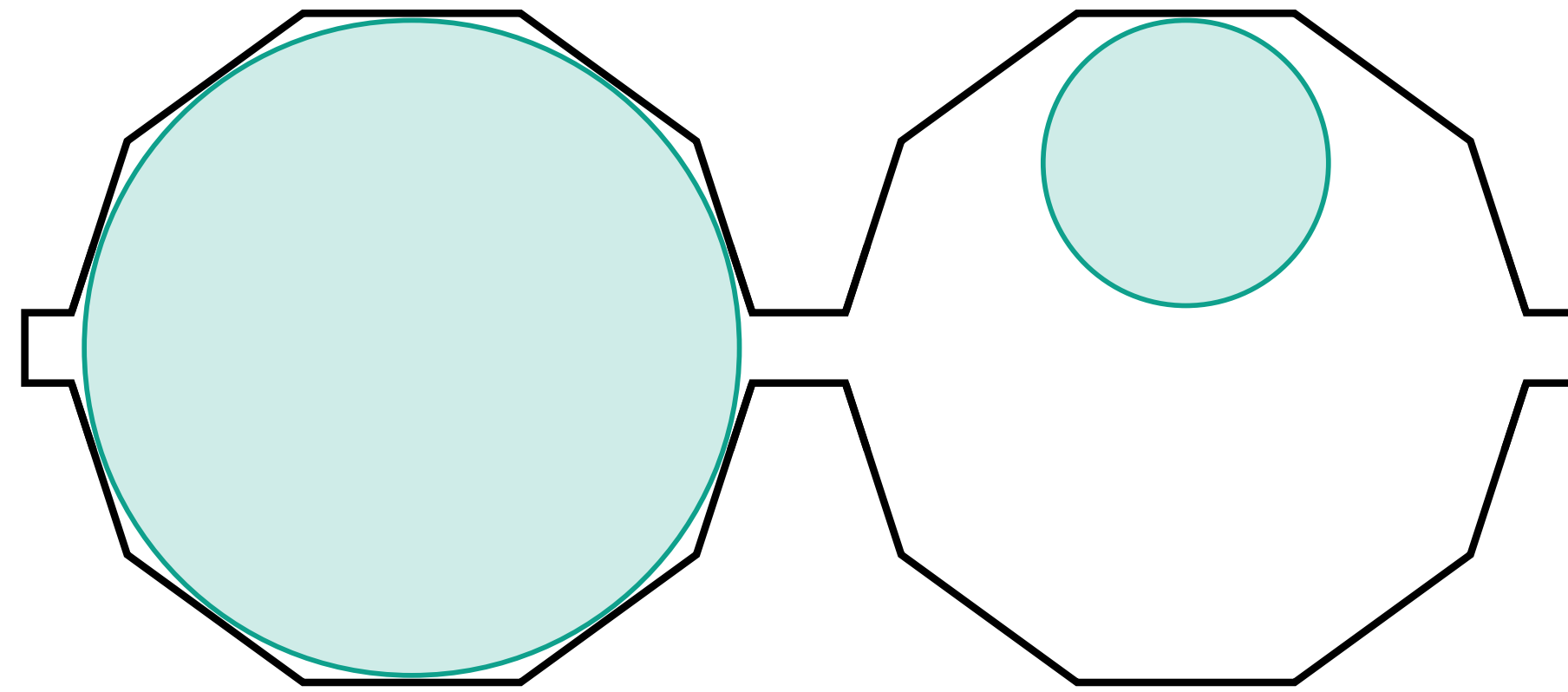
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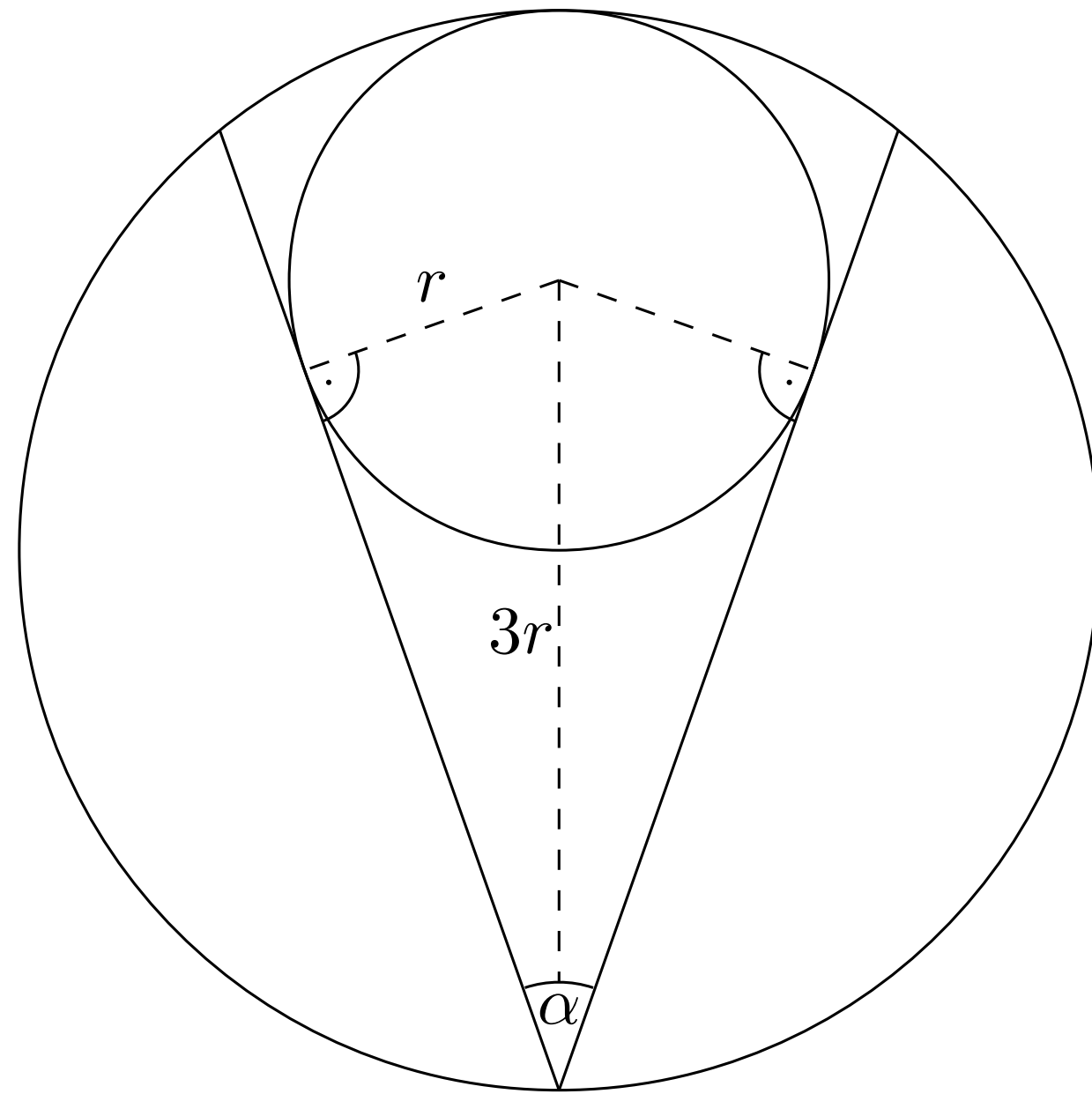
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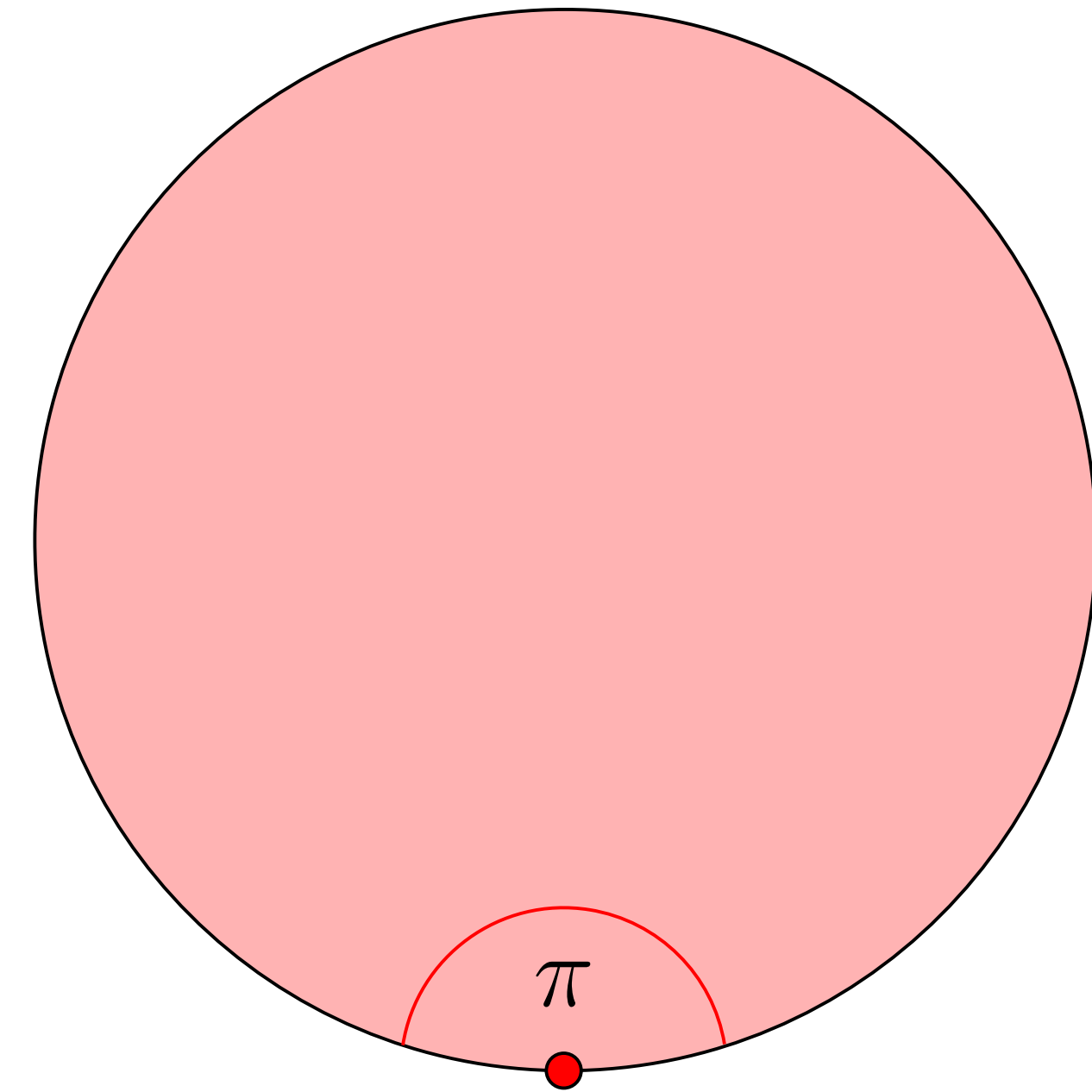
# Duality Gap



# Duality Gap



$$\text{ICP} = \alpha = 2 \arcsin \frac{1}{3}$$



$$\text{AAGP} = \pi$$

$$\frac{\text{AAGP}}{\text{ICP}} = \frac{\pi}{2 \arcsin \frac{1}{3}} \approx 4.622$$

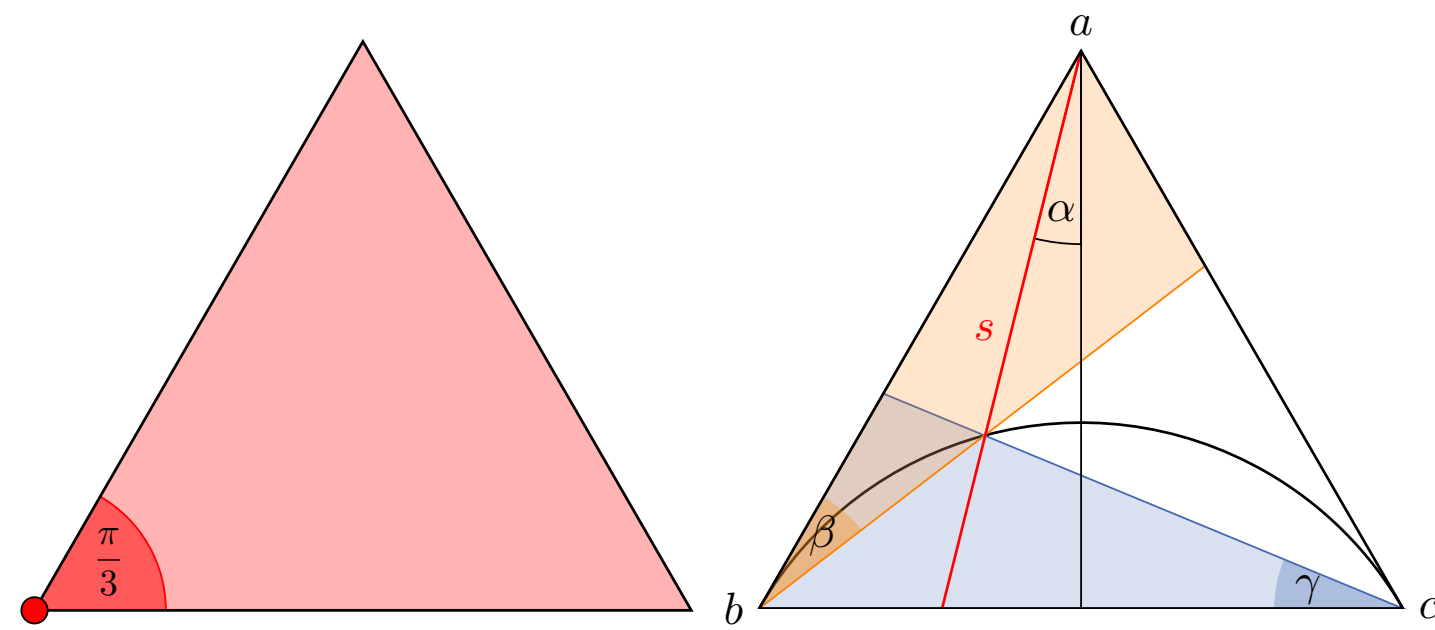
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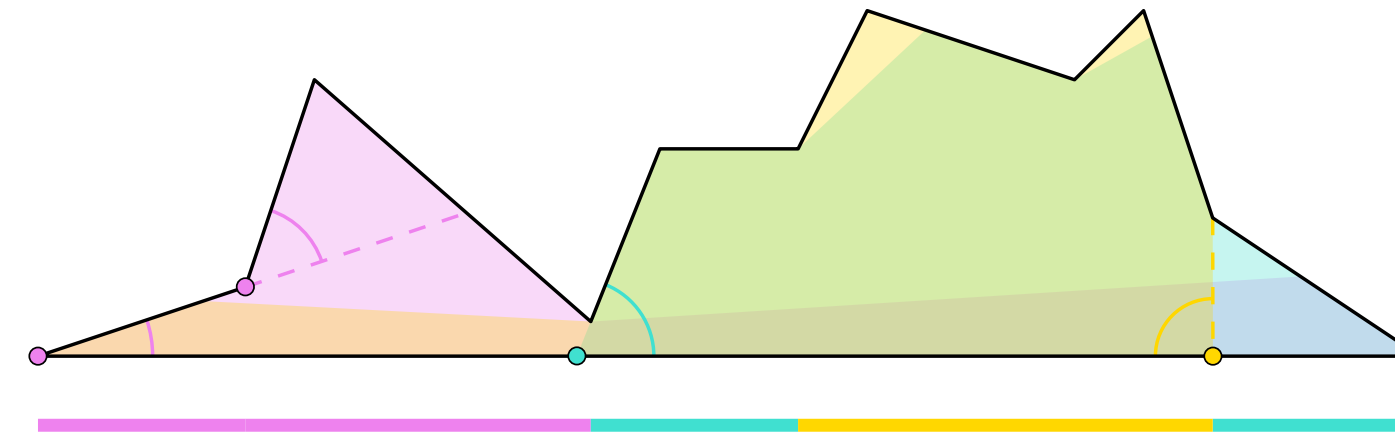
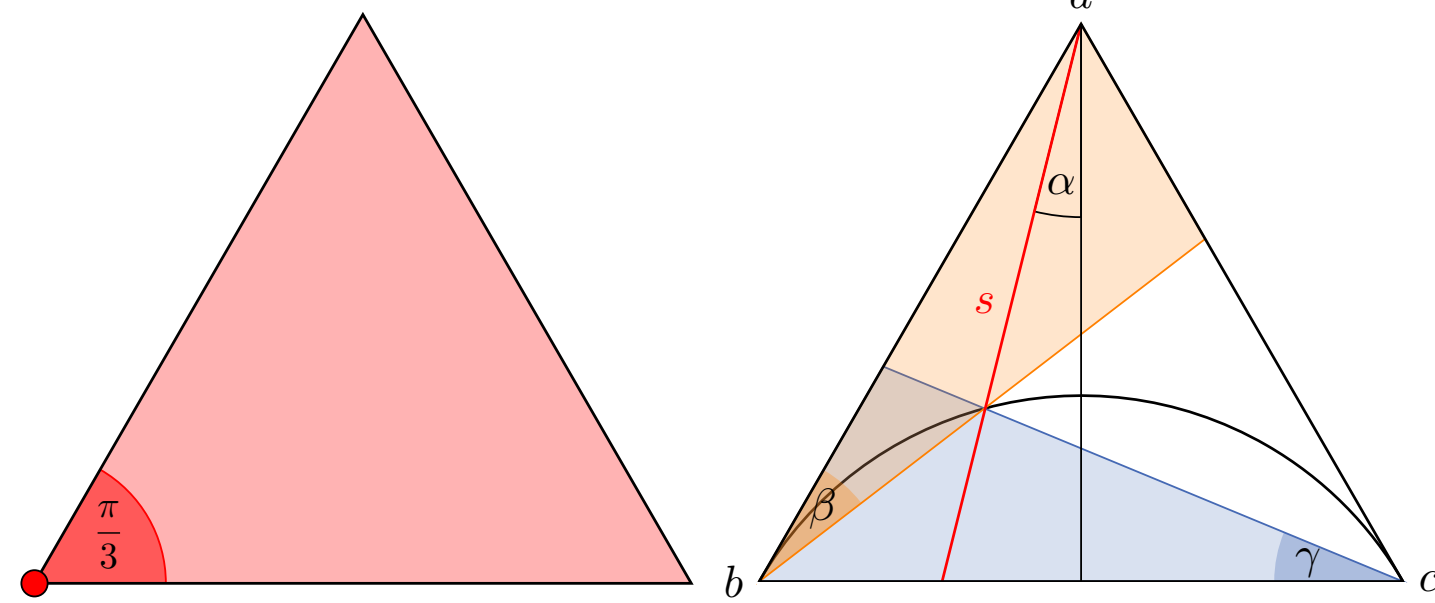
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An optimal covering of an equilateral triangle has an angle of

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There is an upper bound of

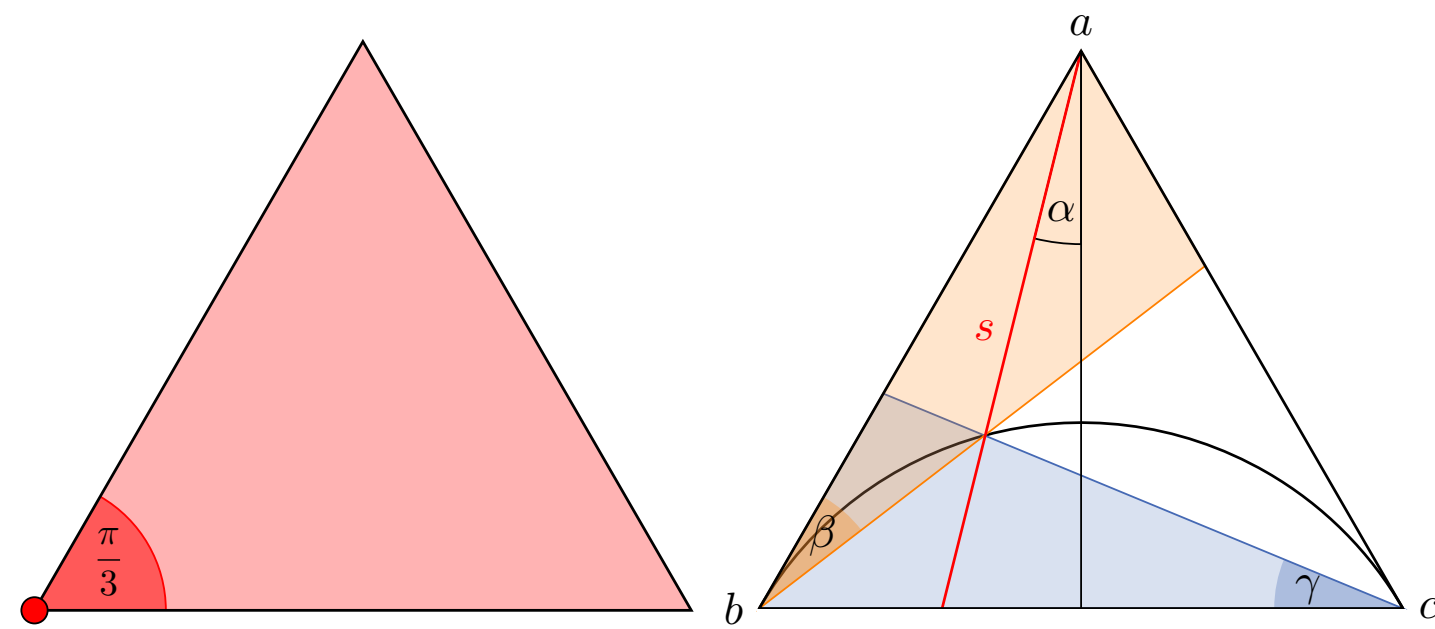
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sufficient to cover any given histogram polygon.

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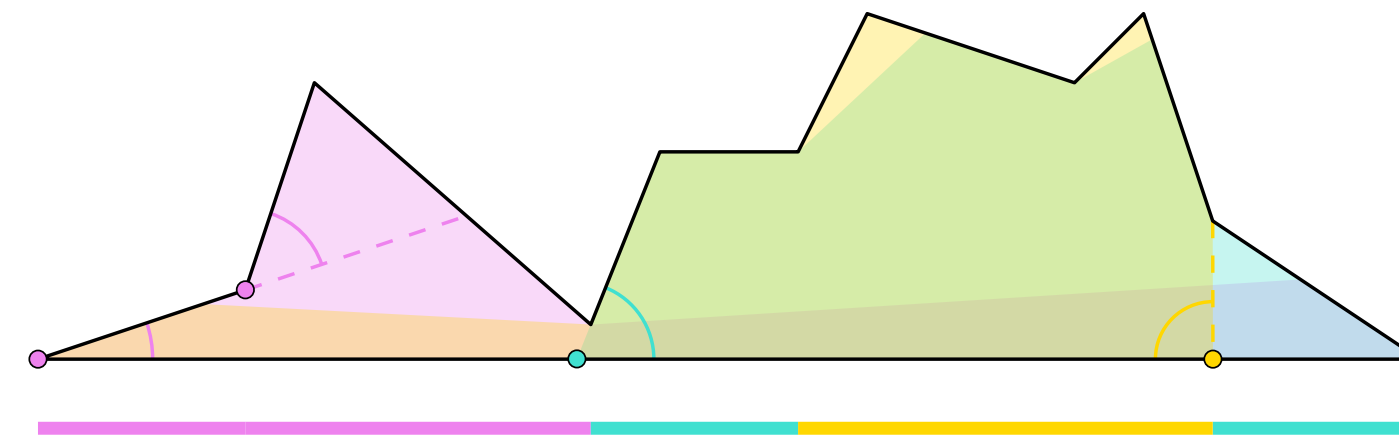
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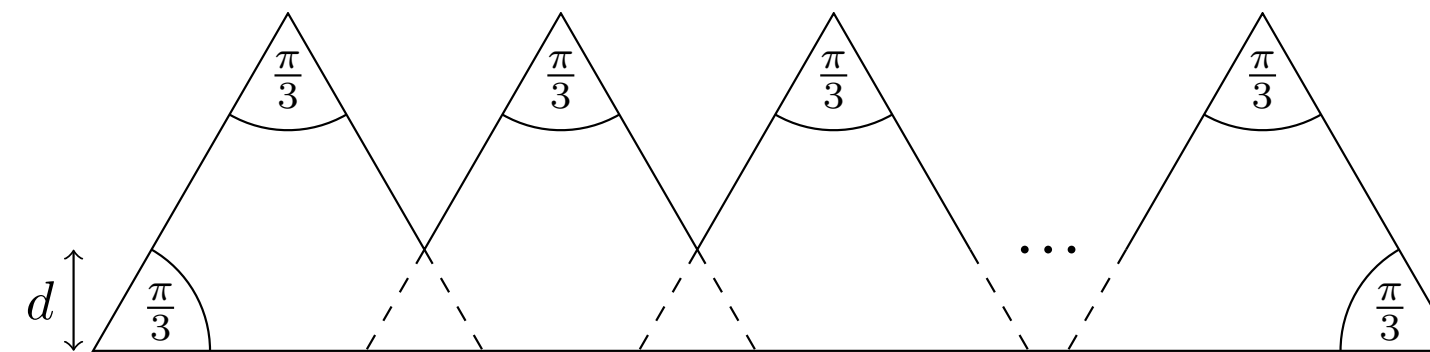
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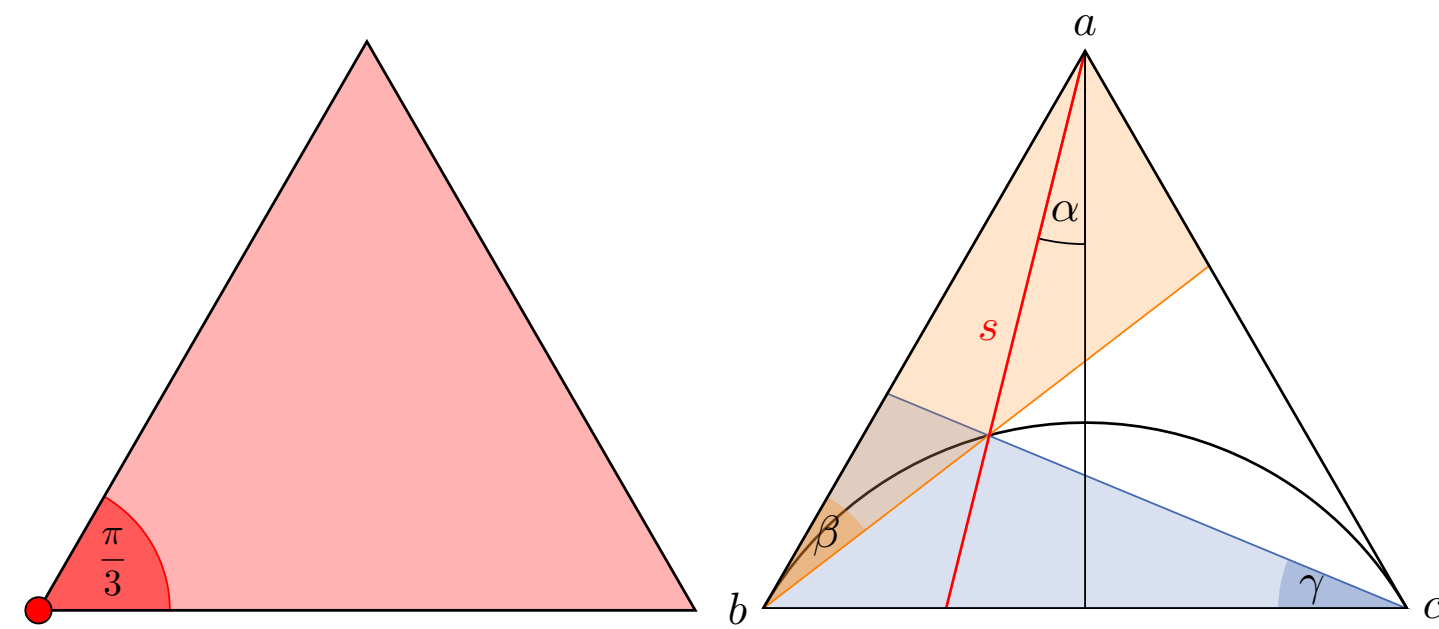
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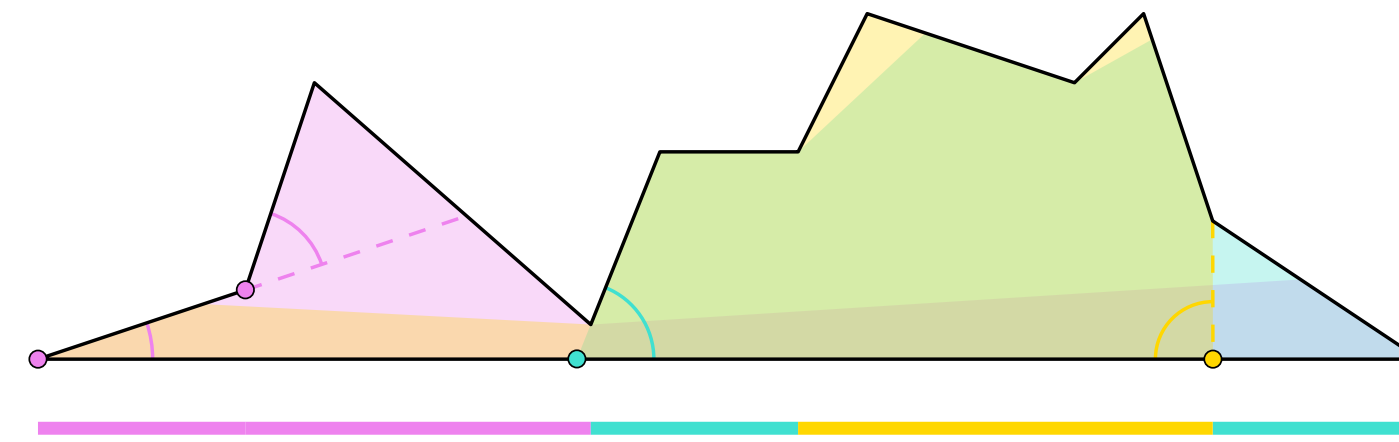
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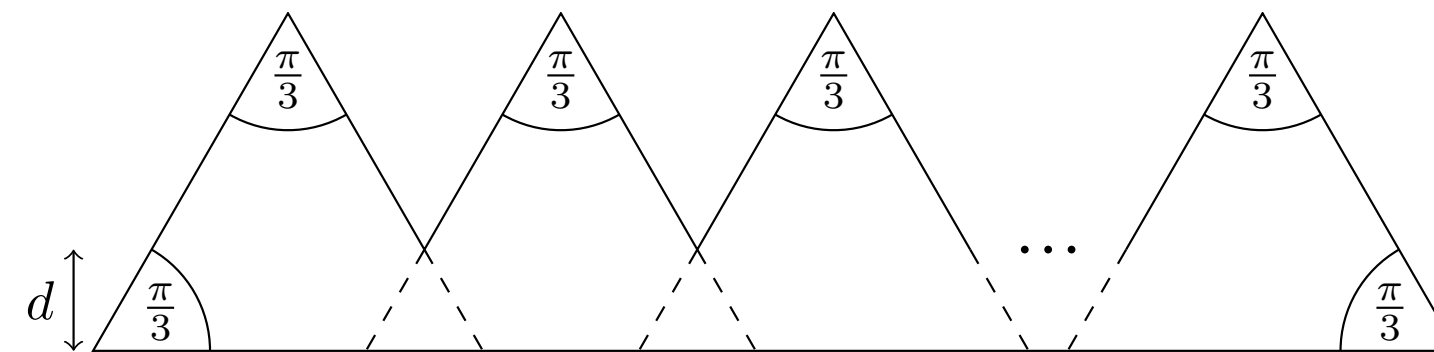
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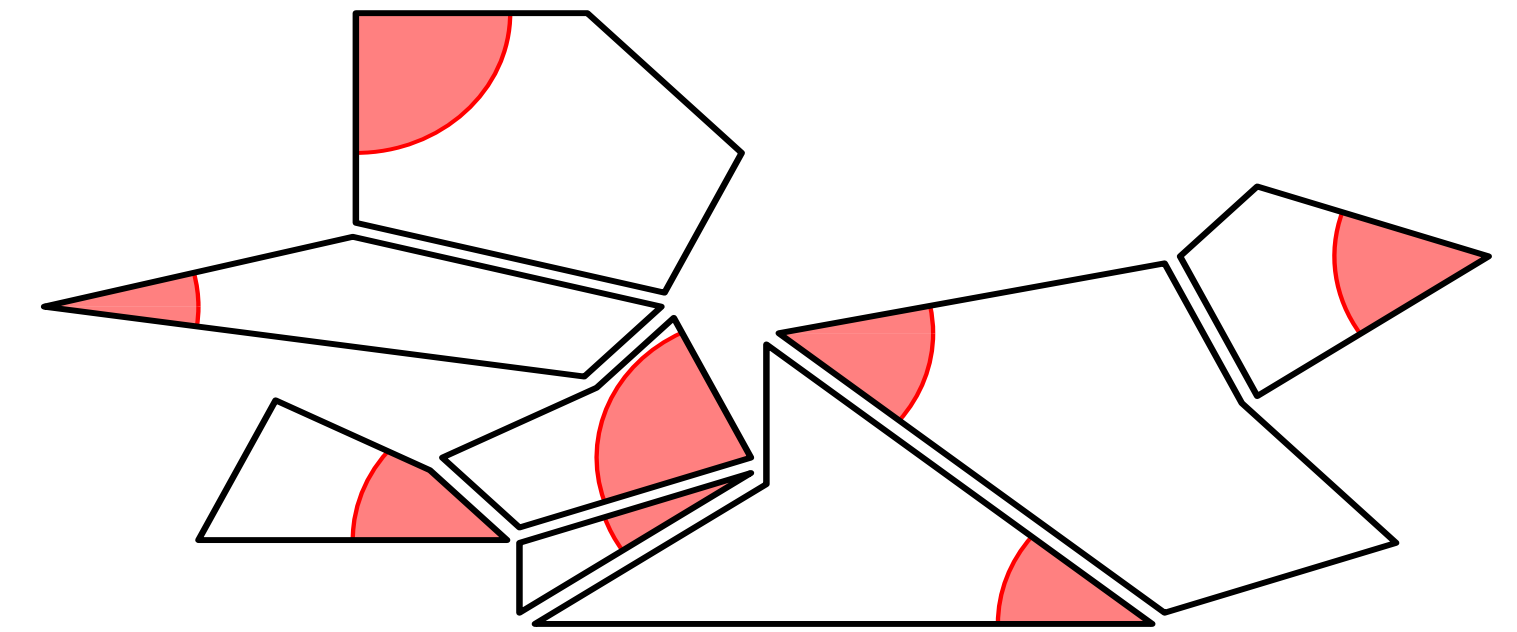
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For simple polygons, we presented an upper bound of

$$(n - 2) \frac{\pi}{4}$$

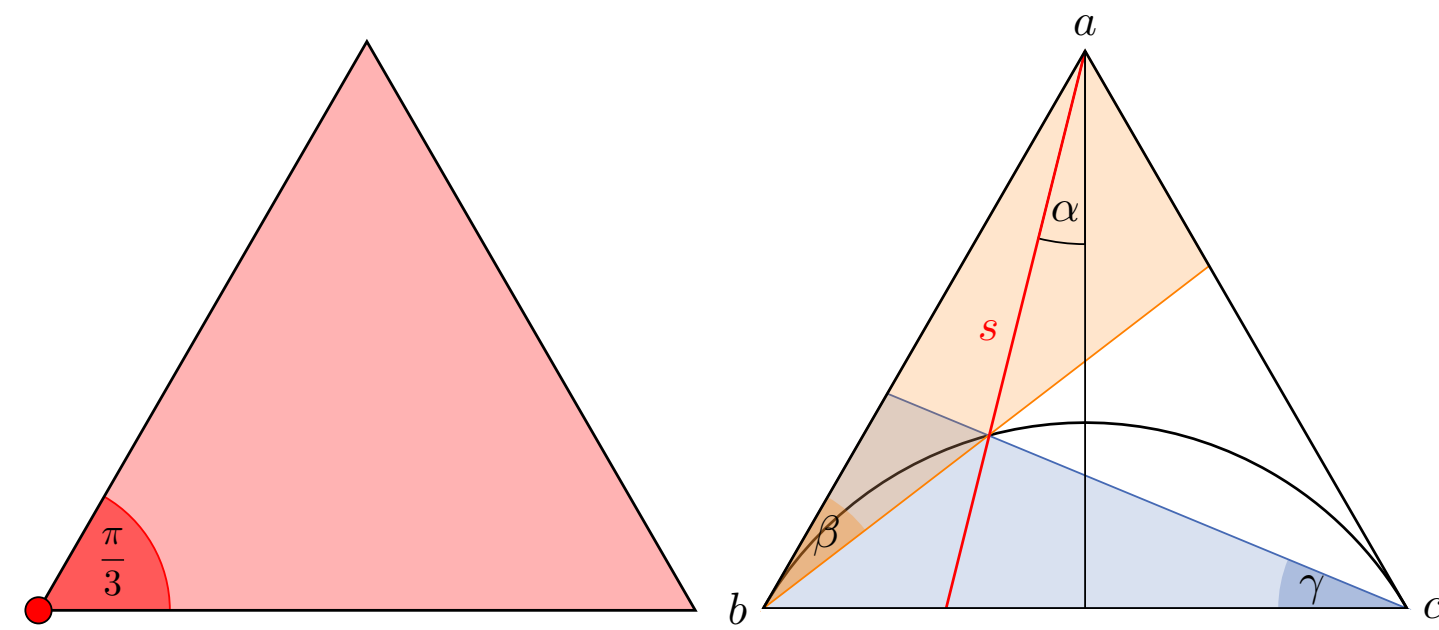
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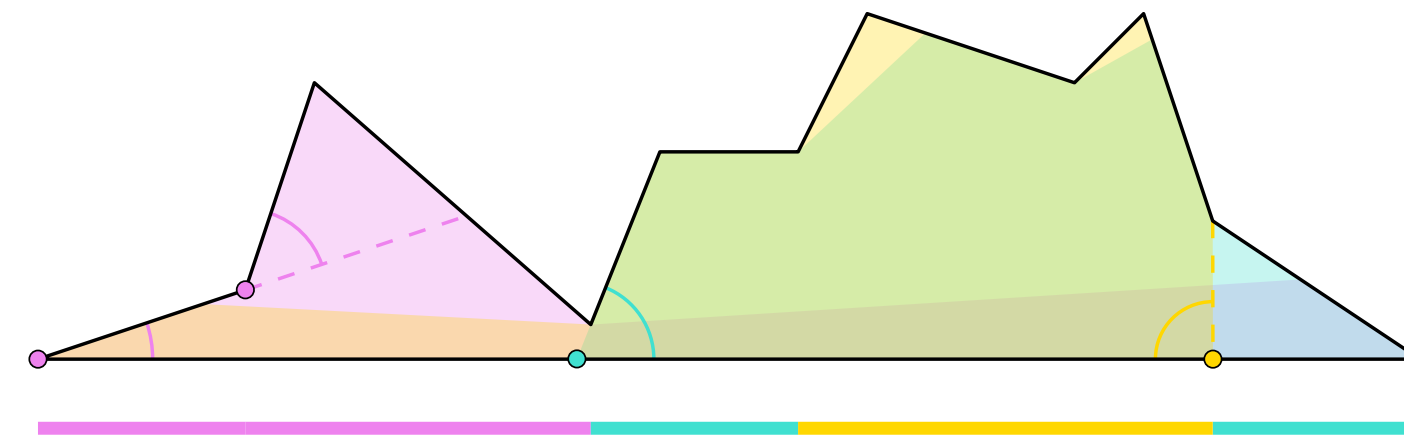
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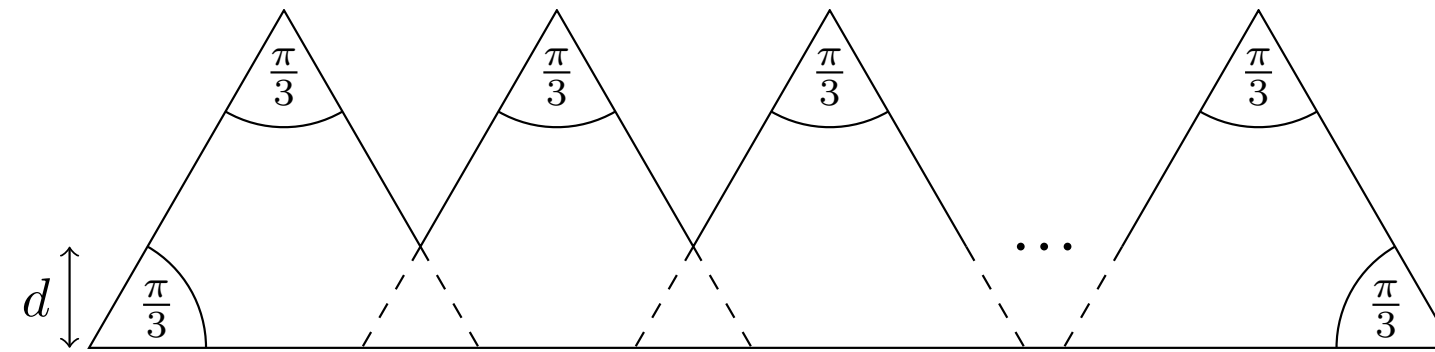
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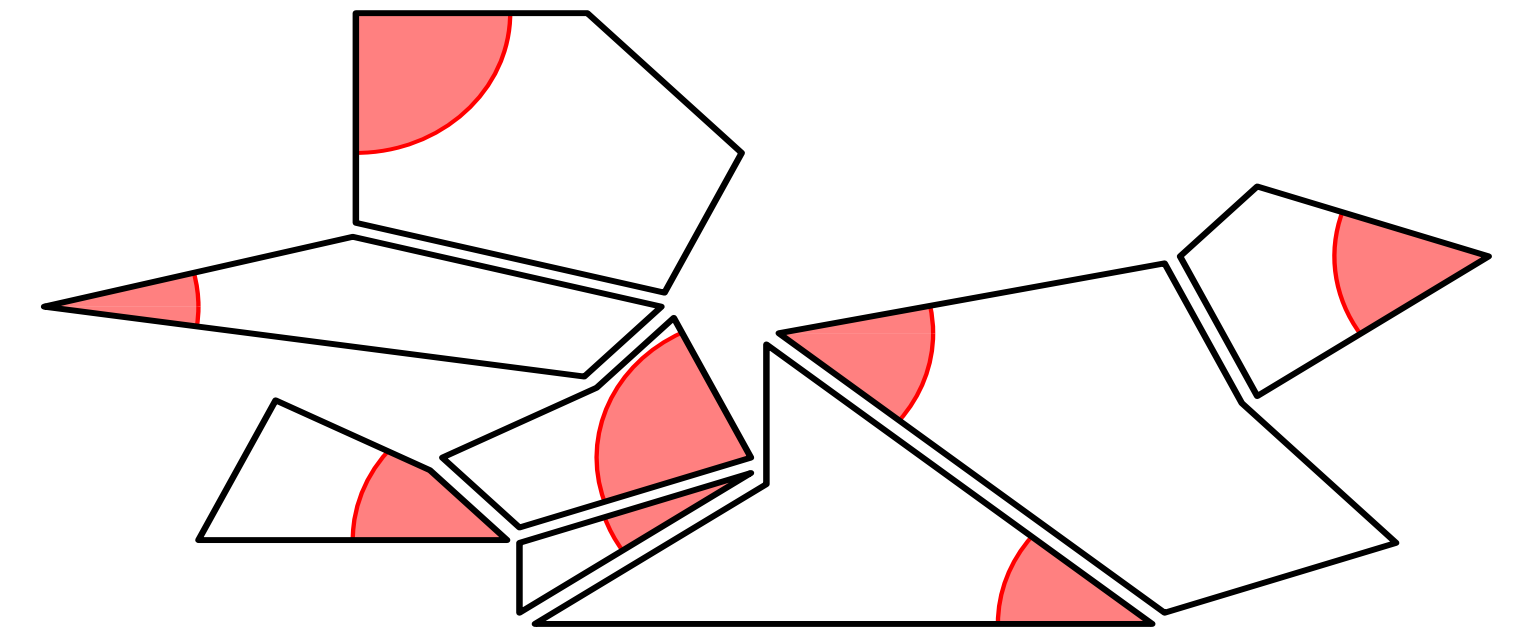
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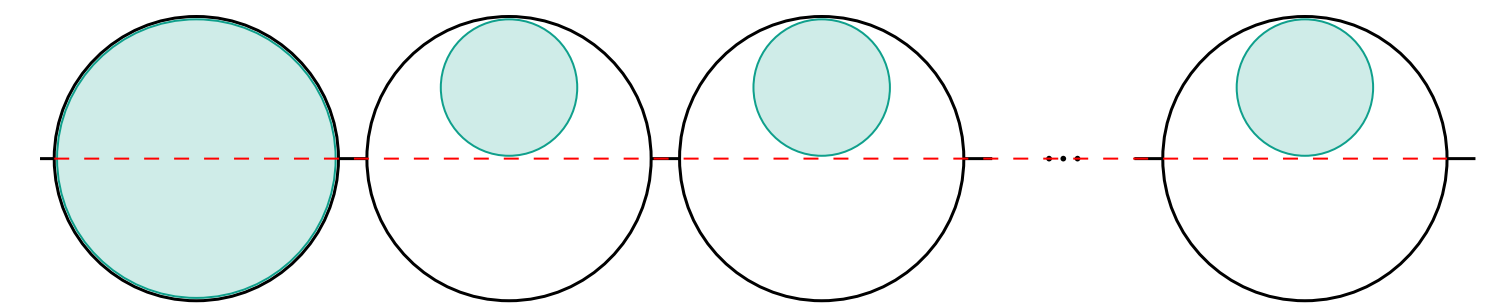


A dual problem deals with independent circle packing.

We determined a duality gap of

$$\approx 4.662$$

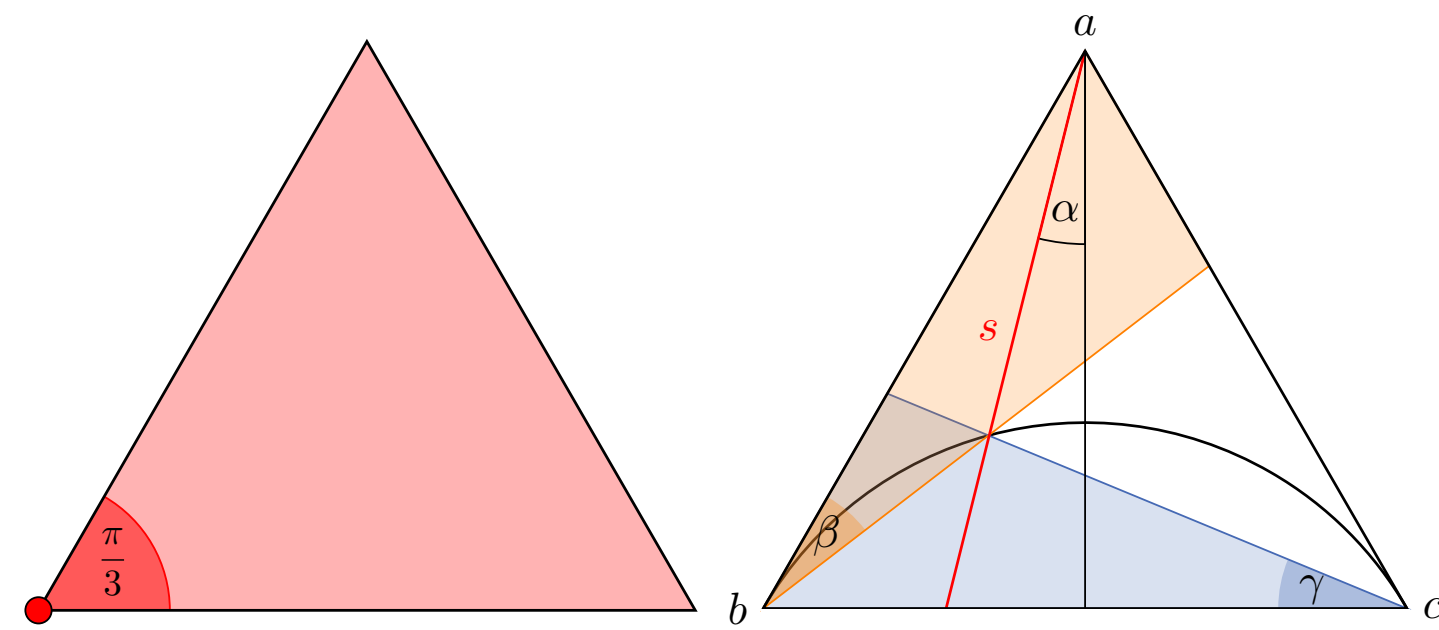
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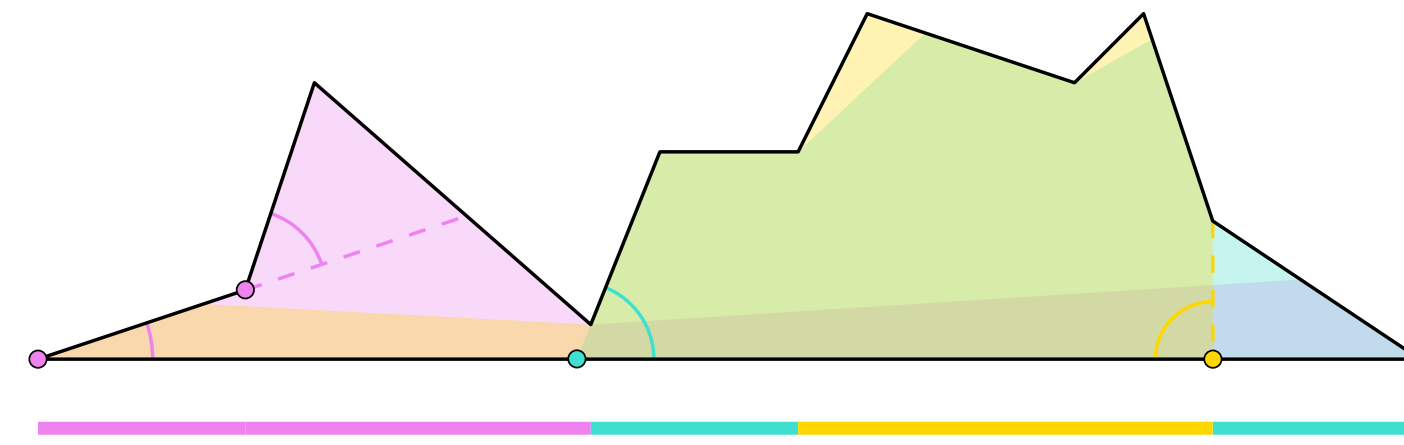
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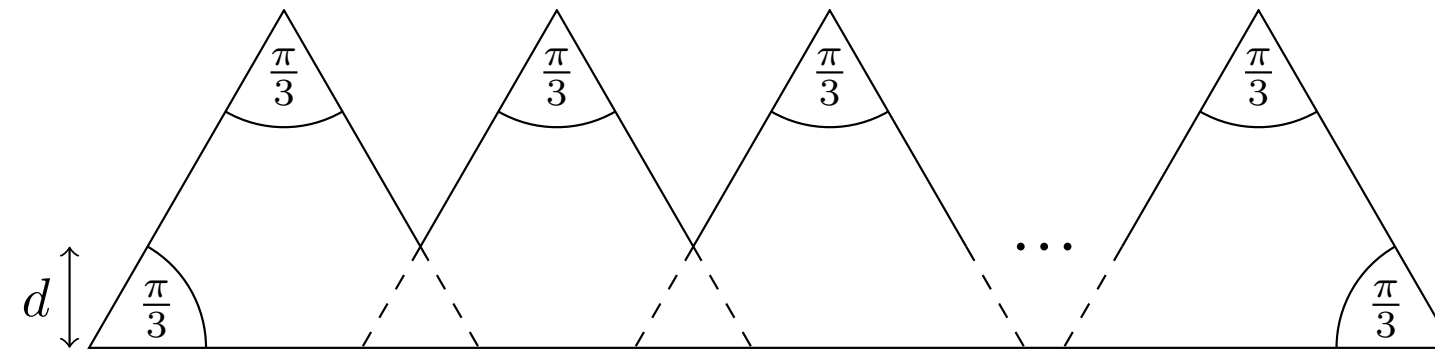
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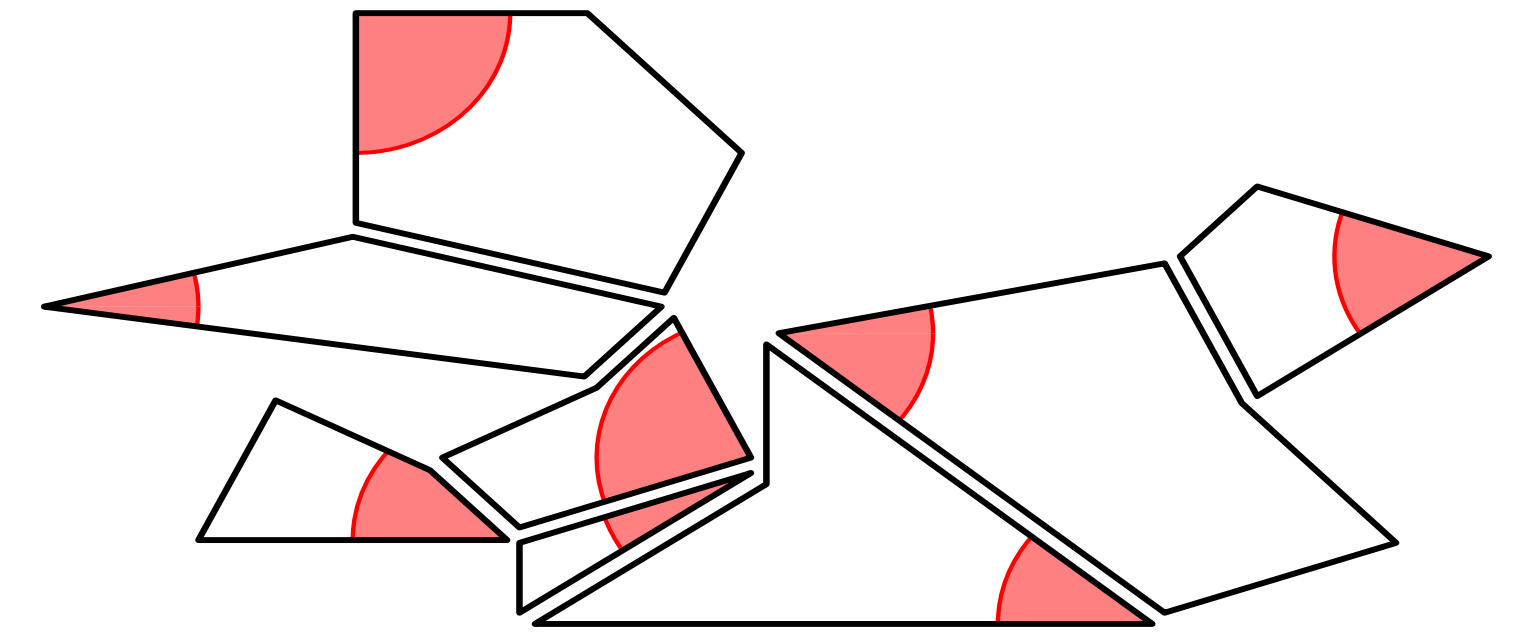
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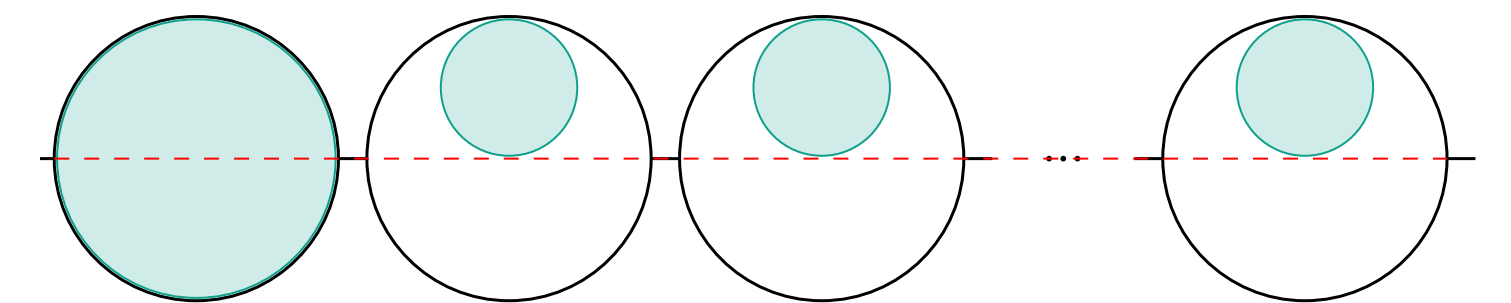


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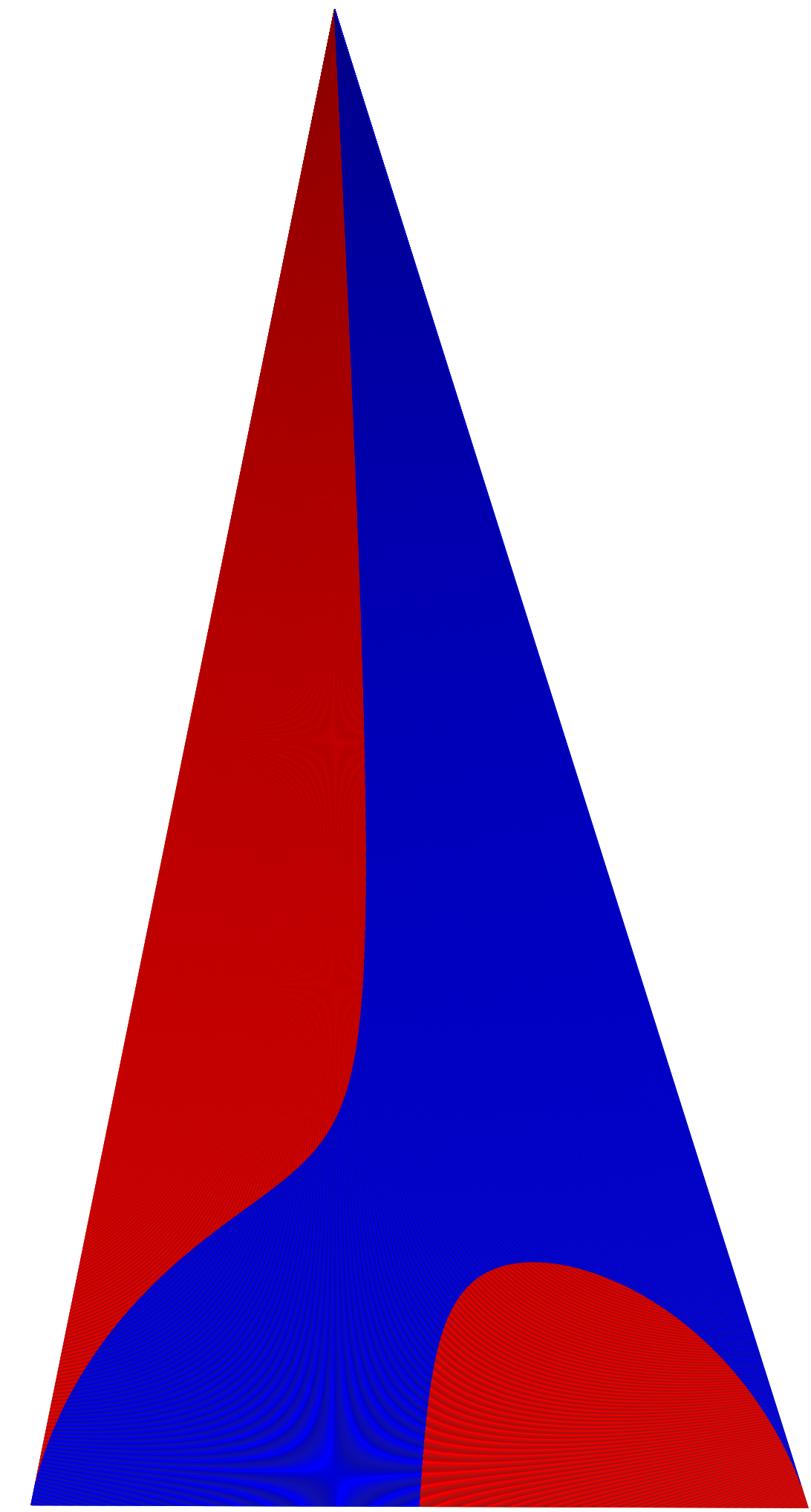
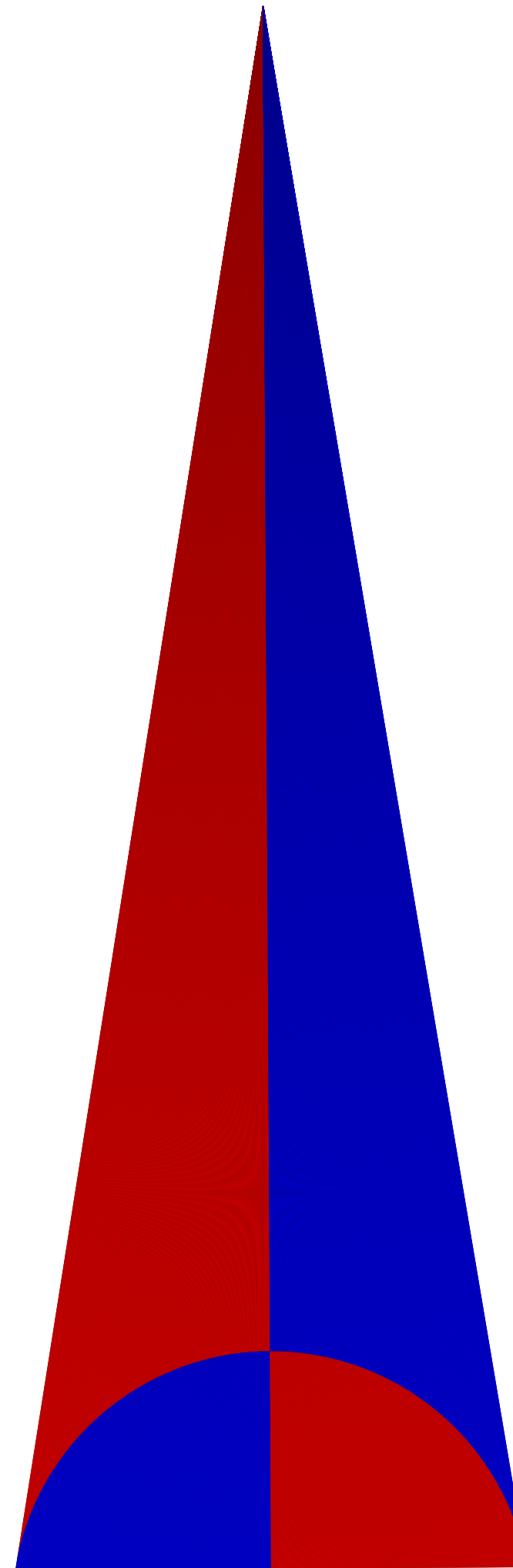
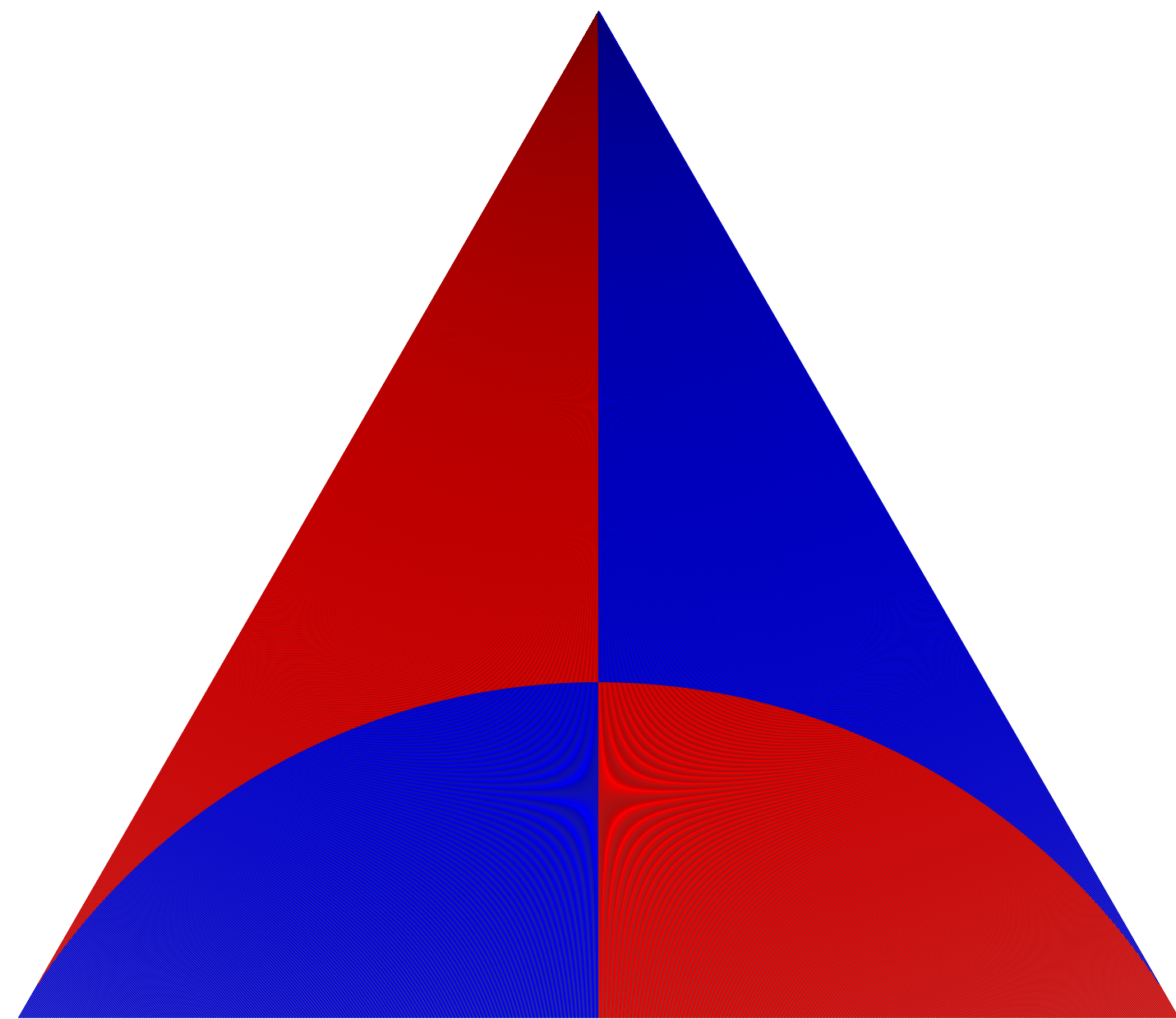
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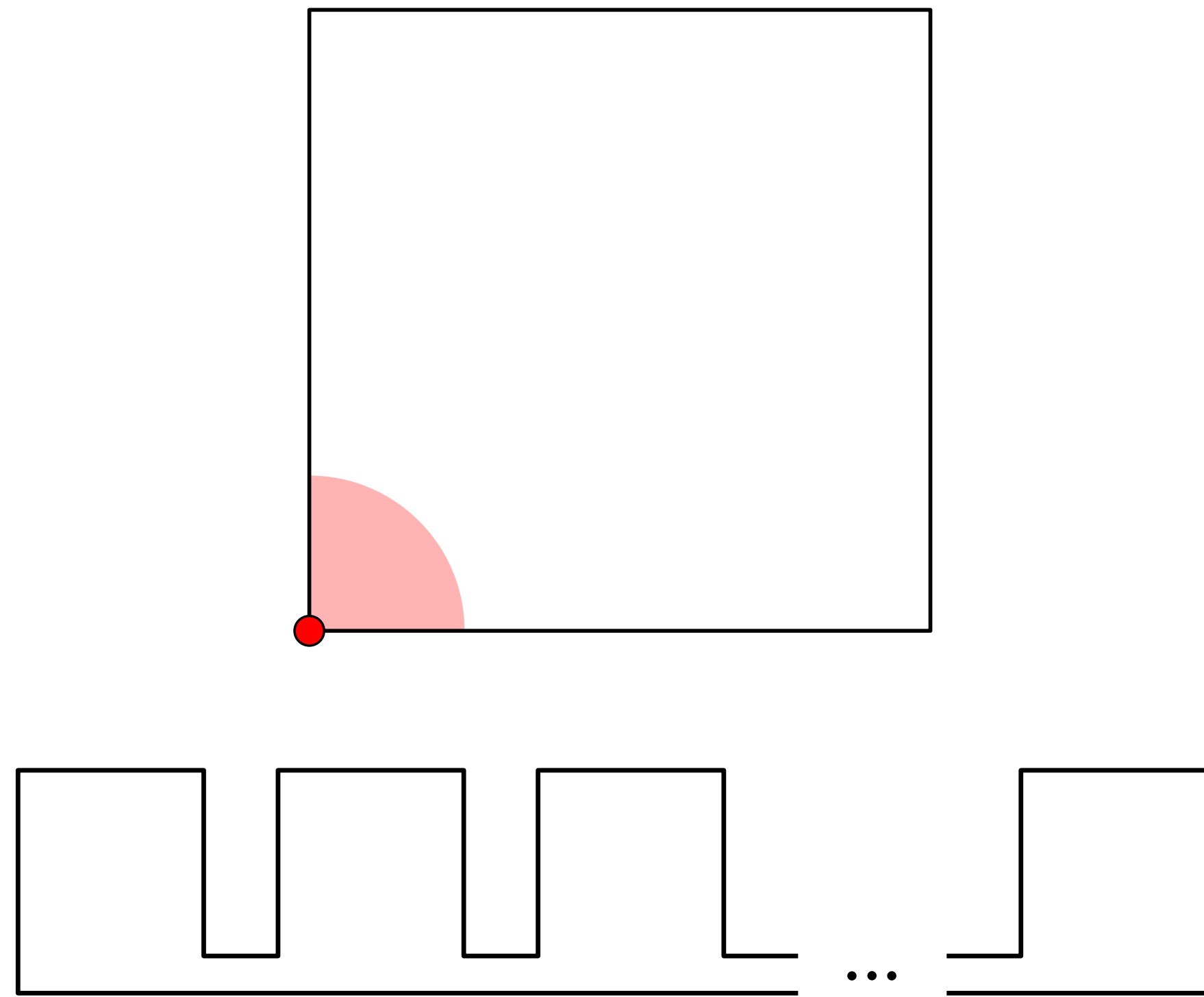




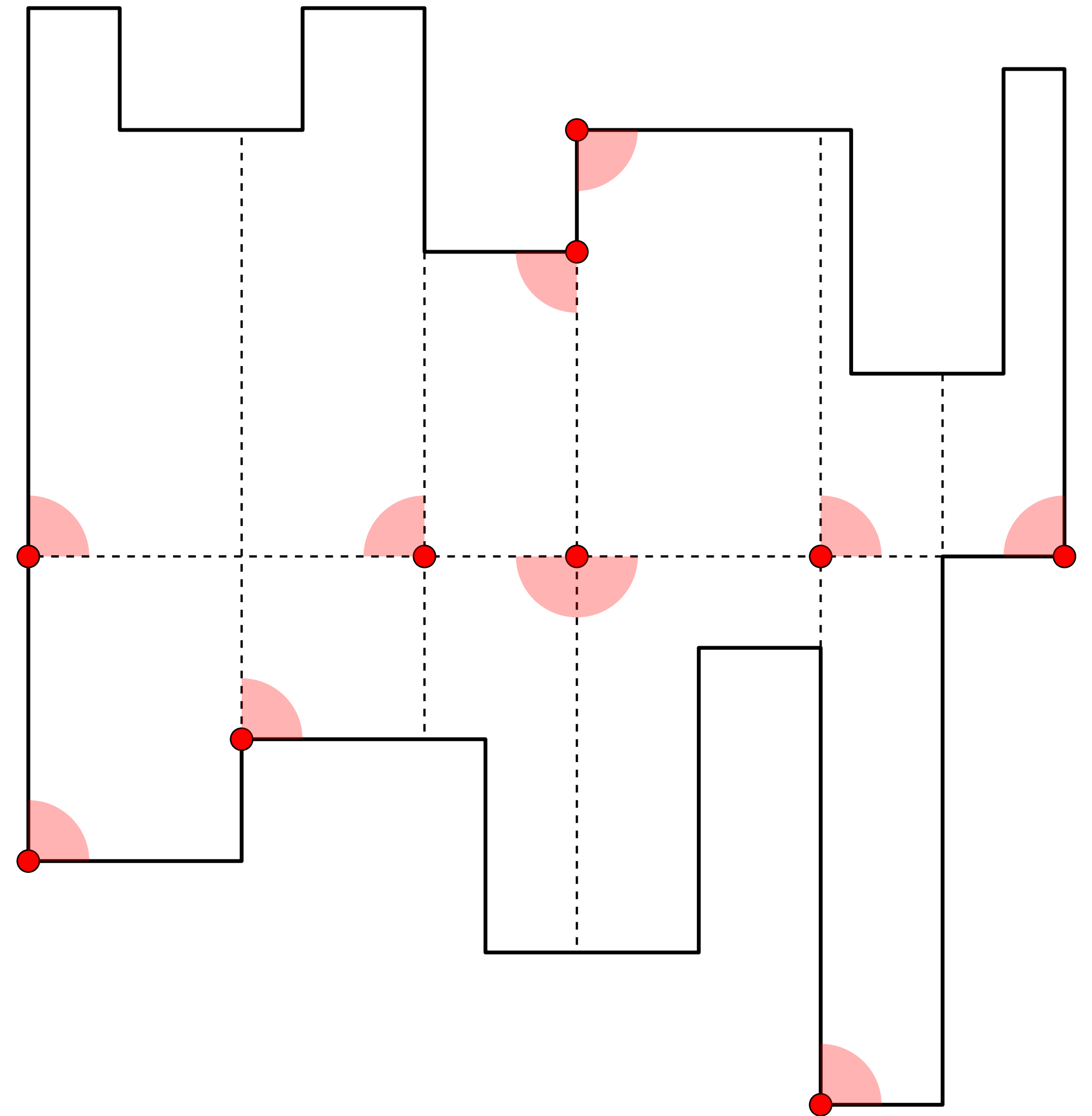
# Future Work



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Assumed tight upper bound:  $\left\lfloor \frac{n}{4} \right\rfloor \frac{\pi}{2}$

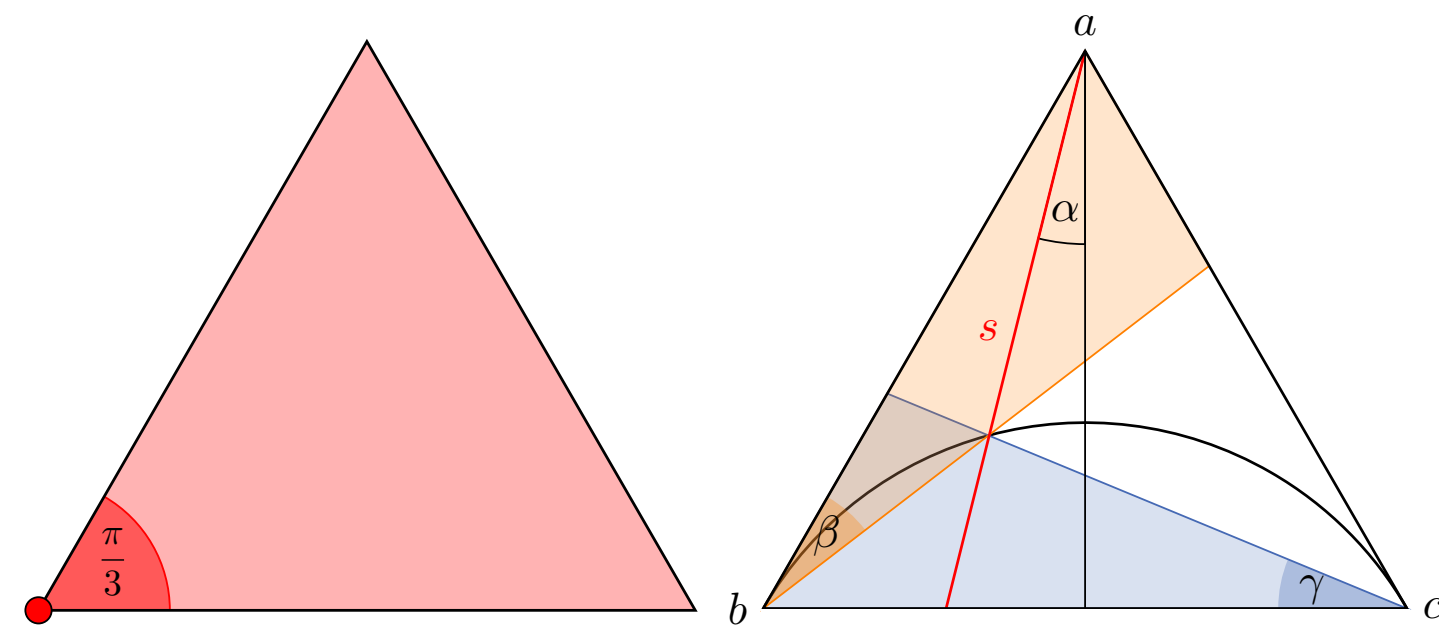




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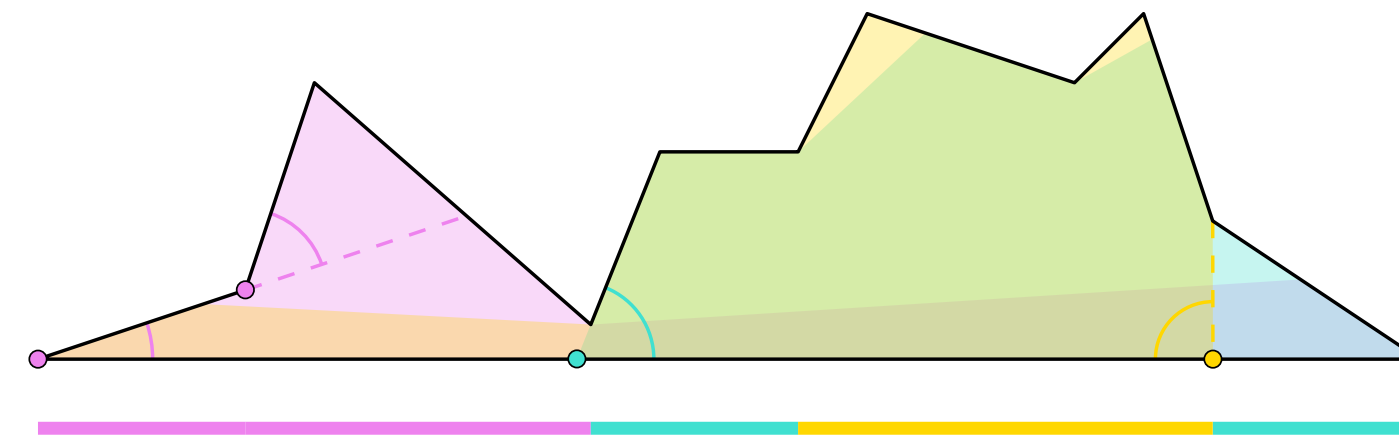
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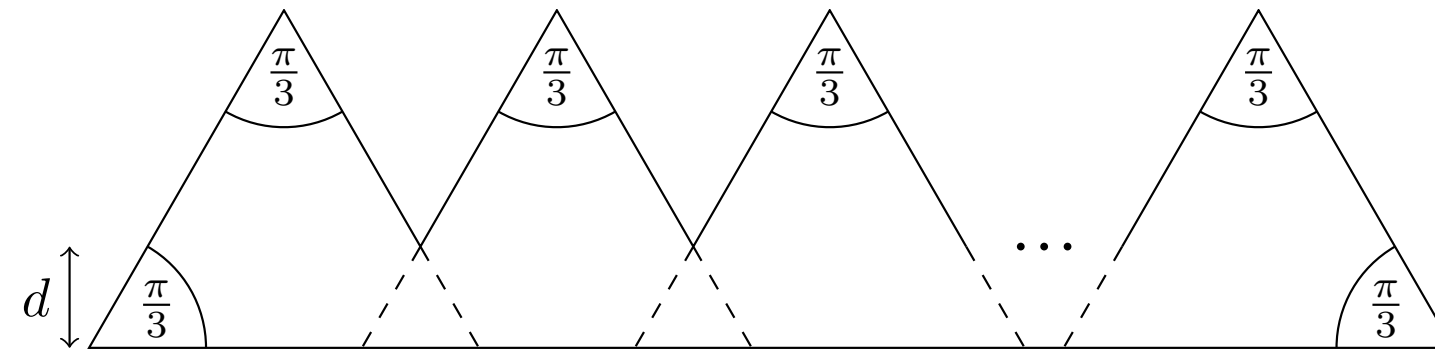
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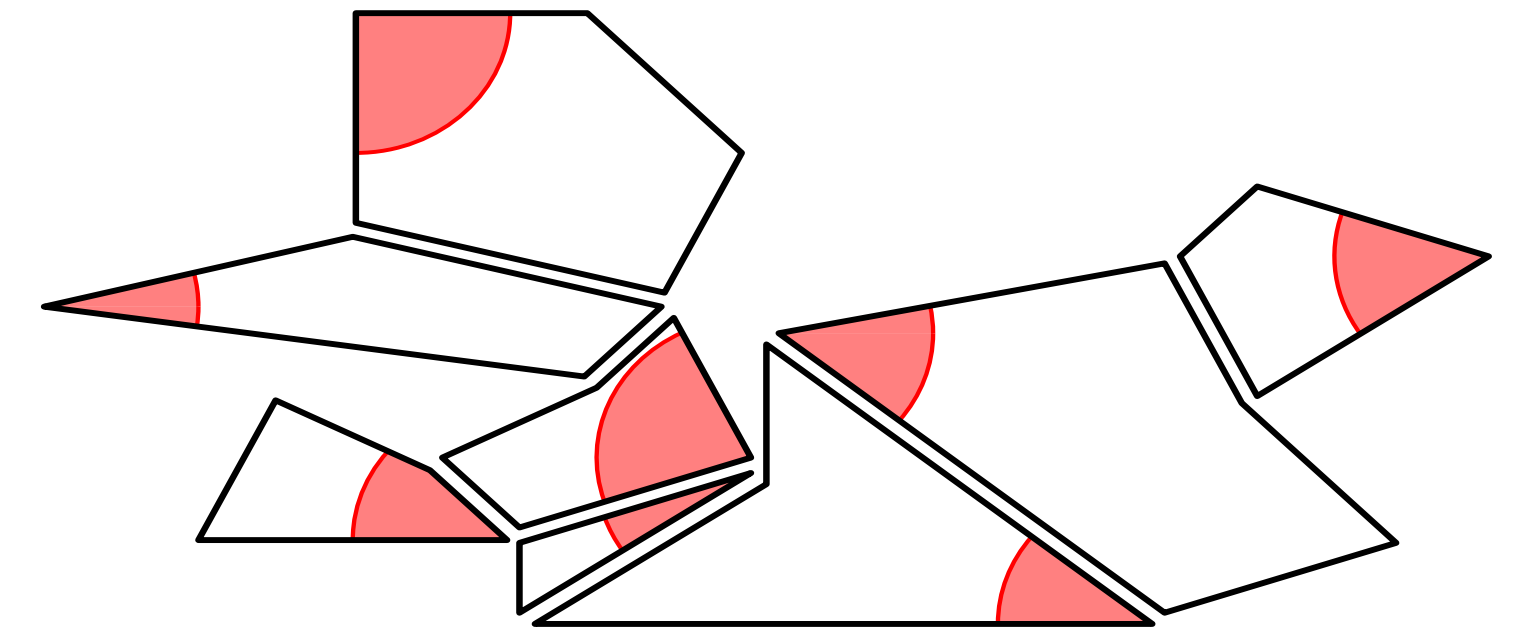
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